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Optimal Selection of the Coefficient Matrix in State-Dependent Control Methods

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This paper presents a state-dependent method for solving nonlinear, control-affine, finite-time optimal control problems. The problem is set in a state-dependent form and solved as a sequence of time-varying linear quadratic regulators, for which an approach based on state transition matrices has been developed. The novelty of the method consists in the possibility of managing a number of different state-dependent coefficient factorizations of the uncontrolled dynamics, for use with various state-dependent control methods. At each iteration, a nonlinear programming problem is solved to select the state-dependent combination that maximizes the controllability, so minimizing the control effort and, likely, the objective function. Test problems are presented to show the effectiveness of the method.

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I. Introduction

Optimal control problems are solved with indirect or direct methods. The former stems from the calculus of variations\(^1\),\(^2\) the latter use a nonlinear programming (NLP) optimization.\(^3\),\(^4\) Both methods require the solution of a complex set of equations (Euler–Lagrange differential equations or Karush–Kuhn–Tucker algebraic equations, respectively) for which iterative numerical methods are used. These iterative procedures implement some form of Newton’s method to find the zeros of a nonlinear function. They are initiated by providing an initial guess solution. Guessing an appropriate initial solution is not trivial, and requires a deep knowledge of the problem at hand. In indirect methods, the initial value of the Lagrange multiplier has to be provided, whose lack of physical meaning makes it difficult to formulate a good guess. In direct methods, the initial trajectory and control have to be guessed at discrete points over the whole time mesh.

Due to their inherent complexity, optimal control problems are often solved within simplified schemes, where some form of approximation is carried out. These methods allow circumventing the original, nonlinear Euler–Lagrange equations, and thus they do not require guessing the initial Lagrange multipliers. Nonetheless, the price to pay is the loss of optimality, since suboptimal solutions are derived. State-dependent methods\(^5\),\(^6\) belong to this category. The state-dependent Riccati equations (SDRE)\(^7\)–\(^10\) is likely the most known example of approximated method due to its simplicity and effectiveness in many applications\(^11\)–\(^16\). This method treats the original infinite-horizon, nonlinear optimal control problem as an infinite-horizon linear-quadratic regulator (LQR), pointwise. A number of LQR problems are solved sequentially at each time instant in which the time domain is discretized. This is done by using state-dependent matrices, which are evaluated pointwise at each time step. With the SDRE the closed-loop control is treated, the control law being function of the present state. The SDRE method can also be used to solve finite-horizon optimal control; one way consists in choosing the state-dependent matrices as functions of the time-to-go.\(^17\)

An alternative approach for the finite-horizon optimal control is represented by the approximated sequence of Riccati equations (ASRE)\(^18\),\(^19\). This method requires again representing the dynamics and the objective function through state-dependent coefficient (SDC) matrices. A number of finite-horizon, time-varying linear-quadratic regulator (TVLQR) problems are solved iteratively. The SDC matrices are evaluated with the solution at the previous iteration. It is worth remarking the differences between SDRE and ASRE. With ASRE, the open-loop optimal guidance problem is treated, and the end-to-end solution is found iteratively. That is, as the SDRE solves a sequence of pointwise LQR at different times labels \((t_k, t_{k+1}, \ldots)\), the ASRE iterates on the solution defined over the whole time domain \([t_i, t_f]\); \(t_i, t_f\) are the initial, final times, respectively. In brief, i) the SDRE solves
the infinite-horizon, closed-loop optimal control \textit{sequentially} by using pointwise LQR; ii) the ASRE solves the finite-horizon, open-loop optimal control \textit{iteratively} by using TVLQR. The TVLQR can be solved with a state transition matrix approach, so avoiding dealing with the differential Riccati equation.

The ASRE method can be implemented to solve a variety of problems.\textsuperscript{20–22} To initiate the algorithm, a certain state-dependent representation for the dynamics must be provided. As such, the degree of freedom represented by the initial guess for the Lagrange multiplier is replaced by the choice of the system factorization. Indeed, although one would expect that different SDC matrices lead to the same result, this is not the case, as it is clearly shown in this paper. Moreover, when different SDC matrices are provided, different results in terms of system trajectory and cost function are found. In principle, one has to run the algorithm a number of times, with different state-dependent representations, and choose the best, feasible solution \textit{a posteriori}.

In this paper the focus is, instead, on modifying the ASRE algorithm to allow managing a set of SDC factorizations in a simultaneous way. Given a set of parent SDC matrices, these are optimally combined to find the one that improves the algorithm convergence, the cost function, and the system controllability. Indeed, it is shown that the latter feature can be directly linked to the solution cost. The higher is the controllability, the less control effort the system will use to enforce the boundary conditions and to minimize the cost function. This is achieved by maximizing the controllability Gramian at each algorithm iteration, which involves solving a NLP optimization problem. More specifically, the minimum singular value of the Gramian is maximized. It is shown that this is the most suitable indicator for the purposes of this paper. In \textsuperscript{23}, an approach to manage multiple SDC factorizations in the SDRE is proposed, where the vector of weighting coefficients is found by minimising the objective function at each time step.

The remainder of the paper is organized as follows. In Section II the background notions are recalled and the motivations are given. The solution method employing the state transition matrix is presented, and implementation issues are discussed. In Section III, the controllability for state-dependent systems is treated, and the metric selected to measure the controllability is discussed. In Section IV, the modified ASRE (MASRE from now on) is formulated and implemented. Section V presents three application cases used to validate the MASRE concept. The achieved results are discussed and final remarks are given in Section VI.
II. Background

An optimal control problem can be stated in a variety of forms, which differ in terms of generality. The most general definition accommodates path constraints, variable final time, control saturation, interior-point constraints, etc. The problems treated in this paper are nonautonomous, nonlinear in the state, and affine (i.e., linear) in the control. The initial state is supposed given, and the final state can be either (in part) specified or unknown. The time span in which the problem is studied is fixed, and both the states and controls are unconstrained.

II.A. Statement of the problem

Given a set of $n$ first-order differential equations

$$\dot{x} = f(x, t) + g(x, t) u,$$  \hspace{1cm} (1)

with $f : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^n$ and $g : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^{n \times m}$, the $m$ control functions $u(t)$ must be determined within initial, final time $t_i, t_f$, such that the performance index

$$J = \varphi(z(t_f), t_f) + \int_{t_i}^{t_f} L(x, u, t) \, dt,$$  \hspace{1cm} (2)

is minimized; $L : \mathbb{R}^{n+m+1} \rightarrow \mathbb{R}$, $z = (x_{q+1}, \ldots, x_n)$ is the vector of state components not specified at final time ($0 \leq q \leq n$), and $\varphi : \mathbb{R}^{(n-q)+1} \rightarrow \mathbb{R}$. The initial condition is simply

$$x(t_i) = x_i,$$  \hspace{1cm} (3)

whereas the final condition is allowed to take three different forms. These specify the hard constrained problem (HCP), the soft constrained problem (SCP), and the mixed constrained problem (MCP), with the final state fully specified, not specified, and partly specified, respectively. The final conditions for each problem, as well as the term $\varphi$ in (2) are defined in Table 1.

Table 1. Definition of final state and costate conditions for the three problems (in MCP $j = 1, \ldots, q < n$ and $k = q + 1, \ldots, n$).

