# Scenario-based Economic Dispatch with Uncertain Demand Response

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Abstract—This paper introduces a new computational framework to account for uncertainties in day-ahead electricity market clearing process in the presence of demand response providers. A central challenge when dealing with many demand response providers is the uncertainty of its realization. In this paper, a new economic dispatch framework that is based on the recent theoretical development of the scenario approach is introduced. By removing samples from a finite uncertainty set, this approach improves dispatch performance while guaranteeing a quantifiable risk level with respect to the probability of violating the constraints. The theoretical bound on the level of risk is shown to be a function of the number of scenarios removed. This is appealing to the system operator for the following reasons: (1) the improvement of performance comes at the cost of a quantifiable level of violation probability in the constraints; (2) the violation upper bound does not depend on the probability distribution assumption of the uncertainty in demand response. Numerical simulations on (1) 3-bus and (2) IEEE 14-bus system (3) IEEE 118-bus system suggest that this approach could be a promising alternative in future electricity markets with multiple demand response providers.

Index Terms—Demand response provider, scenario approach, stochastic economic dispatch

Nomenclature				
ISO	independent system opeartor			
LSE	load serving entity			
DRP	demand response provider			
$P_{G,i}$	power generation of a generator			
$\overline{P_{DR,j}}$	maximum DR commitment of a DRP			
$P_{DR,j}$	accepted DR commitment of a DRP			
$P_{base}$	baseline for end-consumers			
$\pi_{s,j}$	DRP's incentive price to end-consumers			
$\pi_{DR,j}$	a DRP's offer price			
$\delta_j$	DR ratio for DRP			
$P_{DR,j}^{real}$	realized DR commitment			
ξ	decision variable vector			
$P_{Lk}$	total load on bus $k$			
$\alpha$	friction factor for re-dispatch cost			

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- $\epsilon$  constraint violation probability
- $\beta$  confidence level of  $\epsilon$
- $U^s, U^r$  uncertainty set for  $\delta$  in stochastic/robust model  $N_k$  number of samples in the uncertainty set
  - p number of samples removed
  - d number of decision variables

#### I. INTRODUCTION

This paper explores a promising alternative for managing uncertainties in power system optimization. The context of this paper lies in the day-ahead scheduling of resources with many demand response providers. There has been an increasing need for demand response (DR) since the integration of renewables such as solar and wind requires more flexibility from the demand side. As an intermediate between the independent system operator (ISO) and end-consumers, demand response providers (DRPs) are permitted to collect demand response commitments from end-consumers through various means of incentives, and are treated as "virtual generators" in most wholesale market clearing processes [1]–[3].

However, unlike conventional generators whose outputs can be fully controlled, "virtual generation" (demand reduction commitments) from DRPs is much less controllable [4], [5]. This uncertainty primarily comes from the fact of bottom-up aggregation from individual end-consumers, whose behavior is highly uncertain.

There has been a large body of literature investigating different optimization tools to manage the uncertainties in the economic dispatch problem. Stochastic approaches focus on maximizing expected social welfare [6]–[8]. However, this would require prior knowledge of the probability distribution of the underlying uncertainty. More recently, a number of papers investigate the robust optimization-based approach to unit commitment and economic dispatch [9], [10]. It provides a worst case optimal dispatch given an uncertainty set. However, the conservativeness of the robust optimization and the computational burden require substantial efforts to overcome.

As an algorithmic implementation of chance-constrained optimization, the scenario approach [11], [12] is introduced in this paper in the day-ahead market (DAM) clearing process. The key advantage of using this approach is that it allows the trade-off between robustness (constraint violation probability) and performance (objective function value) [13].

The contributions of this paper are the following:

(1) It is the first paper that applies the scenario approach to solving the power system economic dispatch problem

in the presence of uncertain demand response providers. Compared with traditional methods such as stochastic and robust optimization which only provide one solution, by sampling from the uncertainty set and gradually removing those scenarios, the scenario approach provides a curve of realization cost versus samples removed or constraint violation probability (Fig. 6(a)(b)).

- (2) Unlike traditional methods, the scenario approach can be used in the economic dispatch without any assumption on the distribution of the underlying uncertainty. Therefore, the scenario approach could potentially help avoid social welfare loss due to the lack of knowledge of the true underlying uncertainty distribution.
- (3) We investigate the impact of several different scenario removal algorithms on the violation of the categories of constraints (such as power adequacy, branch flow etc.). It sheds lights on system operators' choice of which subset of scenarios to be removed based on their priorities.

This paper is organized as follows. Section II formulates a model to describe the uncertainty of DRPs in day-ahead market. Multiple economic dispatch methods including the scenario approach are introduced in Section III. Some numerical examples are given in Section IV to show the impact of DR uncertainty on day-ahead market, how scenario approach works, as well as ISO's preference among DRPs when congestion exists in the system. We conclude with our findings in Section V.

#### II. PROBLEM FORMULATION

In this section the clearing process of day-ahead market in the presence of DRPs is illustrated. We briefly describe the strategy by which a DRP makes its offer, including price and "capacity". The impact of a DRPs' uncertainty on the day-ahead scheduling is formulated, and an estimate of such uncertainty is made using empirical data obtained from the Electric Reliability Council of Texas (ERCOT).

# A. DRP as a Supplier in Day-ahead Market

In many organized regional markets, demand response (DR) is allowed to contribute to energy balance, capacity resource, and ancillary services [1]–[3], [14], [15]. There are two types of DR: price-based and incentive-based. The main difference lies in the manner of encouraging end-consumers for their demand reduction: the former uses direct price signals (such as real-time or peak-valley prices), while the latter uses indirect or hybrid incentives [16].

