

# An efficient implementation of the algorithm computing the Borel-fixed points of a Hilbert scheme

Paolo Lella  
Dipartimento di Matematica  
Via Carlo Alberto 10  
10123 Torino - Italy  
paolo.lella@unito.it

## ABSTRACT

Borel-fixed ideals play a key role in the study of Hilbert schemes. Indeed each component and each intersection of components of a Hilbert scheme contains at least one Borel-fixed point, i.e. a point corresponding to a subscheme defined by a Borel-fixed ideal. Moreover Borel-fixed ideals have good combinatorial properties, which make them very interesting in an algorithmic perspective. In this paper, we propose an implementation of the algorithm computing all the saturated Borel-fixed ideals with number of variables and Hilbert polynomial assigned, introduced from a theoretical point of view in the paper “Segment ideals and Hilbert schemes of points”, *Discrete Mathematics* 311 (2011).

## Keywords

Borel-fixed ideals, Hilbert scheme, Hilbert polynomial

## 1. INTRODUCTION

The Hilbert scheme  $\mathbf{Hilb}_{p(t)}^n$  parametrizes all the subschemes and all the families of subschemes of the projective space  $\mathbb{P}^n$  with Hilbert polynomial  $p(t)$ . Borel-fixed ideals are a basic tool for the direct study of Hilbert schemes, because

- each component and each intersection of components of a Hilbert scheme contains at least one point defined by a Borel-fixed ideal;
- the Borel-fixed ideals have a strong combinatorial property which makes them very convenient also from an algorithmic perspective.

For instance, Hartshorne in his thesis [9] proved the connectedness of Hilbert scheme constructing sequences of deformations of Borel-fixed ideals (he called them *balanced* ideals) which lead from any point of the Hilbert scheme to the point determined by the unique saturated lexicographic associated to the Hilbert polynomial  $p(t)$ .

More recently, many papers [2, 3, 4, 5, 6, 10, 11] by Bertone, Cioffi, Marinari, Roggero and the author showed

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. To copy otherwise, to republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee.

ISSAC 2012, July 22–25, 2012, Grenoble, France.  
Copyright 2012 ACM 978-1-4503-1269/12/07 ...\$10.00.

how to use Borel-fixed ideals for a local study of the Hilbert scheme, mainly constructing families of ideals sharing with a fixed Borel-fixed ideal the same basis of the quotient space.

Therefore it was very important to have an algorithm computing for each Hilbert scheme  $\mathbf{Hilb}_{p(t)}^n$  all the points defined by Borel-fixed ideals. An algorithm for computing all the Borel-fixed ideals in  $\mathbb{K}[x_0, \dots, x_n]$  with Hilbert polynomial  $p(t)$  based on a combinatorial approach has been proposed in [5] and another algorithm with slight differences to this, and published later, is presented in [12].

In this paper, we describe a concrete implementation of the algorithm, that turns out to be very efficient as we will show with an experimental analysis.

## 2. NOTATION AND GENERAL SETTING

We will consider a field  $\mathbb{K}$  of characteristic 0 and for any polynomial ring  $\mathbb{K}[x] := \mathbb{K}[x_0, \dots, x_n]$  we will order the variables as  $x_n > \dots > x_0$ . Following the notation of [7], we will refer to the Hilbert polynomial of a homogeneous ideal  $I \subset \mathbb{K}[x_0, \dots, x_n]$  as the Hilbert polynomial of the graded module  $\mathbb{K}[x_0, \dots, x_n]/I$ , i.e.

$$p(t) = \dim_{\mathbb{K}} \mathbb{K}[x_0, \dots, x_n]_t / I_t, \quad t \gg 0.$$

A homogeneous ideal  $I$  is said Borel-fixed if it is fixed by the action of the Borel subgroup of upper triangular matrices. Looking at the action of the elements of the Borel subgroup  $\text{Id} + E_{i,j}$ , where  $i < j$  and  $E_{i,j}$  is a matrix with all entries equal to 0 except for the entry of the  $i$ -th row and  $j$ -th column equal to 1, it is possible to prove that the ideal  $I$  has a nice combinatorial property: in fact an ideal  $I$  is Borel-fixed if and only if it is a monomial ideal and

$$x^\alpha \in I \implies \frac{x_j}{x_i} x^\alpha \in I, \quad \forall x_i \mid x^\alpha, \quad x_j > x_i.$$

As done in [5, 10], we define the elementary Borel moves:

- $e_j^-$  as the element  $\frac{x_{j-1}}{x_j}$  in the field of fraction  $\mathbb{K}(x)$  of  $\mathbb{K}[x]$ ;
- $e_i^+$  as the element  $\frac{x_{i+1}}{x_i} \in \mathbb{K}(x)$ ;

and for any monomial  $x^\alpha$  we will say that  $e_i^-$  (resp.  $e_j^+$ ) is admissible on  $x^\alpha$  if  $e_i^-(x^\alpha) = \frac{x_{i-1}}{x_i} x^\alpha \in \mathbb{K}[x]$  (resp.  $e_j^+(x^\alpha) = \frac{x_{j+1}}{x_j} x^\alpha \in \mathbb{K}[x]$ ). We will use additive notation to denote the composition of an elementary move with itself, for instance  $2e_j^- = e_j^- \circ e_j^- = \left(\frac{x_{j-1}}{x_j}\right)^2$ .