<table>
<thead>
<tr>
<th>Problem</th>
<th>$\varphi$</th>
<th>Final state</th>
<th>Final costate</th>
</tr>
</thead>
<tbody>
<tr>
<td>HCP</td>
<td>not defined</td>
<td>$x(t_f) = x_f$</td>
<td>$\lambda(t_f)$ free</td>
</tr>
<tr>
<td>SCP</td>
<td>$\varphi(x(t_f), t_f)$</td>
<td>$x(t_f)$ free</td>
<td>$\lambda(t_f) = \partial \varphi / \partial x$</td>
</tr>
<tr>
<td>MCP</td>
<td>$\varphi(x_k(t_f), t_f)$</td>
<td>$x_j(t_f) = x_{j,f}$</td>
<td>$\lambda_k(t_f) = \partial \varphi / \partial x_k$</td>
</tr>
</tbody>
</table>

Given the Hamiltonian $H(x, \lambda, u, t) = L(x, u, t) + \lambda^T [f(x, t) + g(x, t)u]$, the problem
consists in finding a solution to the Euler–Lagrange equations,
\[ \dot{x} = \frac{\partial H}{\partial \lambda}, \quad \dot{\lambda} = -\frac{\partial H}{\partial x}, \quad \frac{\partial H}{\partial u} = 0, \]  
(4)

where \( \lambda \) is the vector of costates. The differential-algebraic system (4), together with the initial conditions (3) and the final conditions (for both states and costates) in Table 1, defines a differential-algebraic two-point boundary value problem (TPBVP) whose solution provides the functions \( x(t), \lambda(t), u(t), t \in [t_i, t_f] \).

II.B. The ASRE method

Suppose that \( f \) and \( g \) are at least continuous in an open set \( \Omega \in \mathbb{R}^{n+1} \), and that \( f(0, t) = 0, \forall t \in \mathbb{R} \). Under these conditions, the dynamics (1) can be rewritten in the form
\[ \dot{x} = A(x, t) x + B(x, t) u \]  
(5)

where \( A : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^{n \times n} \) is found by mathematical factorization (it is not unique when \( n > 1 \)) and \( B : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^{n \times m} \). Let also the objective function (2) be redefined in the quadratic-like form
\[ J = \frac{1}{2} z(t_f)^T S(t_f) z(t_f) + \frac{1}{2} \int_{t_i}^{t_f} \left[ x^T Q(x, t) x + u^T R(x, u, t) u \right] dt, \]  
(6)

where \( Q : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^{n \times n} \) and \( R : \mathbb{R}^{n+m+1} \rightarrow \mathbb{R}^{m \times m} \). In analogy with \( \varphi \) in Table 1, \( z \) and \( S \) are not defined in HCP, \( z = (x_{q+1}, \ldots, x_n)^T \) and \( S : \mathbb{R}^{(n-q) \times (n-q)+1} \rightarrow \mathbb{R}^{(n-q) \times (n-q)} \) in MCP, and \( z = (x_1, \ldots, x_n)^T \) and \( S : \mathbb{R}^{(n \times n)+1} \rightarrow \mathbb{R}^{n \times n} \) in SCP. The nonlinear dynamics (5) and the performance index (6), together with the initial condition (3) and the final conditions in Table 1, define an optimal control problem, which is equivalent to that set up by (1)–(2). Yet, this new problem is solved iteratively by reducing (5)–(6) to a series of TVLQR. These are defined by evaluating the SDC matrices \( A, B, Q, R, \) and \( S \) using the solution at the previous iteration. This method is known as the approximating sequence of Riccati equations (or ASRE) method.\textsuperscript{18,19}

The initial step consists in solving Problem 0, defined by
\[ \dot{x}^{(0)} = A^{(0)}(t) x^{(0)} + B^{(0)}(t) u^{(0)}, \]  
\[ J = \frac{1}{2} (z^{(0)}(t_f))^T S^{(0)} z^{(0)}(t_f) + \frac{1}{2} \int_{t_i}^{t_f} \left[ (x^{(0)})^T Q^{(0)}(t) x^{(0)} + (u^{(0)})^T R^{(0)}(t) u^{(0)} \right] dt, \]  
(8)
where

$$A^{(0)}(t) = A(x_i, t), \quad B^{(0)}(t) = B(x_i, t), \quad Q^{(0)}(t) = Q(x_i, t), \quad R^{(0)}(t) = R(x_i, 0, t), \quad S^{(0)} = S(z_i, t_f).$$

Problem 0 is a standard TVLQR as the arguments of the state-dependent matrices are all given except for the time (it is a LQR if (5) is not time-dependent). This problem is solved to yield $x^{(0)}(t)$ and $u^{(0)}(t)$, $t \in [t_i, t_f]$. At a generic, subsequent iteration, Problem $k$ is defined by

$$\dot{x}^{(k)} = A^{(k)}(t) \, x^{(k)} + B^{(k)}(t) \, u^{(k)},$$

$$J = \frac{1}{2}(z^{(k)}(t_j))^T S^{(k)} z^{(k)}(t_f) + \frac{1}{2} \int_{t_i}^{t_f} \left[ (x^{(k)}(t))^T Q^{(k)}(t) x^{(k)}(t) + (u^{(k)}(t))^T R^{(k)}(t) u^{(k)}(t) \right] \, dt,$$

where

$$A^{(k)}(t) = A(x^{(k-1)}(t), t), \quad B^{(k)}(t) = B(x^{(k-1)}(t), t), \quad Q^{(k)}(t) = Q(x^{(k-1)}(t), t), \quad R^{(k)}(t) = R(x^{(k-1)}(t), u^{(k-1)}(t), t), \quad S^{(k)} = S(z^{(k-1)}(t_f), t_f).$$

This is inherently a TVLQR regardless of the time-dependence in (5) (note that $x^{(k-1)}$ and $u^{(k-1)}$ are the solutions of Problem $k - 1$). Solving Problem $k$ yields $x^{(k)}(t)$ and $u^{(k)}(t)$, $t \in [t_i, t_f]$. Iterations continue until a certain convergence condition is satisfied. In the present implementation of the algorithm, the convergence is reached when

$$\varepsilon = \|x^{(k)} - x^{(k-1)}\|_\infty = \max_{t \in [t_i, t_f]} \{ |x^{(k)}_j(t) - x^{(k-1)}_j(t)|, j = 1, \ldots, n \} \leq \text{tol}$$

where ‘tol’ is a prescribed tolerance; i.e., iterations terminate when the difference between each component of the state, evaluated for all times, changes by less than ‘tol’ between two consecutive iterations. The sequence of solutions $x^{(k)}$, $u^{(k)}$ is proven to converge to the solution of the original problem (5)–(6) provided that $A(x, t)$ and $B(x, t)$ are Lipschitz continuous with respect to their arguments.\(^{18}\)

II.C. Solution of the TVLQR by the state transition matrix

The sequence of TVLQR is solved by exploiting the structure of their Euler–Lagrange equations, so avoiding dealing with the matrix differential Riccati equation. This approach, in part, is described in [2], and differs from that implemented in [18, 19]. Suppose the following dynamics are given,

$$\dot{x} = A(t) \, x + B(t) \, u,$$
together with the quadratic objective function

\[
J = \frac{1}{2}(z(t_f))^T S z(t_f) + \frac{1}{2} \int_{t_i}^{t_f} [x^T Q(t) x + u^T R(t) u] \, dt,
\]

(15)

where \( Q, \ S \) and \( R \) are positive semi-definite and positive definite time-varying matrices, respectively. The necessary conditions (4) for this problem read

\[
\dot{x} = A(t) x + B(t) u, \quad (16)
\]

\[
\dot{\lambda} = -Q(t) x - A^T(t) \lambda, \quad (17)
\]

\[
0 = R(t) u + B^T(t) \lambda. \quad (18)
\]