In this paper we concentrate on an ISO's day-ahead scheduling of energy resources including *incentive-based DRPs*. DRPs aggregate DR commitments from end-consumers, and are treated as "virtual generators" on the supply side in the day-ahead market. Compared to generators owned by generation companies (GenCos), a DRP's offer price is treated as its marginal cost, while the maximum DR commitment is treated as its "capacity".

The clearing process of the day-ahead market with DRPs is shown in Table I, and a three-layer financial structure is visualized in Fig. 1.

# TABLE I CLEARING PROCESS OF DAY-AHEAD MARKET WITH DRPS

- (1) DRPs incentivize their end-consumers to reduce their electricity consumption for certain periods in the next day, and collect their DR commitments:
- (2) DRPs decide their offer sets  $(\pi_{DR,j}, \overline{P_{DR,j}})$ , GenCos prepare cost curves for generators, and LSEs predict their local demand;
- (3) The ISO clears the day-ahead market. Locational marginal prices (LMPs), accepted generations and DR commitments  $P_{DR,j}$  are determined;
- (4) DRPs receive payments according to the *realized* DR amount  $P_{DR,j}^{real}$  multiplied by the LMP; GenCos are paid according to accepted generation multiplied by the LMP
- (5) End-consumers pay electricity bills to the LSE, and receive rewards from the DRP according to their power reduction in real-time.

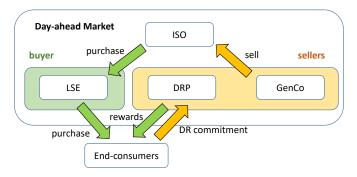


Fig. 1. Three-Layer Financial Structure of DAM with DRP

# B. Offer Set for a DRP

Unlike conventional generators, a DRP's offer set  $(\pi_{DR,j}, \overline{P_{DR,j}})$  cannot be fully determined by physical principles. A model is developed in this subsection to describe a DRP's decisions in steps (1), (2) in Table I, as well as the inner relationship between the offer price  $\pi_{DR,j}$  and "capacity"  $\overline{P_{DR,j}}$  in a DRP's offer set.

Fig. 2(a) shows the aggregated inherent demand curve for DRP j's end-consumers. Load baseline  $P_{base,j}$  is defined as the total electricity consumption of the end-consumers without any DR events; theoretically it equals the quantity on the demand curve corresponding to the retail price  $\pi_{RR}$ . It has been shown in [17] that, by claiming a certain reward  $\pi_{s,j}$  per unit of power reduction, total electricity consumption shrinks as if the retail price increases by the same amount  $\pi_{s,j}$ . Therefore, the maximum DR commitment  $\overline{P_{DR,j}}$  a DRP can collect equals the shrinkage of the load.

If a linear aggregated demand curve is assumed,  $\overline{P_{DR,j}}$  can be calculated as

$$\overline{P_{DR,j}} = min(P_{base,j}, \frac{\pi_{s,j}}{\pi_j^{max} - \pi_{RR}} P_{base,j}), \qquad (1)$$

where  $\pi_i^{max}$  is the Y-axis intercept of the demand curve.

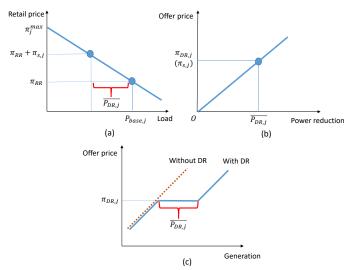


Fig. 2. DRP's Offering Strategy: (a) Inherent demand curve for end-consumers (b) DRP's offer curve (c) Market supply curve with the DRP

Fig. 2(a) is redrawn in 2(b) to illustrate the relationship between the incentive price  $\pi_{s,j}$  and the maximum DR commitment  $\overline{P_{DR,j}}$ . In practice, a DRP is able to choose any  $\pi_{s,j}$  and collect the corresponding  $\overline{P_{DR,j}}$  on the curve. Currently GenCos are required to submit the price-generation curve to the ISO [18], [19]. By assuming that DRPs also need to reveal their cost to the ISO, and that the incentive price is the only variable cost of DRPs, the offer price  $\pi_{DR,j}$  should be equal to the incentive price  $\pi_{s,j}$ . By replacing  $\pi_{s,j}$  with  $\pi_{DR,j}$ , Fig. 2(b) represents the *offer curve* of the DRP, which represents the relationship between its offer price and maximum DR commitment.

A DRP determines its offer by choosing an appropriate incentive price  $\pi_{s,j}$  to its end-consumers. For simplicity, the strategy to choose the optimal  $\pi_{s,j}$  is not discussed in this paper. The maximum DR commitment  $\overline{P_{DR,j}}$  is collected based on the inherent demand curve, and the DRP's offer price  $\pi_{DR,j}$  is set equal to the incentive price. Unlike conventional generators which reveal their full cost curves to the ISO, only one point  $(\pi_{DR,j}, \overline{P_{DR,j}})$  from the DRP's offer curve (2(b)) is submitted to the ISO during each day-ahead scheduling.

Fig. 2(c) shows the change of the total supply curve after a DRP's offer is added into the day-ahead market. The total supply curve at the price  $\pi_{DR,j}$  is shifted horizontally rightward by  $\overline{P_{DR,j}}$ . By choosing different points in its offer curve Fig. 2(b) to formulate the offer set, a DRP can influence the shape of the supply curve in Fig. 2(c), and therefore obtains different rewards in return after the market is cleared.

#### C. Uncertainty of DR

In practice, a DRP's total demand reduction from the baseline usually deviates from its scheduled DR commitment  $P_{DR,j}$  decided by the ISO in step (4), Table I. Common factors that lead to the uncertainty of DR include consumers' uncertain behaviors, ambient condition change, load dynamics,

etc. [4], [5]. The power imbalance or any constraint violation caused by such phenomena in real-time may result in additional re-dispatch cost to the system by means of ramping up expensive generators and the abandonment of renewable resources, etc. Therefore, it is crucial to find a suitable model to describe the uncertainty of DR.