$\mathcal{P}(n, r)$  will denote the poset of all monomials in  $\mathbb{K}[x]_r$  with the Borel partial order  $\leq_B$  given by the transitive closure of the order relations

$$\mathbf{e}_i^+(x^\alpha) >_B x^\alpha >_B \mathbf{e}_j^-(x^\alpha)$$

$\forall x^\alpha \in \mathbb{K}[x]_r, \forall \mathbf{e}_i^+, \mathbf{e}_j^-$  admissible, and we call Borel set any subset of a poset  $\mathcal{P}(n, r)$  closed under increasing elementary moves. By definition, given a Borel-fixed ideal  $I \subset \mathbb{K}[x]$ , the monomial basis of each homogeneous piece  $I_t$  of fixed degree will define a Borel set, thus we will write  $\{I_t\}$  meaning the Borel set of  $\mathcal{P}(n, t)$  defined by the monomials in  $I_t$ . On the other hand, given a Borel set  $\mathcal{B} \subset \mathcal{P}(n, r)$ , we will denote by  $\langle \mathcal{B} \rangle$  the Borel-fixed ideal generated by the monomial of  $\mathcal{B}$  and by  $\langle \mathcal{B} \rangle^{\text{sat}}$  its saturation.

For any monomial  $x^\alpha$  we denote

- $\max x^\alpha = \max\{x_i \text{ s.t. } x_i \mid x^\alpha\}$ ;
- $\min x^\alpha = \min\{x_j \text{ s.t. } x_j \mid x^\alpha\}$ .

For the order chosen on the variables, it will not be misleading to write only the index of the variable instead of the variable itself.

For any Borel set  $\mathcal{B} \subset \mathcal{P}(n, r)$  we will denote by  $\mathcal{B}^c$  the complement set  $\mathcal{P}(n, r) \setminus \mathcal{B}$ . Obviously  $\mathcal{B}^c$  is closed under decreasing elementary moves and we will call such a set *order set*, being its dehomogenization (imposing  $x_0 = 1$ ) an order ideal.

For any subset  $S$  of  $\mathcal{P}(n, r)$ ,  $S_{(\geq i)}$  will denote the subset of  $S$  of the monomials with minimum greater than or equal to  $i$ :

$$S_{(\geq i)} := \{x^\alpha \in S \mid \min x^\alpha \geq i\}.$$

Obviously  $S_{(\geq 0)} = S$ .

As well known, given any Borel-fixed ideal  $I$ , its saturation  $I^{\text{sat}}$  is generated by the monomials obtained substituting  $x_0 = 1$  in the monomials that generate  $I$ . Moreover,  $x_0$  is not a nilpotent element for any Borel-fixed ideals.

**PROPOSITION 2.1.** *Let  $I \subset \mathbb{K}[x]$  be a Borel-fixed ideal. The linear form  $x_0$  is regular for  $I$ . Thus, for any shift of degree  $t$ , there is the short exact sequence induced by the multiplication by  $x_0$*

$$0 \longrightarrow \frac{\mathbb{K}[x]}{I}(t-1) \xrightarrow{\cdot x_0} \frac{\mathbb{K}[x]}{I}(t) \longrightarrow \frac{\mathbb{K}[x]}{(I, x_0)}(t) \longrightarrow 0. \quad (1)$$

This regular sequence says that the ideal  $(I, x_0)$  has Hilbert polynomial  $\Delta p(t) = p(t) - p(t-1)$ . We will perform hyperplane section repeatedly so we define  $\Delta^0 p(t) := p(t)$  and recursively  $\Delta^k p(t) := \Delta^{k-1} p(t) - \Delta^{k-1} p(t-1)$ . It is very easy to manipulate Hilbert polynomial considering the Gotzmann decomposition.

**DEFINITION 2.2.** *An admissible Hilbert polynomial  $p(t)$  has a unique Gotzmann decomposition*

$$p(t) = \binom{t+a_1}{a_1} + \dots + \binom{t+a_r-(r-1)}{a_r}, \quad (2)$$

$a_1 \geq \dots \geq a_r$ . The number  $r$  of terms in this sum is said Gotzmann number of  $p(t)$ .