From equation (18) it is possible to get

\[
u = -R^{-1}(t) B^T(t) \lambda, \quad (19)\]

which can be substituted into (16)–(17) to yield

\[
\begin{pmatrix}
\dot{x} \\
\dot{\lambda}
\end{pmatrix} =
\begin{bmatrix}
A(t) & -B(t) R^{-1}(t) B^T(t) \\
-Q(t) & -A^T(t)
\end{bmatrix}
\begin{pmatrix}
x \\
\lambda
\end{pmatrix}.
\]

(20)

Since (20) is a system of linear differential equations, the solution can be written as

\[
x(t) = \phi_{xx}(t_i, t) x_i + \phi_{x\lambda}(t_i, t) \lambda_i, \quad (21)
\]

\[
\lambda(t) = \phi_{\lambda x}(t_i, t) x_i + \phi_{\lambda\lambda}(t_i, t) \lambda_i, \quad (22)
\]

where \( x_i, \lambda_i \) are the initial state, costate, respectively, and the functions \( \phi_{xx}, \phi_{x\lambda}, \phi_{\lambda x}, \) and \( \phi_{\lambda\lambda} \) are the components of the state transition matrix, which can be found by integrating the following dynamics

\[
\begin{bmatrix}
\dot{\phi}_{xx} & \dot{\phi}_{x\lambda} \\
\dot{\phi}_{\lambda x} & \dot{\phi}_{\lambda\lambda}
\end{bmatrix} =
\begin{bmatrix}
A(t) & -B(t) R^{-1}(t) B^T(t) \\
-Q(t) & -A^T(t)
\end{bmatrix}
\begin{bmatrix}
\phi_{xx} & \phi_{x\lambda} \\
\phi_{\lambda x} & \phi_{\lambda\lambda}
\end{bmatrix},
\]

(23)

with initial conditions

\[
\phi_{xx}(t_i, t_i) = \phi_{\lambda\lambda}(t_i, t_i) = I_{n \times n}, \quad \phi_{x\lambda}(t_i, t_i) = \phi_{\lambda x}(t_i, t_i) = 0_{n \times n}. \quad (24)
\]

If both \( x_i \) and \( \lambda_i \) were given, it would be possible to compute \( x(t) \) and \( \lambda(t) \) through (21)–(22), and therefore the optimal control function \( u(t) \) with (19). As only \( x_i \) is given, the
issue is computing $\lambda_i$ for the three problems defined above.

**II.C.1. Hard Constrained Problem**

In a HCP ($x_f$ fully given, $S$ not defined), the value of $\lambda_i$ can be found by writing (21) at final time

$$x_f = \phi_{xx}(t_i, t_f) x_i + \phi_{x\lambda}(t_i, t_f) \lambda_i,$$

and solving for $\lambda_i$; i.e.,

$$\lambda_i(x_i, x_f, t_i, t_f) = \phi_{xx}^{-1}(t_i, t_f) [x_f - \phi_{xx}(t_i, t_f) x_i].$$

**II.C.2. Soft Constrained Problem**

In a SCP ($x_f$ not specified, $S$ a $n \times n$ positive definite matrix), the transversality condition $\lambda(t_f) = \partial \varphi / \partial x$ in Table 1 applied to (15) reads

$$\lambda(t_f) = S x(t_f),$$

which can be used to find $\lambda_i$. This is done by writing (21)–(22) at final time and using (27),

$$x(t_f) = \phi_{xx}(t_i, t_f) x_i + \phi_{x\lambda}(t_i, t_f) \lambda_i,$$

$$S x(t_f) = \phi_{x\lambda}(t_i, t_f) x_i + \phi_{\lambda\lambda}(t_i, t_f) \lambda_i.$$

Equations (28)–(29) represent a linear algebraic system of $2n$ equations in the $2n$ unknowns $\{x(t_f), \lambda_i\}$. The system can be solved by substitution to yield

$$\lambda_i(x_i, t_i, t_f) = (\varphi_{\lambda\lambda}(t_i, t_f) - S(t_f) \phi_{x\lambda}(t_i, t_f))^{-1} [S(t_f) \phi_{xx}(t_i, t_f) - \phi_{x\lambda}(t_i, t_f)] x_i.$$

**II.C.3. Mixed Constrained Problem**

Let the state be decomposed as $x = (y, z)$, where $y = (x_1, \ldots, x_q)$ are the the $q$ known components at final time, whereas $z = (x_{q+1}, \ldots, x_n)$ are the remaining $n-q$ free components at final time. The costate is decomposed accordingly, i.e., $\lambda = (\xi, \eta)$, and $S$ is $(n-q) \times (n-q)$. The MCP is solved by partitioning the state transition matrix in a suitable form such that,
at final time, equations (21)–(22) read
\[
\begin{pmatrix}
  \mathbf{y}(t_f) \\
  \mathbf{z}(t_f)
\end{pmatrix} =
\begin{bmatrix}
  \phi_{yy} & \phi_{yz} \\
  \phi_{zy} & \phi_{zz}
\end{bmatrix}
\begin{pmatrix}
  \mathbf{y}_i \\
  \mathbf{z}_i
\end{pmatrix} +
\begin{bmatrix}
  \phi_{y\xi} & \phi_{y\eta} \\
  \phi_{z\xi} & \phi_{z\eta}
\end{bmatrix}
\begin{pmatrix}
  \mathbf{\xi}_i \\
  \mathbf{\eta}_i
\end{pmatrix},
\]
(31)

\[
\begin{pmatrix}
  \mathbf{\xi}(t_f) \\
  \mathbf{\eta}(t_f)
\end{pmatrix} =
\begin{bmatrix}
  \phi_{\xi\gamma} & \phi_{\xi z} \\
  \phi_{\gamma y} & \phi_{\gamma z}
\end{bmatrix}
\begin{pmatrix}
  \mathbf{y}_i \\
  \mathbf{z}_i
\end{pmatrix} +
\begin{bmatrix}
  \phi_{\xi\xi} & \phi_{\xi\eta} \\
  \phi_{\gamma\xi} & \phi_{\gamma\eta}
\end{bmatrix}
\begin{pmatrix}
  \mathbf{\xi}_i \\
  \mathbf{\eta}_i
\end{pmatrix},
\]
(32)

where the dependence on \( t_i, t_f \) is omitted for brevity. From the first row of (31) it is possible to get
\[
\mathbf{\xi}_i = \phi_{yy}^{-1}[\mathbf{y}_f - \phi_{yy}\mathbf{y}_i - \phi_{yz}\mathbf{z}_i] - \phi_{y\gamma}^{-1}\phi_{y\eta}\mathbf{\eta}_i,
\]
(33)

which can be substituted in the second row of (31) to yield
\[
\mathbf{z}(t_f) = [\phi_{zy} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{yy}]\mathbf{y}_i + [\phi_{zz} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{yz}]\mathbf{z}_i + \phi_{z\xi}\phi_{y\xi}^{-1}\mathbf{y}_f + [\phi_{z\eta} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{y\eta}]\mathbf{\eta}_i.
\]
(34)

Equations (33)–(34), together with the transversality condition \( \mathbf{\eta}(t_f) = \mathbf{S}(t_f)\mathbf{z}(t_f) \) (see Table 1), can be substituted in the second row of (32) to compute \( \mathbf{\eta}_i \); i.e.,
\[
\mathbf{\eta}_i(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f) = [\phi_{\gamma\eta}]^{-1}\mathbf{w}(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f),
\]
(35)

where \( \phi_{\gamma\eta} = \phi_{\gamma\xi}\phi_{y\xi}^{-1}\phi_{y\eta} - \phi_{z\eta} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{y\eta} \) and
\[
\mathbf{w}(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f) = [\mathbf{S}(\phi_{zy} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{yy}) - \phi_{yy} + \phi_{\gamma\xi}\phi_{y\xi}^{-1}\phi_{y\eta}]\mathbf{y}_i + [\mathbf{S}(\phi_{zz} - \phi_{z\xi}\phi_{y\xi}^{-1}\phi_{yz}) + \phi_{\gamma\xi}\phi_{y\xi}^{-1}\phi_{y\eta}]\mathbf{z}_i + [\mathbf{S}(\phi_{z\xi}\phi_{y\xi}^{-1}) - \phi_{\gamma\xi}\phi_{y\xi}^{-1}]\mathbf{y}_f.
\]
(36)

Once \( \mathbf{\eta}_i \) is known, the remaining part of the initial costate, \( \mathbf{\xi}_i \), is computed through (33), and therefore the full initial costate is obtained as a function of the initial condition, given final condition, initial, and final time; i.e, \( \mathbf{\lambda}_i(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f) = (\mathbf{\xi}_i(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f), \mathbf{\eta}_i(\mathbf{x}_i, \mathbf{y}_f, t_i, t_f)) \).