To generalize the behavior of a group of consumers for a certain DRP j, the demand response (DR) ratio is defined as the ratio between the DRP's realized and scheduled DR amounts

$$\delta_j = \frac{P_{DR,j}^{real}}{P_{DR,j}}, j = 1, 2, ..., N.$$
 (2)

In this paper, we make similar assumptions as [4] that the realized DR amount  $P_{DR,j}^{real}$  follows a truncated normal distribution. Since  $P_{DR,j}$  is a constant value submitted to the ISO,  $\delta_j$  should also be a truncated normal distribution  $\delta_j \sim N(\mu_j, \sigma_j^2, \delta_j^{min}, \delta_j^{max})$  (Fig. 3).

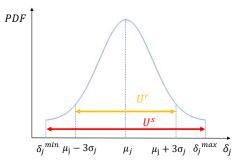


Fig. 3. Truncated Normal Distribution for Demand Response Ratio

Although no previous research has ever provided experimental data for incentive-based DRP, it has been shown in [17] that price-based and incentive-based DR are essentially equivalent; Therefore, price-based DR data of a commercial consumer in ERCOT area [20] is taken as a substitution. By assuming a virtual DRP having only one such consumer (j=1), the uncertainty for this particular user can be calculated to be  $\hat{\mu}_1 = 1$  and  $\hat{\sigma}_1 = 0.67$ .

The above assumptions may lead to an over-estimation of the uncertainty in incentive-based DRPs. First, consumer(s) using price-based DR are liable to become numb to too frequent changes of real-time prices, while incentive-based DR usually has less DR events and consumers are more focused on each DR event. In addition, the end-consumer in [20] is not assigned to any DRP; it means that the consumer cannot receive timely notifications from the DRP about the upcoming DR events. However, this estimation of  $\delta_j$  is still valuable, and will be used in Section IV on a treatment group to show the impact of uncertainty levels of DR on market clearing results.

# III. ECONOMIC DISPATCH METHODS IN DAY-AHEAD MARKET

In this section multiple economic dispatch models are introduced to describe the clearing process of the dayahead market. The ISO's goal is to maximize social welfare (or minimize total supply cost). Its decision variable  $\xi = [P_G, P_{DR}]$  consists of two components: accepted generation  $P_G = [P_{G,i}], i = 1, 2, ..., M$ , and accepted DR commitment  $P_{DR} = [P_{DR,j}], j = 1, 2, ..., N$ .

Recalling the definition of *demand response ratio* in (2), let the vector  $\delta = [\delta_1, \delta_2, ..., \delta_N]$  describe the uncertainty of all DRPs in the power system. As mentioned above,  $\delta_j$  follows a truncated normal distribution  $\delta_j \sim N(\mu_j, \sigma_j^2, \delta_j^{min}, \delta_j^{max})$ . If not specified differently, DR ratios for different DRPs are independent of each other.

#### A. Deterministic Model

The deterministic model widely used in today's day-ahead market neglects all the uncertainties in the power system. In our problem, this is equivalent to suppose that all realized DR commitments are equal to their scheduled amount ( $\delta_j=1,j=1,2,...,N$ ). Therefore the deterministic optimization problem is described as follows:

$$\min_{\xi} l(\xi)$$
s.t.  $f_1(\xi) = 0$ 

$$f_2(\xi) \le 0$$

$$\xi \le \xi \le \overline{\xi}.$$
(3)

The above  $l(\xi)$  represents the *total cost on supply side* including generators and DRPs. As a "virtual generator", the cost of a DRP is assumed to be the multiple of its offer price  $\pi_{DR,j}$  and the accepted DR commitment  $P_{DR,j}$ , namely

$$l(\xi) := \sum_{i=1}^{M} (a_i P_{G,i}^2 + b_i P_{G,i}) + \sum_{j=1}^{N} P_{DR,j} \pi_{DR,j}.$$
 (4)

Equality constraint  $f_1(\xi) = 0$  instead represents the power balance requirement of supply and demand. To this purpose,

$$f_1(\xi) = -\left(\sum_{i=1}^M P_{Gi} + \sum_{j=1}^N P_{DR,j}\right) + P_L.$$
 (5)

Eventually,  $f_2(\xi) \leq 0$  is meant to represent branch flow constraints, by defining

$$f_2(\xi) = H \cdot P_{inj}(\xi) - F^{max},\tag{6}$$

where  $P_L$  represents the total system load, H is the distribution factor matrix, and  $P_{inj}(\xi)$  is a vector indicating net branch flow injections on each bus.

In addition, decision variables  $\xi = [P_G \ P_{DR}]$  should satisfy their capacity limits  $\underline{\xi} \leq \xi \leq \overline{\xi}$ , which are equivalent to

$$\frac{P_{G,i} \le P_{G,i} \le \overline{P_{G,i}}, i = 1, 2, ..., M}{0 \le P_{DR,j} \le \overline{P_{DR,j}}, j = 1, 2, ..., N,}$$
(7)

#### B. Stochastic Model

The stochastic model focuses on minimizing the *expectation* of supply-side cost in the power system considering the uncertainties involved. Particularly, this paper concentrates on the uncertainty of realized DR commitment  $P_{DR,j}^{real}$  caused by random behaviors of end-consumers (measured by DR ratios). Therefore, uncertainty set  $U^s$  in the stochastic model is defined, as the collection of  $N_k$  independent and identically distributed (i.i.d.) extractions of  $\delta$ , where  $\delta$  is the vector whose elements are the DR ratios from the DRPs  $(\delta_j)$ . In practice,  $U^s$  is created by directly using historical data. It can also be formulated by Monte Carlo sampling from the known probability distribution of  $\delta$ .