It can be easily proved that the Gotzmann decomposition of  $\Delta p(t)$  can be obtained from the decomposition of  $p(t)$  decreasing by 1 each index  $a_i$  and discarding the binomial

coefficients with a negative index below. Moreover we define an inverse operator  $\Sigma$  that associates to  $p(t)$  the Hilbert polynomial obtained increasing by 1 all the indices  $a_i$ .

**EXAMPLE 2.1.** *The Hilbert polynomial  $p(t) = 3t + 1$  has Gotzmann decomposition*

$$\binom{t+1}{1} + \binom{t+1-1}{1} + \binom{t+1-2}{1} + \binom{t-3}{0}.$$

Hence

$$\begin{aligned} \Delta p(t) &= \binom{t}{0} + \binom{t-1}{0} + \binom{t-2}{0} = 3 \\ \Sigma p(t) &= \binom{t+2}{2} + \binom{t+2-1}{2} + \binom{t+2-2}{2} + \binom{t+1-3}{1} = \\ &= \frac{3}{2}t^2 + \frac{5}{2}t - 1. \end{aligned}$$

We conclude the background materials with the properties linking the Hilbert polynomial to the regularity of an ideal.

**DEFINITION 2.3.** *A coherent sheaf  $\mathcal{F}$  over  $\mathbb{P}^n$  is said  $m$ -regular if for every  $i > 0$*

$$H^i(\mathcal{F}(m-i)) = 0. \quad (3)$$

The Castelnuovo-Mumford regularity of  $\mathcal{F}$  is the smallest integer  $m$ , for which  $\mathcal{F}$  is  $m$ -regular.

**THEOREM 2.4 (Gotzmann's Regularity Theorem).** *Let  $A$  be any  $\mathbb{K}$ -algebra and let  $Z \subset \text{Proj } A[x_0, \dots, x_n]$  be any subscheme with Hilbert polynomial  $p(t)$ , whose Gotzmann number is  $r$ . Then the sheaf of ideals  $\mathcal{I}_Z$  is  $r$ -regular.*

**PROPOSITION 2.5** ([8, PROPOSITION 2.6]). *Let  $I$  be a saturated ideal in  $\mathbb{K}[x]$ . The regularity of  $I$  is equal to the Castelnuovo-Mumford regularity of the sheaf of ideals  $\mathcal{I}$  obtained from the sheafification of  $I$ .*

**PROPOSITION 2.6** ([8, 1]). *The regularity of a Borel-fixed ideal  $I \subset \mathbb{K}[x_0, \dots, x_n]$  is equal to the maximal degree of one of its generators.*

### 3. MAIN PROPERTIES

Now we recall the technical properties that we need to guarantee the correctness of the algorithm, which are already proved in [5] and that we here adapt at the current notation.

**LEMMA 3.1.** *Let  $J \subset \mathbb{K}[x]$  be a saturated Borel-fixed with Hilbert polynomial  $p(t)$  whose Gotzmann number is  $r$ . Let  $x^\beta$  be a minimal monomial for  $\leq_B$  of  $\{J_r\} \subset \mathcal{P}(n, r)$  such that  $\min x^\beta = x_0$ . Then the ideal  $I = \langle \{J_r\} \setminus \{x^\beta\} \rangle^{\text{sat}}$  is Borel-fixed with Hilbert polynomial  $\bar{p}(t) = p(t) + 1$ .*

**PROOF.** First, note that by definition of minimal monomial,  $\{I_r\}$  is still a Borel set. Called  $q(t)$  the volume polynomial of  $J$ , we show that  $I$  has volume polynomial  $\bar{q}(t) = q(t) - 1$  applying Gotzmann's Persistence Theorem [8, Theorem 3.8], i.e. proving that

$$\dim_{\mathbb{K}} J_r - \dim_{\mathbb{K}} I_r = \dim_{\mathbb{K}} J_{r+1} - \dim_{\mathbb{K}} I_{r+1} = 1.$$

By construction  $\dim_{\mathbb{K}} J_r - \dim_{\mathbb{K}} I_r = 1$ . The Borel condition ensures that  $x^\beta x_0 \in J_{r+1} \setminus I_{r+1}$  and there are no other elements, because  $x^\beta x_0$  is the only monomial that cannot be generated from the monomials in  $I_r$  by multiplication of

a single variable. In fact let us consider the monomial  $x_i x^\beta$ ,  $i > 0$ . Since  $x_0 \mid x^\beta$  the following identity holds:

$$x_i x^\beta = \frac{x_i}{x_{i-1}} \cdots \frac{x_1}{x_0} x_0 x^\beta = \mathbf{e}_{i-1}^+ \circ \cdots \circ \mathbf{e}_0^+(x^\beta) x_0$$

and for each  $i$ ,  $\mathbf{e}_{i-1}^+ \circ \cdots \circ \mathbf{e}_0^+(x^\beta)$  belongs to  $I_r$ , by the minimality of  $x^\beta$ .  $\square$