II.D. Implementation of the iterative scheme

The ASRE method illustrated in Section II.B has been coded into a numerical framework by implementing the solution scheme shown in Section II.C. The overall flowchart of the algorithm is shown in Figure 1. The original, fully nonlinear problem (1)–(2) is factorized into the state-dependent form (5)–(6), and then solved iteratively with the three-step procedure in the ‘Solver’ box. At each iteration the convergence condition (13) is checked, and either another iteration is performed or the solution is found. As the integration of (23) requires evaluating the SDC matrices for different times, an uniform time grid of \( N \) points is constructed, \( t_i = t_1, t_2, \ldots, t_N = t_f \), and a cubic spline interpolation is used to evaluate
these matrices within the $N-1$ time intervals. The solution is then available only on this grid, i.e., $x_j = x(t_j), \lambda_j = \lambda(t_j), u_j = u(t_j), j = 1, \ldots, N$, and a cubic spline may be again used if the solution over an intermediate time has to be retrieved.

$$\dot{x} = f(x, t) + g(x, t) \, u$$
$$J = \varphi(x(t_1), t_1) + \int_{t_1}^{t_f} L(x, u, t) \, dt$$

$$x = A(x, t) \, x + B(x, t) \, u$$
$$J = \frac{1}{2} x^T S \, z + \frac{1}{2} \int_{t_1}^{t_f} [x^T Q(x, t) x + u^T R(x, u, t) u] \, dt$$

**Problem Statement**

$$\dot{x} = f(x, t) + g(x, t) \, u$$

**Factorization**

$$x = A(x, t) \, x + B(x, t) \, u$$

$$J = \frac{1}{2} x^T S \, z + \frac{1}{2} \int_{t_1}^{t_f} [x^T Q(x, t) x + u^T R(x, u, t) u] \, dt$$

**Problem 0**

Define $A = A^{(0)}, B = B^{(0)}, Q = Q^{(0)}, R = R^{(0)}$ [Eq. (9)]

**Problem $k$**

Define $A = A^{(k)}, B = B^{(k)}, Q = Q^{(k)}, R = R^{(k)}$ [Eq. (12)]

**Iter**

$k = 0$

$k > 0$

**Solver**

Compute: i) $\phi_{xx}, \phi_{x\lambda}, \phi_{\lambda x}, \phi_{\lambda\lambda}$ [Eqs. (23)-(24)]

ii) $\lambda$ [Eq. (26) HCP/Eq. (30) SCP/Eq. (35) MCP]

iii) $x^{(k)}(t), \lambda^{(k)}(t), u^{(k)}(t)$ [Eqs. (21)-(22), (19)]

**Error**

$$\epsilon = ||x^{(k)} - x^{(k-1)}||_\infty$$

$\epsilon < \text{tol}$

$\epsilon > \text{tol}$

Stop

Solution

$$x(t) = x^{(k)}(t), \lambda(t) = \lambda^{(k)}(t), u(t) = u^{(k)}(t)$$

**Figure 1. Flowchart for the implementation of the ASRE scheme.**

### II.E. Dependence on Factorization

The ASRE algorithm, as coded in Figure 1, has been used to solve a number of nonlinear optimal control problems. This algorithm requires factorizing the nonlinear dynamics with a state-dependent matrix. Finding a good factorization can be non-trivial, mostly because, for systems with at least two states, there exists an infinite number of factorizations. Indeed, given a dynamical system

$$\dot{x}_1 = f_1(x_1, x_2),$$
$$\dot{x}_2 = f_2(x_1, x_2),$$

(37)
it is possible, for instance, to use the SDC matrix $A_1(x)$,

$$
\begin{pmatrix}
\dot{x}_1 \\
\dot{x}_2
\end{pmatrix} = \begin{bmatrix}
\frac{f_1(x_1, x_2)}{x_1} & 0 \\
0 & \frac{f_2(x_1, x_2)}{x_2}
\end{bmatrix} \begin{pmatrix}
x_1 \\
x_2
\end{pmatrix},
$$

(38)

or $A_2(x)$,

$$
\begin{pmatrix}
\dot{x}_1 \\
\dot{x}_2
\end{pmatrix} = \begin{bmatrix}
0 & \frac{f_1(x_1, x_2)}{x_2} \\
\frac{f_2(x_1, x_2)}{x_1} & 0
\end{bmatrix} \begin{pmatrix}
x_1 \\
x_2
\end{pmatrix}.
$$

(39)

Not only, as there exists a family of linear combinations,

$$
A_3(\alpha, x) = (1 - \alpha) A_1(x) + \alpha A_2(x), \quad \forall \alpha \in \mathbb{R},
$$

(40)

which produces infinite factorizations of (37). This procedure can be applied to any system with $n > 1$ states. One could think that the factorization choice does not affect the final solution, but actually this is not the general case, as shown by the following, simple example.

II.E.1. Example 1

The problem consists in minimizing

$$
J = \frac{1}{2} \int_{t_i}^{t_f} \left( x_1^2 + x_2^2 + u^2 \right) dt,
$$

(41)

under the dynamics (taken from [24])

$$
\begin{aligned}
\dot{x}_1 &= x_2, \\
\dot{x}_2 &= x_1 x_2 + u,
\end{aligned}
$$

(42)

with boundary conditions $x(t_i) = (2, 1)^T$, $x(t_f) = (4, 1)^T$, and $t_i = 0$, $t_f = 5$.

Two possible parent factorizations of the nonlinear, uncontrolled dynamics (42) are

$$
A_1(x) = \begin{bmatrix}
0 & 1 \\
0 & x_1
\end{bmatrix}, \quad A_2(x) = \begin{bmatrix}
0 & 1 \\
x_2 & 0
\end{bmatrix}.
$$

(43)

The remaining SDC matrices for this problem are $B = [0, 1]^T$ (in (5)) and $Q = I_{2 \times 2}$, $R = 1$ (in (6)); $S$ in (6) is not defined, since Example 1 is an HCP. This problem has been solved with the ASRE scheme by using $A_1$, $A_2$ in (43) and also $A_3$ in (40) with $\alpha = 0.5$. 

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Optimal Selection of the Coefficient Matrix in State-Dependent Control Methods – F.T., M.M., and F.B.-Z.
The three different solutions are shown in Figure 2, and their features (cost and iterations) are summarized in Table 2. The qualitative and, most importantly, quantitative difference among these solutions makes it clear that the SDC factorization of the uncontrolled dynamics plays a crucial role in the present, as well as any other, state-dependent methods.

| Table 2. Features of the three solutions to Example 1 (tol = 10^{-6}). |
|-----------------|---|---|
| A(x)            | J  | Iter |
| A_1             | 52.24 | 12   |
| A_2             | 64.91 | 56   |
| A_3             | 53.13 | 21   |

Figure 2. Trajectories of the three solutions to Example 1.