The stochastic optimization problem is described as follows:

$$\min_{\xi} E[l(\xi, \delta)]$$

$$s.t. \quad \tilde{f}_1(\xi) \le 0$$

$$f_2(\xi, \delta) \le 0, \forall \delta \in U^s$$

$$\underline{\xi} \le \xi \le \overline{\xi}.$$
(8)

Compared to (4), the cost function  $l(\xi, \delta)$  is influenced by the realized DR ratio. Therefore, the objective function in stochastic model is chosen as the average cost

$$E[l(\xi, \delta)] = E[\sum_{i=1}^{M} (a_i P_{Gi}^2 + b_i P_{Gi}) + \sum_{j=1}^{N} \underbrace{\delta_j P_{DR,j}}_{P_{DR,j}} \pi_{DR,j}]. \quad (9)$$

In addition, the uncertainty of the DRPs may also cause energy inadequacy to the system. Chance constraints are used to describe the requirement of maintaining energy adequacy with a probability higher than a certain level  $(1-\gamma)$ 

$$Pr(\sum_{i=1}^{M} P_{Gi} + \sum_{j=1}^{N} \delta_{j} P_{DR,j} > P_{L}) \ge 1 - \gamma.$$
 (10)

In order to solve the chance constrained optimization problem at low computational cost, in (8) the constraint (10) is updated by the following linear approximation, obtained from the semi-analytical method proposed in [21]:

$$-(\sum_{i=1}^{M} P_{Gi} + \sum_{j=1}^{N} g_j(\delta_j, \gamma) P_{DR,j}) + P_L \le 0$$

$$\underbrace{\tilde{f}_1(\varepsilon)}$$

$$(11)$$

where  $g_j := \mu_j + \sigma_j \Phi^{-1}(1 - \gamma)$  is always a constant value if the distribution of  $\delta_j$  is known and  $\gamma$  is pre-determined.

In addition, branch flow constraints are required to be strictly satisfied for all the scenarios of  $\delta \in U^s$ , leading to

$$\underbrace{H \cdot P_{inj}(\xi, \delta) - F^{max}}_{f_2(\xi, \delta)} \le 0, \forall \delta \in U^s. \tag{12}$$

#### C. Robust Model

The following robust model is preferred when the ISO pays more attention to minimizing the supply-side cost for the worst case of the uncertain parameter  $\delta$  (13). In the robust model, energy adequacy and branch flow constraints need to be satisfied for all the scenarios in an uncertainty set  $U^r$ . The problem is described as

$$\min_{\xi} \max_{\delta \in U^r} l(\xi, \delta)$$
s.t.  $f_1(\xi, \delta) \le 0, \forall \delta \in U^r$ 

$$f_2(\xi, \delta) \le 0, \forall \delta \in U^r$$

$$\xi \le \xi \le \overline{\xi},$$
(13)

where the quadratic cost function  $l(\xi, \delta)$  and branch flow constraint function  $f_2(\xi, \delta)$  have the same form as in (9) and (12). Energy adequacy constraint is described as

$$f_1(\xi, \delta) = -\left(\sum_{i=1}^{N} P_{Gi} + \sum_{j=1}^{M} \delta_j P_{DR, j}\right) + P_L \le 0.$$
 (14)

The uncertainty set  $U^r$  is determined by the  $3-\sigma$  rule, that is,  $U^r$  is a box whose sides are intervals  $(\delta_j^{min}, \delta_j^{max})$  with  $\delta_j^{min} = \mu_j - 3\sigma_j$  and  $\delta_j^{max} = \mu_j + 3\sigma_j$ .

## D. Scenario Approach Model

The scenario approach introduced by Campi and Garatti [11], [12] is a refinement of the convex robust model. By selecting a finite number of scenarios from the uncertainty set  $U^s$ , this approach improves dispatch performance while guaranteeing a quantifiable risk level. Similar to the robust model (13), the scenario approach in economic dispatch is described as follows:

$$\min_{\xi} \max_{\delta \in U'} l(\xi, \delta)$$

$$s.t. \quad f_1(\xi, \delta) \le 0, \forall \delta \in U'$$

$$f_2(\xi, \delta) \le 0, \forall \delta \in U'$$

$$\xi < \xi < \overline{\xi}.$$
(15)

where  $l(\xi, \delta)$ ,  $f_1(\xi, \delta)$  and  $f_2(\xi, \delta)$  are as in (13).

The only difference lies in the uncertainty set U'. References [11], [12], [22], [23] have proposed multiple approaches in formulating U' in scenario approach; In this paper, the number of scenarios  $(N_k)$  in the original uncertainty set  $U^s$  is assumed to be finite, which is reasonable especially when  $U^s$  is obtained from historical data. U' is formulated by removing a certain number of scenarios p from  $U^s$ . It can be shown that formulation (16) is equivalent to (15) (see e.g. references [11], [24], [25]).

$$\min_{\xi,h} h$$

$$s.t. \quad l(\xi,\delta) \le h, \forall \delta \in U'$$

$$f_1(\xi,\delta) \le 0, \forall \delta \in U'$$

$$f_2(\xi,\delta) \le 0, \forall \delta \in U'$$

$$\xi \le \xi \le \overline{\xi}$$
(16)

The reason for transforming (15) to (16) is to reformulate it so that it can be easily handled by standard optimization solvers. Moreover, this new formulation allows for easy removal of scenarios by simply removing the corresponding constraints.

As a brief explanation to (16), note that formulation (15) is equivalent to

$$\min_{\xi,h} h$$

$$s.t. \quad \max_{\delta \in U'} l(\xi, \delta) \le h$$

$$f_1(\xi, \delta) \le 0, \forall \delta \in U'$$

$$f_2(\xi, \delta) \le 0, \forall \delta \in U'$$

$$\xi \le \xi \le \overline{\xi},$$

since the choice of  $\xi$  that allows one to reduce h as much as possible is that returning the minimum of  $\max_{\delta \in U'} l(\xi, \delta)$ ; so  $\max_{\delta \in U'} l(\xi, \delta) \leq h$  is clearly equivalent to  $l(\xi, \delta) \leq h, \forall \delta \in U'$ .

It is perhaps worth noticing that in (16), by removing p scenarios from the uncertainty set, a total number of 3p constraints are removed since there are three constraints  $l(\xi,\delta) \leq h, \ f_1(\xi,\delta) \leq 0$  and  $f_2(\xi,\delta) \leq 0$  corresponding to each scenario  $\delta$ .

The main result associated to (15) is as follows. Let  $\epsilon \in (0,1)$  be a "violation" parameter and  $\beta \in (0,1)$  be a confidence parameter. Then, Theorem 2.1 of [12] claims that if the following inequality is satisfied

$$\binom{p+d-1}{p} \sum_{i=0}^{p+d-1} \binom{N_k}{i} \epsilon^i (1-\epsilon)^{N_k-i} \le \beta \tag{17}$$

(being d the dimensionality of  $\xi$ ), then with probability no more than  $1 - \beta$ , the optimal solution of (15), say  $\xi^*$  satisfies all the constraints in (13) but at most an  $\epsilon$  fraction of violation  $(f(\xi^*, \delta) > 0)$  with probability at most  $\epsilon$ ).

The highlight of this theorem is its universality. The result holds as long as the objective function is convex, regardless of the interpretation of the objective function, constraints or probability distribution of uncertain variables. In addition, a careful choice of removing scenarios can potentially lead to a higher performance with a quantifiable risk level in the sense of probability of violation of constraints. In practice, the finite uncertainty set  $U^s$  can be formulated using historical data.

Removing scenarios from the uncertainty set may help to improve the performance of the objective function. At the same time, it increases the system risk level, which is mainly reflected in three aspects:

- (i) Power inadequacy: Violation of  $f_1(\xi^*, \delta) \leq 0$  indicates energy inadequacy in certain scenarios  $\delta \in U^s$ .
- (ii) Branch flow violation: Violation of  $f_2(\xi^*, \delta) \leq 0$  means that the optimal solution  $\xi^*$  causes overflow in certain scenarios  $\delta \in U^s$ .
- (iii) h-violation: Violation of  $l(\xi^*, \delta) \leq h^*$  means that higher supply cost (lower social welfare) is possible when applying the "optimal" solution  $\xi^*$  in certain scenarios in  $U^s$ ;  $h^*$  is an over-optimistic estimation.

The pleasant message converged by Theorem 2.1 of [12] is that the risk is quantifiable since one is  $1 - \beta$  confident that any violation happens with probability less than  $\epsilon$ , where  $\epsilon$  is computed from (17).

An important feature of Theorem 2.1 of [12] is that it holds true irrespective of the scenario removal algorithm, i.e. it holds true irrespective on the strategy through which scenario to remove are chosen. As an example, reference [12] has defined the "optimal" scenario removal algorithm as the algorithm that "among all possible eliminations of p scenarios out of total N" implement the one that gives the best improvement on the objective function. However, reference [12], [26] also explain that such an "optimal" algorithm can be computationally intractable. Then, one can resort to suboptimal algorithms to remove scenarios, such as the greedy algorithm, which removes scenarios sequentially; sometimes even the greedy algorithm is computationally unaffordable [12], [26], and other algorithms must be used.

In this paper, the following two algorithms for removing scenarios in the economic dispatch with uncertain DRPs are proposed:

- (1) "Min" algorithm. This algorithm is based on the fact that under-reaction of DR would cause energy inadequacy of the power system in real-time. Usually more expensive generators need to be ramped up by the system operator, which leads to a higher realization cost. Therefore, this algorithm assumes that scenarios with low total DR  $min(\sum_{j=1}^{N} \delta_{j}^{real} P_{DR,j})$  are removed with high priority.
- (2) "Center" algorithm. This algorithm is based on the fact that DRPs' deviations from their scheduled amounts (expectation) may lead to branch flows violations. In scenarios when branch flows are violated, expensive auxiliary generators are used by the ISO to satisfy branch flow constraints. Therefore, this algorithm assumes scenarios with largest DR deviations  $max(\sum_{j=1}^{N}|\delta_{j}^{real}-\mu_{j}|P_{DR,j})$  need to be removed with priority.

These two suboptimal algorithms are tractable vis-a-vis computational cost since they require neither sequential solution nor extra optimization.

Since these two algorithms focus on different aspects of performance improvement, they will have impacts on different types of constraint violations (Type (i),(ii),(iii)). Further discussions on this topic are given in Section IV-B.

The abbreviation of scenario selection "percentage/algorithm" is used in the remainder of this paper. For instance, 20%/min represents the action of removing the 20% smallest  $\delta$ 's following "min" algorithm from the uncertainty set.

# E. Realization Cost

Uncertainty in the realization of DR may lead to multiple constraint violations in real-time, as we have discussed in Section III-D. Therefore, in real-time dispatch (namely, the balancing market), ISO may need to ramp up/down some expensive reserves or auxiliary generators in order to satisfy the power balance as well as branch flow limits.

In our modeling of real-time balancing dispatch, auxiliary generators that have constant marginal cost are assumed to be on the buses that have DRPs. Therefore, each auxiliary generator is responsible for making up the deviation of the realized DR from its day-ahead scheduling level. Real-time balancing cost  $W_{rtd}$  is computed as the auxiliary generators' marginal cost multiplied by the absolute value of ramp up/down power (which is equal to the deviation of DR at the same bus),

$$W_{rtd}(\xi^*, \delta) = \sum_{j}^{N} \pi_{aux,j} |(\delta_j^{real} - \mu_j) P_{DR,j}|, \qquad (18)$$

and the total realization cost is the summation of expected realized dispatch cost and the expectation of balancing cost

$$W_{real} = E_{\delta}[l(\xi^*, \delta) + W_{rtd}(\xi^*, \delta)]. \tag{19}$$

Here  $\pi_{aux,j}$  is the marginal cost for auxiliary generator j, and  $\xi^*$  is the optimal value of  $\xi$  calculated by different models in Section III A-D.