### III. Controllability of State-Dependent Dynamics

A property that ASRE requires is controllability of the pair \( \{A(x,t), B(x,t)\} \) in (5). A system is controllable if and only if there always exists an admissible control that is able to transfer the state from a condition to another in a finite time. When dealing with state-dependent dynamics, the SDC matrices must be selected such that the state-dependent system is controllable for all the possible trajectories. In [26], necessary and sufficient conditions for the existence of SDC matrices are shown. It would be desirable to choose the factorization that grants the maximum controllability of the state-dependent dynamics, since the solution performance is strictly related to the controllability of the system. An increased controllability leads to a reduction of the control effort and likely to a reduced cost. In this way it could be possible to identify the optimal factorization.
However, the controllability of (5) does not lead to the controllability of (1), as there is no trivial relation between the controllability of \( \{A(x, t), B(x, t)\} \) and that of \( \{f(x, t), g(x, t)\} \).\(^{24}\) The controllability of the state-dependent dynamics is increased with the purpose of easing the convergence of the ASRE method.

A linear system is controllable when the controllability matrix, \( K_c \), is full-rank; i.e.,

\[
\text{rank}(K_c) = \text{rank}\begin{bmatrix} B & AB & A^2B & \cdots & A^{n-1}B \end{bmatrix} = n. \tag{44}
\]

Since \( A, B \) in (5) are in general state- and time-dependent, \( K_c \) is state- and time-dependent as well, and thus the controllability check (44) must be done for all states and times. This is deemed not viable. Another way to study the controllability of state-dependent systems is based on evaluating their distance from being uncontrollable, which can be done by estimating the amplitude of the controllability radius,\(^{28,29}\)

\[
\rho(A, B) = \min_{\lambda \in \mathbb{C}} \sigma_{\min}\left(\begin{bmatrix} A - \lambda I & B \end{bmatrix}\right), \tag{45}
\]

where \( \sigma_{\min}(\cdot) \) returns the smallest singular value of its matrix argument.

In this paper, the notion of controllability Gramian is used.\(^{25}\) This is defined as

\[
P(t_i, t_f) = \int_{t_i}^{t_f} \Phi(t_i, \tau)B(\tau)B^T(\tau)\Phi^T(t_i, \tau)d\tau, \tag{46}
\]

where \( \Phi(t_i, \tau) \) is the solution of

\[
\dot{\Phi}(t, \tau) = A(t)\Phi(t, \tau), \quad \Phi(t, t) = I, \tag{47}
\]

with \( A(t) \) as in (14). A system is controllable if \( P(t_i, t_f) \) is nonsingular; i.e., if \( \det P(t_i, t_f) \neq 0 \). Notice that this condition for controllability is seeking controllability along the entire trajectory of the system, and is not a combination of pointwise controllability for a given set of values of the states. This is, however, a binary condition: it can be used to check whether a system is controllable or not, but not to measure its controllability. As the higher controllability provides the less control effort, the focus is to choose the factorization that maximizes the controllability, and therefore a metric is introduced by defining a figure of merit.

In [10,30,31], the SDC are chosen by maximizing the absolute value of the determinant of \( K_c K_c^T \) (note that \( K_c \) is not necessarily square. When \( P(t_i, t_f) \) is considered, it could seem reasonable to adopt the absolute value of its determinant as a merit figure of controllability. Actually, such a choice would be misleading since it is easy to prove that a nearly singular matrix can show a determinant much greater than zero.\(^{28}\) Alternatively, the minimum
eigenvalue of $P$ may be used to measure the controllability: a matrix is close to singularity when its smallest eigenvalue is close to zero, regardless of the determinant value. However, the eigenvalues suffer of ill-conditioning as small perturbations of a matrix can lead to big variations of its eigenvalues. To overcome these issues, in the remainder, the Gramian singular values have been used. These are positive real numbers that can be used in place of eigenvalues. More specifically, $\sigma_{\text{min}}(P)$ is considered to measure the system controllability.

In Figure 3, the trend of $\sigma_{\text{min}}(P)$ for the solutions to Example 1 is shown, at each iteration, for the three SDC matrices considered. It can be noticed that the case requiring the lowest number of iterations corresponds to $A_1$, which presents also the highest $\sigma_{\text{min}}(P)$. From the other two cases ($A_2$, $A_3$) it can be inferred that the lower $\sigma_{\text{min}}(P)$, the higher number of iterations, and vice versa. Not only, as from Table 2 it can be seen that the highest $\sigma_{\text{min}}(P)$ involves the best (i.e., lowest) value of cost function.

In summary, different SDC representations lead to different behaviors of the ASRE. The factorizations resulting the most controllable are also those entailing the lowest cost functions and numbers of iterations. Given a nonlinear dynamics, the selection of the state-dependent matrix $A(x)$ is a degree of freedom that can be exploited to enforce the best performances in terms of algorithm convergence and control effort.

![Figure 3. Values of $\sigma_{\text{min}}(P)$ during the iterations for the three factorization cases of Example 1.](image)

**IV. A Combined Method**

It has been shown that the SDC of the uncontrolled dynamics affect the algorithm convergence, the number of iterations, the solution trajectory, and the objective function value. In principle, once a nonlinear problem is given, it would be possible to find a number of
different SDC matrices and to run the algorithm in Figure 1 several times. Then, the best solution in terms of cost function could be extracted among those feasible. This means using the ASRE algorithm as a black-box tool, not relying on any notion of controllability.

In this paper the focus is instead on handling all the possible factorizations at once, and to optimally combine them to increase the controllability of the resulting system. Given a set of parent SDC matrices, the spirit is to treat them in the same fashion as $A_1, A_2$ are combined to define $A_3$ in (40). The aims are: 1) to increase the controllability, 2) to ease the algorithm convergence, 3) to improve the solution cost. This is done by maximizing $\sigma_{\text{min}}(P)$ at each iteration, which involves solving a nonlinear programming problem. If needed, the approximate solution so found is used to tackle the TPBVP defined by the original dynamics. It is shown that the approximate solution so found is a good initial guess, and it is only refined to satisfy the Euler–Lagrange equations (4).

IV.A. Managing multiple factorizations

Given a set of $n_f$ parent factorizations of $f(x, t)$ in (1) having the form $f(x, t) = A_i(x, t)x$, $i = 1, \ldots, n_f$, the combination that grants the maximum controllability is looked for. The family of SDC matrices (generated by the parents $A_i$) can be written as

$$A(\alpha, x, t) = \left[ \prod_{j=1}^{n_f-1} (1 - \alpha_j) \right] A_1(x, t) + \sum_{i=2}^{n_f-1} \alpha_{i-1} \left[ \prod_{j=i}^{n_f-1} (1 - \alpha_j) \right] A_i(x, t) + \alpha_{n_f-1} A_{n_f}(x, t),$$

(48)

where $\alpha = (\alpha_1, \ldots, \alpha_{n_f-1})^T$ is a vector of weighting coefficients. The family (48) produces the state-dependent form (5), i.e.,

$$\dot{x} = A(\alpha, x, t)x + B(x, t)u,$$

(49)

if the following inequalities are satisfied,

$$\|\alpha\|_1 \leq 1, \quad \text{and} \quad \alpha_j \geq 0, \ j = 1, \ldots, n_f - 1.$$

(50)

Conditions (50) define the region of existence for $\alpha$, which is required in order to ensure a convex domain (see Figure 4). More precisely, the restrictions in (50) guarantee a one-to-one correspondence between $\alpha$ and $A(\alpha, x, t)$ in (48). As $\alpha$ is a vector of bounded coefficients, $A(\alpha, x, t)$ is Lipschitz continuous when so are the parent factorizations $A_i(x, t)$. Other relations can be formulated in place of (48). In [23], a linear combination is used to define a family of factorizations in the SDRE method; in [10,30,31], such combination is done at scalar level by treating each nonlinear element of the original dynamics independently from the others. Equation (48) is preferred as it is more general. It allows considering...
families of nonlinearities in a unified way. This is done to maintain some physical insight into the problem. Indeed, in typical problems dynamics and nonlinearities are related to vector quantities that in the proposed approach would be treated as vectors and not as independent scalar variables.

\[ \alpha_1 \geq 0 \]

\[ \alpha_2 \geq 0 \]

\[ \parallel \alpha \parallel_1 \leq 1 \]

\[ \alpha_3 \geq 0 \]

Figure 4. Region of existence for \( \alpha \) in one, two, and three dimensions \((n_f = 2, 3, 4\), respectively).  

IV.B. Optimal combination of state-dependent factorizations

The expression (48) allows us to handle a set of SDC matrices simultaneously. This is used to modify the ASRE method to accommodate for \( n_f \) factorizations. However, the iterative method requires the dynamics to be written in the form (5), which can be achieved by specifying the value of \( \alpha = (\alpha_1, \ldots, \alpha_{n_f})^T \) in (49). It is important to recall that the parent factorizations \( A_i \) have been introduced for the purpose of increasing the system controllability. Thus, the weights \( \alpha_j \) have to be chosen such that the minimum singular value of the Gramian associated to (49) is maximized. This is done by solving, at each iteration, the following Nonlinear Programming (NLP) problem

\[
\max_{\alpha} \sigma_{\min}(P(\alpha)) \quad \text{subject to} \quad \begin{cases} 
\parallel \alpha \parallel_1 \leq 1 \\
\alpha_j \geq 0, \ j = 1, \ldots, n_f - 1
\end{cases}
\]

(51)

where \( P(\alpha) \) is the Gramian obtained by plugging (48) into (46)–(47). The degrees of freedom \( \alpha_j \) are thus exploited optimally to improve the existing algorithm. This is the ultimate goal of the MASRE. The price to pay is the solution of the algebraic problem (51) at each iteration. This holds as the state-dependent matrices \( A, B \) in (46)–(47) are evaluated with the solution at the previous step. Indeed, in analogy with Eqs. (9) and (12), the SDC matrices are

\[
A^{(0)}(\alpha, t) = A(\alpha, x_i, t), \quad \text{in Problem 0;} \tag{52}
\]

\[
A^{(k)}(\alpha, t) = A(\alpha, x^{(k-1)}, t), \quad \text{in Problem k.} \tag{53}
\]
IV.C. Implementation of the modified scheme

The ASRE scheme in Figure 1 has been modified to manage multiple factorizations and to choose their optimal combination that yields the best controllability. Figure 5 reports the flow chart of the MASRE algorithm. Once the problem is stated, a number of parent SDC matrices for the uncontrolled dynamics are defined, which are then combined to generate the whole family through (48). The general problem is then defined by (49) and (6). This is treated in the same fashion as in Section II except that a value for \( \alpha \) is needed this time. In order to obtain the optimal value of \( \alpha \), the NLP is solved in the ‘Optimization’ box. The first guess value for \( \alpha \) is the value found at the previous iteration (an initial guess \( \alpha^{(0)} \) is provided to run the NLP the first time). The approximate solution \( (x(t), \lambda(t), u(t)) \) is found once the convergence condition is verified. This condition is the same as in the ASRE, where a single SDC matrix is used. In the MASRE, where the SDC matrix may change at each iteration, the conditions for the convergence of the sequence of Riccati equations may be violated. However, numerical experiments suggest that algorithm convergence is achieved once \( \alpha \) in (48) stabilizes about an optimal value, such that the sequence of Riccati equations approaches convergence with a fixed SDC matrix (see Section V).

IV.D. Refinement of the approximate solution

The approximate solution is expressed through functions evaluated at discrete locations of a time grid made up of \( N \) points (see Section II.D). If needed, this solution can be refined by solving the Euler–Lagrange equations (4). In particular, the approximate solutions can be used as a first guess solutions for a TPBVP solver that deals with Eq. (4) and the boundary conditions as in (3) and Table 1 (see Figure 5). This step is particularly easy when the original problem (2) possesses a quadratic-control objective function. This condition, together with the control affinity in (1), makes it possible to yield the control as a function of the Lagrange multiplier from the third of (4). This makes the Euler–Lagrange equations a purely differential TPBVP. In the present implementation of the method, this step is accomplished through a symbolic manipulation of (1)–(2). (If a TPBVP solver implementing a collocation method is used, the approximate solution discretized over the time grid is already put in the proper form).

In next section it is shown through examples that the solution obtained with the MASRE represents a good approximation of the optimal solution. The latter is labelled \( x^{(\text{opt})}(t) \), \( \lambda^{(\text{opt})}(t) \), and \( u^{(\text{opt})}(t) \), and the associated cost index \( J^{(\text{opt})} \). It is not straightforward to assert whether \( J^{(\text{opt})} \) is better or worse than the state-dependent objective function \( J \), as \( x^{(\text{opt})}(t) \) verifies a “different” dynamics than those imposed by the MASRE algorithm.
Problem Statement
\[ \dot{x} = f(x, t) + g(x, t) u \]
\[ J = \phi(x(t_f), t_f) + \int_{t_i}^{t_f} L(x, u, t) dt \]

Parents
\[ A_1(x, t), A_2(x, t), \ldots, A_n(x, t) \]

Factorization
Define \( B(x, t), Q(x, t), R(u, x, t), S(x(t_i), t_i) \)

\[ k = k + 1 \]

Problem 0
Define \( A(\alpha) = A^{(0)}(\alpha), [\text{Eq. (48), (52)}] \)
\[ B = B^{(0)}, Q = Q^{(0)}, R = R^{(0)} [\text{Eq. (9)}] \]

\[ k = 0 \]

Problem 1
Define \( A(\alpha) = A^{(1)}(\alpha), [\text{Eq. (48), (53)}] \)
\[ B = B^{(1)}, Q = Q^{(1)}, R = R^{(1)} [\text{Eq. (12)}] \]

Iter \( k > 0 \)

Solver
Compute:
1) \( \phi_{xx}, \phi_{x\lambda}, \phi_{\lambda x}, \phi_{\lambda\lambda} \) [Eqs. (23)–(24)]
2) \( \lambda \) [Eq. (26) HCP/\(Q(\alpha^{(k)}), R^{(k)} SC\)/\(S(\alpha^{(k)}), SCP\)/\(MCP\)]
3) \( x^{(k)}(t), \lambda^{(k)}(t), u^{(k)}(t) \) [Eqs. (21)–(22), (19)]

Optimization
\[ \alpha^{(k)} = \arg \max \sigma_{\min}(P(\alpha)) \quad [\text{Eq. (51)}] \]

Error
\[ \varepsilon = ||x^{(k)} - x^{(k-1)}||_\infty \]

\[ \varepsilon < \text{tol} \]

Stop
\[ \varepsilon > \text{tol} \]

Approximate Solution
\[ x(t) = x^{(k)}(t), \lambda(t) = \lambda^{(k)}(t), u(t) = u^{(k)}(t) \]

Optimal Solution
\[ x^{(\text{opt})}(t), \lambda^{(\text{opt})}(t), u^{(\text{opt})}(t), J^{(\text{opt})} \]

Figure 5. Flowchart for the implementation of the MASRE scheme.
V. Application cases

In this section the MASRE scheme presented in Section IV is applied to three sample problems. The properties of the solution method are discussed. It is shown that the approximate solutions reproduce well the optimal solution. The latter is found by solving the Euler–Lagrange TPBVP using the approximate solutions as first guess (see Figure 5).