#### IV. NUMERICAL EXAMPLES

Three power systems with DRPs are analyzed in this section in order to compare the performance of the scenario approach with other economic dispatch methods. In addition, the trade-off between performance and system risk level, as well as the impact of distribution of  $\delta$  on social welfare is also discussed.

# A. 3-bus System with one DRP

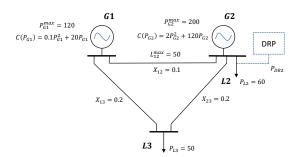


Fig. 4. 3-bus System with the DRP

The 3-bus system shown in Fig. 4 contains two generators, and one DRP on bus 2. All subscripts j are dropped here since there is only one DRP. Recalling Fig. 2(a), a linear aggregated inherent demand curve for end-consumers is assumed, with baseline  $P_{base} = P_{L2} = 60$ , retail price  $\pi_{RR} = 100$ , and Y-axis intercept  $\pi^{max} = 400$ .

Historical data of DR ratio  $U^s$  is created by randomly selecting  $N_k=1000$  points from the truncated normal distribution ( $\mu=1,\sigma=0.1,\delta^{max}=1.5$  and  $\delta^{min}=0.5$ ).  $U^r$  follows the  $3-\sigma$  rule ( $\delta^{max}=0.7,\delta^{min}=1.3$ ). The reliability factor  $\gamma=0.8$  in the stochastic model,  $\beta=10^{-5}$  in the scenario approach, and  $\pi_{aux,j}=150, \forall j$ . The size of the uncertainty set is 1000.

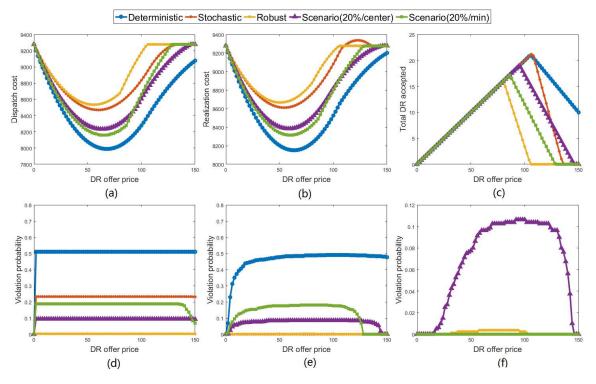


Fig. 5. Dispatch Results Under Different Models (a) Dispatch cost (b) Realization cost (c) DR bids accepted (d) Power inadequacy prob. (Type (i)) (e) Branch flow violation prob. (Type (ii)) (f) h-violation prob. (Type (iii))

#### B. Simulation Results for Economic Dispatch

Simulation results for different economic dispatch models are shown in Fig. 5, and the results at one snapshot are shown in Table II. For all subgraphs in this figure, the x axis represents the different choices of offer price  $\pi_{DR}$  for the DRP. Higher offer price  $\pi_{DR}$  results in more profit per unit DR, but the offer is less likely to be accepted by the ISO when  $\pi_{DR}$  exceeds the marginal cost of other generations.

Economic dispatch performance is measured by the dayahead dispatch cost and the realization cost, while system risk is measured by violation probability. It can be concluded that the robust model is usually most conservative, since it takes the first inflection of DR acceptance in Fig. 5(c), and has the highest dispatch and realization cost in Fig. 5(a)(b). However, this level of conservatism leads to the lowest violation of all constraints (energy adequacy Fig. 5(d), branch flow limits Fig. 5(e) and h-violation status Fig. 5(f)). In contrast, deterministic model has the lowest dispatch and realization cost as well as the highest DR amount accepted. However, its neglect of uncertainty results in huge system risk in power adequacy and branch flow constraints. The stochastic model is relatively neutral in both performance and violation. Unlike the robust model, the deterministic and stochastic models don't measure the h-violation status since they purely concentrate on expected cost rather than cost in the worst case.

Fig. 5 also shows the impact of scenario removal algorithms on different constraint violations (Type (i),(ii) and (iii)). Compared with a more *neutral* "center" algorithm, the "min" algorithm is more *optimistic* in the realization of DR, since it

prefers removing low DR scenarios. The optimistic aspect of "min" algorithm is exhibited in Fig. 5(f) by low *h-violation* probability on the dispatch cost. However, in realization, DR has higher probability to be lower than its expectation, which leads to *lower DR acceptance ratio*, and higher power inadequacy violation (Fig. 5(d)). Given the fact than in our test system DR helps to relieve congestion, lower DR realization also leads to *higher branch flow violation* (Fig. 5(e)).

The system operator may prefer to use the "center" over the "min" algorithm in this situation, since with similar levels of realization cost, the "center" algorithm provides higher DR acceptance ratio as well as lower physical violations (power adequacy and branch flows).

# C. Distributionally-Robust Feature of the Scenario Approach

A significant advantage of the scenario approach is that the constraint violation upper bound is guaranteed by the theory (17). Therefore, optimization by the scenario approach doesn't rely on any assumption of the distribution of the underlying uncertainty. In contrast, when using a stochastic model, such an assumption is necessary. Table II shows how an incorrect assumption on the distribution can affect the dispatch result. We take a snapshot at  $\pi_{DR}=100$ , and the case "Sto(uniform)" is created with the "incorrect" assumption that  $\delta$  is uniformly distributed  $\delta \sim U(0,2)$ . Compared with the true model "Sto" ( $\delta$  follows truncated normal distribution), "Sto(uniform)" clearly over-estimates the uncertainty by using the incorrect distribution; therefore, it schedules less efficient conventional generation in the day-ahead market, which leads to higher

TABLE II  $\mbox{Dispatch Results under } \pi_{DR} = 100$ 

$Model^{a}$	Dispatch Cost	Realization Cost	Total DR	Balance Vio.	Branch Flow Vio.	h-Vio.
Dtm	8278.1	8520.9	20.0	0.496	0.488	N/A
Sto	8704.1	8946.8	20.0	0.215	0	N/A
Rob	9127.9	9181.1	4.4	0.0025	0.0013	0
Sce(min)	8577.5	8792.8	17.7	0.104	0.095	0.105
Sce(center)	8556.2	8727.3	14.1	0.210	0.194	0.013
Sto(uniform)	8718.0	8963.3	20	0.232	0	N/A

<sup>&</sup>lt;sup>a</sup> Acronyms for models: Dtm: Deterministic, Sto: Stochastic, Rob: Robust, Sce: Scenario approach.

dispatch cost, realization cost and violation probability.

# D. Trade-off between Performance and Risk

As mentioned before, the scenario approach enables the trade-off between performance and system risk. Fig. 6 illustrates this trade-off by selecting the "center" algorithm to eliminate scenarios, and taking the snapshot of offer price  $\pi_{DR}=100$ .

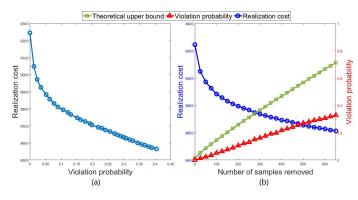


Fig. 6. Trade-off between cost and violation in scenario approach (a) Dispatch cost vs. violation probability (b) Dispatch cost and violation probability vs. number of removed samples

Fig. 6(a) directly shows the inverse proportional relationship between the realization cost and system risk. Fig. 6(b) illustrates that with an increase of the number of scenarios p removed, realization cost plunges immediately but becomes stable when p>600. This indicates a decreasing efficiency in improving the realization cost in scenario approach. On the other hand, the increase of violation probability is close to a linear increase and within the theoretical upper bound (17). The observations in (a)(b) could inspire the decision makers about how to choose an optimal p in the scenario approach.

## E. Influence of $\delta$ Distribution on DR acceptance

As discussed above, the distribution of  $\delta$  has a significant impact on the realization cost as well as the system risk. Simulation results of higher uncertainty in the DRP ( $\hat{\mu}=1,\hat{\sigma}=0.67,\delta^{max}=1.8,\,\delta^{min}=0.2$ , calculated in Section II-C) on 3-bus system, are shown in Fig. 7. Compared to Fig. 5(a)(b), this figure shows a shrinking of DR acceptance and increase of total realization cost. Due to less reliability for the DRP, less DR commitment is taken by the ISO in order to maintain a required system risk level. High violation probability also causes increase in the realization cost.

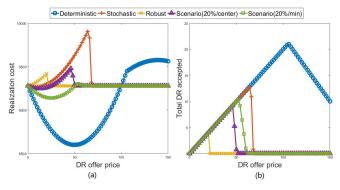


Fig. 7. Dispatch Result Under Different Models ( $\sigma=0.67$ ) (a) Realization cost (b) DR amount accepted

#### F. IEEE 14-bus System with 2 DRPs

In this subsection a more complicated IEEE 14-bus system with two DRPs is analyzed (Fig. 8). We assume that the two largest loads (buses 3 and 4) have DRPs (N=2) that can help to reduce the demand, with  $\pi_{RR}=100,\pi_{c,j}^{max}=400,j=1,2$ . Both DRPs have the same distribution of  $\delta_j$  as that in Section IV-A, but are independent of each other. Line constraint  $L_{24}=30$ , reliability factor  $\gamma=0.8$ , and  $\pi_{aux,j}=150, \forall j$ . The size of the uncertainty set is 1000.

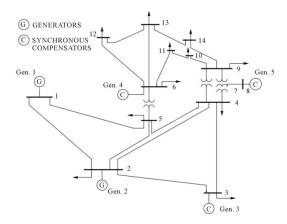


Fig. 8. IEEE 14-bus test system

The "center" algorithm is adopted to select and remove scenarios in our simulation. However, the definition of this algorithm changes with multiple DRPs. Scenarios  $\delta = [\delta_1, \delta_2]$  whose DR commitment has most deviation from the expected

total DR commitment,

$$\max_{\delta} (\sum_{j=1}^{N} |\delta_j - \mu_j| \overline{P_{DR,j}}), \tag{20}$$

would be preferred to be removed. By definition,  $P_{DR,j}$ , j=1,2 should be used in (20) in order to formulate the accepted DR commitment by the ISO. However, since scenarios are removed before the optimization problem is solved by the ISO, and the value of  $P_{DR,j}$  is unknown, we use  $\overline{P_{DR,j}}$ , j=1,2 instead as their rough estimates.

The simulation result in Fig. 9(a)(b) shows that the robust and stochastic models are more conservative than the scenario-approach, with earlier inflection of DR acceptance and higher realization cost. With an increasing number of scenarios removed, the scenario approach (50%) has higher DR acceptance ratio than that of scenario approach (20%).

TABLE III
PROGRAM COMPUTATIONAL TIME

Model	Mean Time(s)	Median Time(s)
Dtm	0.031	0.026
Sto	16.76	13.61
Rob	20.71	17.14
Sce(20%)	14.04	12.98
Sce(50%)	7.63	5.84

The computational cost is very similar among scenario approach (no scenario is removed), stochastic and robust models, since they have a similar number of decision variables as well as number of constraints. By removing scenarios from the uncertainty set, with a decrease in the number of constraints, the scenario approach runs faster than the other two approaches. Table III shows that in the 14-bus system, by removing 50% of the scenarios, the computational time of the scenario approach is only 40% of the stochastic model, and 35% of the robust model.