V.A. Problem 1

The problem set up is that of Section II.E.1. Beside $A_1$ and $A_2$ in (43), two additional SDC matrices have been provided,

$$A_3(x) = \begin{bmatrix} x_2 & 1-x_1 \\ 0 & x_1 \end{bmatrix}, \quad A_4(x) = \begin{bmatrix} x_2 & 1-x_1 \\ x_2 & 0 \end{bmatrix}.$$  \hspace{1cm} (54)

It is worth mentioning that the controllability matrices associated to the four factorizations are

$$K_{c_1}(x) = \begin{bmatrix} 0 & 1 \\ 1 & x_1 \end{bmatrix}, \quad K_{c_2}(x) = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, \quad K_{c_3}(x) = \begin{bmatrix} 0 & 1-x_1 \\ 1 & x_1 \end{bmatrix}, \quad K_{c_4}(x) = \begin{bmatrix} 0 & 1-x_1 \\ 1 & 0 \end{bmatrix},$$

and it is clear that $\text{rank}(K_{c_1}(x)) = \text{rank}(K_{c_2}(x)) = 2 \forall x$, while both $K_{c_3}(x)$ and $K_{c_4}(x)$ become singular at $x_1 = 1$. It is interesting to notice that the factorizations that are not always controllable, $A_3$ and $A_4$, are not “natural” representations of the uncontrolled dynamics, as their derivation relies on some sort of addition and subtraction of the same term.

The MASRE algorithm has been initiated with $A_1$, $A_2$ in (43), $A_3$, $A_4$ in (54), and with $B$, $Q$, $R$ as in Section II.E.1. The iterations are summarized in Table 3 where the error, the Gramian minimum singular value, and the objective function are reported for each iteration. The last column in Table 3 reports the parent factorization which the family converges to (through the weighting vector $\alpha$, whose first guess is $\alpha^{(0)} = (1-\sqrt{2}/2, 1-\sqrt{2}/2, 1-\sqrt{2}/2)^T$); the termination tolerance is set to $10^{-6}$. Nine iterations are needed to converge to an approximate solution. The algorithm is able to detect the non-controllability associated to $A_3$, $A_4$ and to discard them. The optimal matrix turns out to be $A_1$, and this, remembering the results obtained in Section II.E.1, confirms that the algorithm is able to automatically detect the best factorization in terms of controllability and thus cost function.

Figures 6(a) and 6(b) show the first, second, and last iterations in terms of solution trajectory $(x_1, x_2)$ and control profile $u(t)$. The algorithm is able to converge to the final approximate solution $(k = 9)$ although the first iteration lies quite far from it. The approximate solution is then refined by enforcing the Euler–Lagrange TPBVP, and the results are
shown in Figures 6(c) and 6(d). From this step it can be inferred that the developed scheme approximates well the final, optimal solution. The cost of the latter is $J^{(\text{opt})} = 2.606 \times 10^1$, which means a 0.2% difference with respect to $J$ in Table 3, $k = 9$. This refinement makes it clear that the approximate solution is not optimal, but rather suboptimal.

![Figure 6](image-url)
V.B. Problem 2

This problem consists in minimizing the same performance index of Problem 1 (Eq. (41); i.e., $B, Q, R$ as in Section II.E.1) subject to

\[
\begin{align*}
\dot{x}_1 &= x_1^2 x_2 + x_2, \\
\dot{x}_2 &= x_1 x_2 + u,
\end{align*}
\]

(55)

with the same boundary conditions and initial, final time. The algorithm has been fed with four different SDC matrices of the uncontrolled dynamics

\[
A_1(x) = \begin{bmatrix} x_1 x_2 & 1 \\ 0 & x_1 \end{bmatrix}, \quad A_2(x) = \begin{bmatrix} 0 & x_1^2 + 1 \\ 0 & x_1 \end{bmatrix}, \quad A_3(x) = \begin{bmatrix} x_1 x_2 & 1 \\ x_2 & 0 \end{bmatrix}, \quad A_4(x) = \begin{bmatrix} 0 & x_1^2 + 1 \\ x_2 & 0 \end{bmatrix},
\]

(56)

all of which are always controllable; i.e., $\det(K_c) \neq 0$, $i = 1, \ldots, 4$. Differently from the previous problem, where two factorizations were evidently not “natural” representations of the uncontrolled dynamics, this problem is expressed naturally with four factorizations, so that one would not know a priori which SDC matrix should be used to approximate the optimal solution with the ASRE algorithm. MASRE is initiated with the same settings of Problem 1.

The summary of iterations is reported in Table 4, where not all the steps are reported for brevity sake (the algorithm converges in 21 iterations). It can be seen that the algorithm automatically converges toward the parent $A_2$ factorization, but uses $A_4$ in the first two iterations. In particular, $\sigma_{\min}(P)$ is on the order of $10^7$ in the first iteration, but then it drops to $10^3$ in the second; here the algorithm is able to detect a better controllability associated to $A_2$, and therefore it adjusts $\alpha$ to underweight $A_4$ in favor of $A_2$. As before, the first, second, and last iterations in terms of solution trajectory and control profile are shown in Figures 7(a) and 7(b). In order to assess the quality of this solution, the problem has been solved with the original algorithm, employing the four different factorizations separately, and it has been confirmed that the the best solution is the one related to $A_2$. Moreover, the original ASRE solver is not able to provide a solution using $A_1$ or $A_3$: the former yields a sequence of iterations whose error stably oscillates without converging to any solution, while the latter stops as the transition matrix becomes ill-conditioned after a number of iterations. This implies that controllability is at most a necessary conditions for a converging factorization, but the above discussion also stresses the fact that optimizing the controllability allows us to dodge those factorizations that would not lead to any solution.

The approximate solution has been refined by solving the TPBVP associated to (55) and (41). The results are shown in Figure 7(c) and 7(d). Again, this step makes it clear that
the approximate solution represents well the final, optimal solution. The objective function associated to the latter is \( J^{(\text{opt})} = 9.413 \), which is 2.9% less than that in Table 4, \( k = 21 \), so showing once again that the approximate solution is not optimal but only sub-optimal.

**Table 4.** MASRE iterations for Problem 2, CPU time 231.9 s.

<table>
<thead>
<tr>
<th>( k )</th>
<th>Error</th>
<th>( \sigma_{\text{min}}(P) )</th>
<th>( J )</th>
<th>Parent</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2.000</td>
<td>( 1.739 \times 10^7 )</td>
<td>14.442</td>
<td>( A_4 )</td>
</tr>
<tr>
<td>2</td>
<td>1.376</td>
<td>( 2.978 \times 10^3 )</td>
<td>7.876</td>
<td>( A_4 )</td>
</tr>
<tr>
<td>3</td>
<td>1.121</td>
<td>( 1.355 \times 10^4 )</td>
<td>10.858</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>4</td>
<td>( 7.010 \times 10^{-1} )</td>
<td>( 1.615 \times 10^4 )</td>
<td>9.327</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>5</td>
<td>( 2.423 \times 10^{-1} )</td>
<td>( 1.313 \times 10^4 )</td>
<td>9.798</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>6</td>
<td>( 9.391 \times 10^{-2} )</td>
<td>( 1.480 \times 10^4 )</td>
<td>9.767</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>7</td>
<td>( 6.236 \times 10^{-2} )</td>
<td>( 1.495 \times 10^4 )</td>
<td>9.643</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>8</td>
<td>( 2.657 \times 10^{-2} )</td>
<td>( 1.434 \times 10^4 )</td>
<td>9.708</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>9</td>
<td>( 9.479 \times 10^{-3} )</td>
<td>( 1.443 \times 10^4 )</td>
<td>9.710</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>10</td>
<td>( 6.504 \times 10^{-3} )</td>
<td>( 1.452 \times 10^4 )</td>
<td>9.696</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>11</td>
<td>( 2.838 \times 10^{-3} )</td>
<td>( 1.449 \times 10^4 )</td>
<td>9.701</td>
<td>( A_2 )</td>
</tr>
<tr>
<td>21</td>
<td>( 9.011 \times 10^{-7} )</td>
<td>( 1.449 \times 10^4 )</td>
<td>9.701</td>
<td>( A_2 )</td>
</tr>
</tbody>
</table>