Fig. 9(c) shows the influence of DRPs' locations on the accepted DR commitments in economic dispatch. Offer price  $\pi_{DR,1}=\pi_{DR,2}$  is assumed in order to make the two DRPs comparable; the only difference lies in the truth that demand reduction on DRP 2 would help to release the congestion in the power system but DRP 1 would not. With the increase of offer price by both DRPs, the acceptance ratio (defined as  $P_{DR,j}/P_{DR,j}$ , j=1,2) of DRP1 drops much earlier than that of DRP2. DRP2's offer is still accepted by the ISO at a price between 56 and 70, while the offer by DRP1 already gets rejected with the same offering price.

In summary, our analysis shows that the location of the DRP is essential to optimal dispatch results when there is line congestion. DRP offers which can help to ease the line congestion are preferred by the ISO, and have higher chance to get accepted.

## G. IEEE 118-bus System with 2 DRPs

In this subsection, a simulation is performed on an IEEE 118 bus system with two DRPs (Fig. 10). We assume Buses 15 (markered with a red circle) and 59 (markered with a

blue circle) have DRPs (N=2), with  $\pi_{RR}=100,\pi_{c,j}^{max}=300, j=1,2$ . There is no line flow limit in the system;  $\gamma=0.8,\,\pi_{aux,j}=150, \forall j,$  and the number of scenairos in the uncertainty set  $U^s,U^r,U^{test}$  are all 1600.

In the simulation, the DRPs' offers are assumed to be  $(\pi_{DR,1}, \overline{P_{DR,1}}) = (30,13.50)$  and  $(\pi_{DR,2}, \overline{P_{DR,2}}) = (35,48.48)$ . Simulation results are shown in Table IV. Some of the key observations include: (a) The robust-based dispatch exhibits the most conservative DR acceptance; (b) The scenario approach incurs less computational time than the stochastic and robust models; and (c), in terms of cost, it's not always better to remove more scenarios; in this case, too optimistic with DR in "Sce(50%)" results in higher dispatch and realization cost.

#### V. CONCLUDING REMARKS

This paper introduces the new computational framework of scenario approach for accounting for the uncertainty of demand response providers (DRPs) in the day-ahead market. Unlike conventional generators, DRPs provide demand response by bottom-up aggregation, which causes high uncertainty of DR because of the behavior of end-consumers. The scenario approach model is introduced in order to obtain better performance than the robust model in most cases, while maintaining an acceptable system risk level. In addition, the scenario approach's key advantage lies in its ability to trade-off between the performance and system risk, compared with existing methods such as deterministic, robust and stochastic.

Simulation results shows that the uncertainty of DR will lead to less amount of DRPs' offer's being accepted in all dispatch methods (stochastic, robust and scenario approach). Though depending on the algorithm employed for selecting and removing scenarios, the scenario approach provides an improvement in expected realization cost over the robust model. In the economic dispatch, the decision maker is free to select and remove a number of scenarios from the uncertainty set in order to lower the realization cost. However, the removal of some scenarios comes at the cost of a small level of violation probability in the constraints. This paper also provides a preliminary assessment of scenario removal algorithms. This may help the system operator to increase the efficiency of improving the dispatch performance in the scenario approach.

Future work will investigate the trade-off of performance and risk involved with the scenario removal algorithms in economic dispatch in order to take care of both optimality and computational cost. Strategic behavior of multiple DRPs in the market will also be investigated in the context of scenario-based economic dispatch framework.

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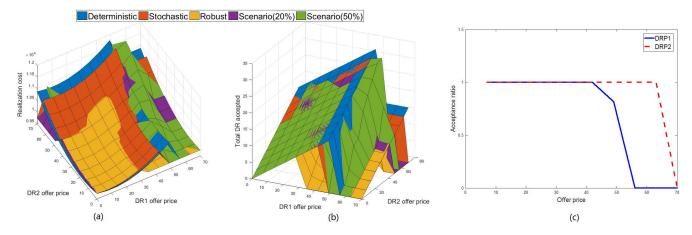


Fig. 9. IEEE 14-bus system with 2 DRPs (a) Realization cost (b) DR offers accepted (c) DR acceptance ratio  $(\pi_{DR,1}=\pi_{DR,2})$ 

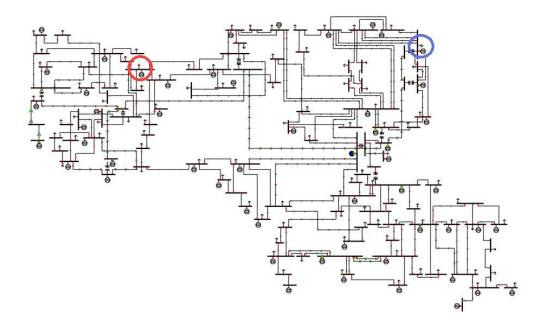


Fig. 10. IEEE 118-bus test system. Red marker: Bus 14, blue marker: Bus 59

TABLE IV
SIMULATION RESULTS ON IEEE 118-BUS SYSTEM

Model	Disp. Cost	Real. Cost	Total Gen	Total DR	Balance Vio.	h-Vio.	Computational Time (s)
Dtm	12562	12812	4180.0	61.97	0.516	-	0.71
Sto	12580	12830	4184.6	61.97	0.167	-	6.58
Rob	12595	12595	4242.0	0	0	0	2.42
Sce(20%)	12595	12595	4242.0	0.34	0.029	0	1.67
Sce(50%)	12812	12814	4188.0	58.16	0.173	0.044	1.52

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