**V.C. Problem 3**

In this problem the controlled, planar Keplerian motion of a spacecraft in polar coordinates is studied. The dynamics are

\[
\begin{align*}
\dot{x}_1 &= x_3, \\
\dot{x}_2 &= x_4, \\
\dot{x}_3 &= x_1 x_4^2 - 1/x_1^2 + u_1, \\
\dot{x}_4 &= -2x_3 x_4/x_1 + u_2/x_1,
\end{align*}
\]

(57)

where \( x_1 \) is the radial distance from the attractor, \( x_2 \) the phase angle, \( x_3 \) the radial velocity, and \( x_4 \) the transversal velocity; \( u_1, u_2 \) are the radial and transversal components of the control acceleration (Eqs. (57) are written in canonical units). The objective function is

\[
J = \frac{1}{2} \int_{t_i}^{t_f} (u_1^2 + u_2^2) \, dt.
\]

(58)

The initial state is \( x_i = (1, 0, 0, 1) \), and \( t_i = 0, t_f = \pi \). Two different problems are solved, according to the definition of \( x_f \). In the HCP, \( x_f = (1.52, \pi, 0, 1.52^{-3/2}) \), while in the MCP \( x_{2,f} \) is unspecified; i.e., the final phase angle is left free. (This set up mimics an Earth–
Mars transfer with the planets moving in circular, coplanar orbits. Eqs. (57) have been factorized by using

\[
A_1(x) = \begin{bmatrix}
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 \\
-x_4^2 & 0 & 0 & x_1 x_4 \\
0 & 0 & -2x_4 & x_1
\end{bmatrix}, \\
A_2(x) = \begin{bmatrix}
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 \\
x_4^2 - \frac{1}{x_1^3} & 0 & 0 & 0 \\
0 & 0 & -2x_4 & x_1
\end{bmatrix}, \\
A_3(x) = \begin{bmatrix}
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 \\
x_4 x_1 - \frac{1}{x_4 x_1^2} & 0 & 0 & x_1 \\
0 & 0 & -2x_4 & x_1
\end{bmatrix}, \\
B(x) = \begin{bmatrix}
0 \\
0 \\
0 \\
0
\end{bmatrix}.
\]

It can be noticed that all the \( A_i \) \((i = 1, 2, 3)\) as well as \( B \) are not Lipschitz continuous when \( x_1 \to 0 \) (and \( x_4 \to 0 \) in \( A_3 \) only). These conditions are however never expected to occur
along the solutions of the present problem, and therefore the MASRE is run with the input as in (59). Moreover, the origin is not an equilibrium point of the system, and the conditions for writing (57) into a state-dependent form as per [18] are violated. The objective function (58) is written in the form (6) with $Q = 0_{4 \times 4}$, $R = I_{2 \times 2}$, and, $S = 0$ (in the MCP only).

The HCP is solved with 15 iterations, and its most relevant figures are reported in Table 5 (in the first iteration it is not possible to ascertain the closest parent factorization). It is interesting to notice that the factorization associated to $A_1$ is discarded because of its non-controllability, although the associated ASRE is able to converge (when fed with $A_1$) to a cost higher than that in Table 5, $k = 15$. This means that the MASRE algorithm prunes away the less performing SDC matrices, and uses the ones that likely produce the best solution. The approximate solutions for some sample iterations are shown in Figures 8(a), 8(c), and 8(e). When refined, the approximate solution leads to the optimal solution in Figures 8(b), 8(d), 8(f). The optimal cost is $J^{(opt)} = 1.830 \times 10^{-1}$, which is 18% less than $J$ in Table 5, $k = 15$. This is mostly due to the difference in the $u_2$ profiles (see Figure 8(f)).

<table>
<thead>
<tr>
<th>$k$</th>
<th>Error</th>
<th>$\sigma_{\min}(P)$</th>
<th>$J$</th>
<th>Parent</th>
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<td>$2.279 \times 10^{-1}$</td>
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</tr>
<tr>
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<td>$3.047 \times 10^{-1}$</td>
<td>$2.254 \times 10^{-1}$</td>
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</tr>
<tr>
<td>12</td>
<td>$1.074 \times 10^{-5}$</td>
<td>$3.047 \times 10^{-1}$</td>
<td>$2.254 \times 10^{-1}$</td>
<td>$A_2$</td>
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<tr>
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<td>$3.047 \times 10^{-1}$</td>
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<td>$2.254 \times 10^{-1}$</td>
<td>$A_2$</td>
</tr>
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</table>

The sequence of iterations for the MCP is reported in Table 6. In this case, the algorithm converges in 13 iterations and uses again $A_2$ as parent factorization. The cost of the refined, optimal solution is $J^{(opt)} = 3.694 \times 10^{-2}$, which is 23% less than that in Table 6, $k = 13$. A set of iterated solutions, as well as the final approximated and the optimal solution are shown in Figure 9. As it would have been expected, by letting the algorithm to determine the final angle, the MCP cost function is lower than that of the HCP, and even the convergence is
Figure 8. Problem 3 (HCP): iterated, approximate, and optimal solution.

faster, since the additional degree of freedom represented by the unspecified final phase is exploited. The optimal solution in Figure 9 can be compared with that in [2], where the minimum-time problem is solved for the same set up. Albeit with proper differences, transfer orbit and the thrust profile are similar.
Figure 9. Problem 3 (MCP): iterated, approximate, and optimal solution.
Table 6. MASRE iterations for Problem 3 (MCP), CPU time 19.9 s.

<table>
<thead>
<tr>
<th>Error</th>
<th>$\sigma_{\text{min}}$</th>
<th>J</th>
<th>Matrix</th>
</tr>
</thead>
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<tr>
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VI. Conclusion

Approximate methods for solving nonlinear optimal control problems aim to circumvent the classic Euler–Lagrange equations by using ad hoc, simpler schemes. As this eases the search for a feasible solution, it does not guarantee optimality, as in general these methods lead to sub-optimal solutions. This is the case of the approximating sequence of Riccati equations, which solves the optimal control problem as a sequence of time-varying linear quadratic regulators. This method depends upon the factorization chosen for the uncontrolled dynamics. It has been shown that even for a simple problem, two different factorizations lead to two different solutions, in terms of system trajectory, objective function, and algorithm iterations. This work makes it possible to automatically select the best factorization, from a set of candidate ones, during the solution process. This is done by maximizing, at each iteration, the controllability Gramian associated to a family of factorizations. The augmented effort consists in solving a nonlinear programming problem at each iteration. The developed method can be used also in other state dependent control schemes. Examples show that the method is effective, and approximates well the true optimal solution, though it would be useful to extend it to accommodate further features (e.g., control saturation, variable final time, non-affinity of control, etc.).
References


Optimal Selection of the Coefficient Matrix in State-Dependent Control Methods – F.T., M.M., and F.B.-Z.