**PROPOSITION 3.2** ([5, PROPOSITION 5.2]). *Consider a saturated Borel-fixed ideal  $J \subset \mathbb{K}[x]$  defining a subscheme with Hilbert polynomial  $p(t)$  whose Gotzmann number is  $r$ . Let  $I = J|_{x_1=x_0=1}$  be its double saturation and let  $\bar{p}(t)$  be the Hilbert polynomial of the subscheme defined by  $I$  in  $\mathbb{P}^n$ . Then*

$$\bar{p}(t) = p(t) - \dim_{\mathbb{K}} I_r + \dim_{\mathbb{K}} J_r. \quad (4)$$

**THEOREM 3.3.** *Let  $p(t)$  be an admissible Hilbert polynomial in  $\mathbb{P}^n$  and let  $r$  be an integer greater than or equal to the Gotzmann number of  $p(t)$ . There is a bijective function between the set  $\mathcal{I}_{p(t)}^n$  of the saturated Borel-fixed ideals of  $\mathbb{K}[x_0, \dots, x_n]$  with Hilbert polynomial  $p(t)$  and the set*

$$\mathcal{B}_{p(t)}^n = \left\{ \begin{array}{l} \mathcal{B} \subset \mathcal{P}(n, r) \text{ Borel set s.t.} \\ |\mathcal{B}_{(\geq i)}^c| = \Delta^i p(r), \forall i = 0, \dots, n-1 \end{array} \right\}.$$

More precisely

$$\begin{array}{ccc} \mathcal{I}_{p(t)}^n & \xleftarrow{1:1} & \mathcal{B}_{p(t)}^n \\ J & \longrightarrow & \{J_r\} \\ \langle \mathcal{B} \rangle^{\text{sat}} & \longleftarrow & \mathcal{B}. \end{array}$$

**PROOF.** First of all note that if the two maps are well-defined, i.e. for each  $J, J_{\geq r} = \langle J_r \rangle$  (Proposition 2.6 and Gotzmann's Regularity Theorem),

$$\begin{array}{ccc} J & \longrightarrow & \{J_r\} \longrightarrow \langle \{J_r\} \rangle^{\text{sat}} = J, \\ \mathcal{B} & \longrightarrow & \langle \mathcal{B} \rangle^{\text{sat}} \longrightarrow \{ \langle \mathcal{B} \rangle_r^{\text{sat}} \} = \mathcal{B}. \end{array}$$

Let  $J \subset \mathbb{K}[x]$  be a Borel-fixed with Hilbert polynomial  $p(t)$  and let  $\mathcal{N} = \mathcal{P}(n, r) \setminus \{J_r\}$ . Obviously  $|\mathcal{N}_{(\geq 0)}| = |\mathcal{N}| = p(t) = \Delta^0 p(r)$ . Using the short exact sequence (1), we determine the Borel ideal  $I = (J, x_0)^{\text{sat}} \cap \mathbb{K}[x_1, \dots, x_n]$  with Hilbert polynomial  $\Delta p(t)$ . Thus being  $\{I_r\} = \{J_r\}_{(\geq 1)} \subset \mathcal{P}(n-1, r)$ ,  $|\mathcal{N}_{(\geq 1)}| = |\{I_r\}^c| = \Delta p(r)$ . Since  $I$  is Borel-fixed in the polynomial ring  $\mathbb{K}[x_1, \dots, x_n]$  we can repeat the reasoning with the hyperplane section defined by  $x_1 = 0$  and so on.

Let us now consider a Borel set  $\mathcal{B} \subset \mathcal{P}(n, r)$ , such that the complement  $\mathcal{N} = \mathcal{B}^c$  satisfies the condition  $|\mathcal{N}_{(\geq i)}| = \Delta^i p(r)$  for every  $i$ . Firstly  $\text{reg}(\langle \mathcal{B} \rangle^{\text{sat}}) \leq r$  by Proposition 2.6, so let us prove that  $\langle \mathcal{B} \rangle^{\text{sat}}$  has Hilbert polynomial  $p(t)$ . We proceed by induction on the degree  $d$  of the Hilbert polynomial. For any  $n$ , if  $\deg p(t) = 0$ , then  $\mathcal{N}_{(\geq i)} = \emptyset$ , for every  $i \geq 1$ , since  $\Delta p(t) = 0$ , that is for any  $x^\beta \in \mathcal{N}$ ,  $\min x^\beta = 0$ . Applying repeatedly Lemma 3.1 starting from the Hilbert polynomial  $\bar{p}(t) = 0$  (corresponding to the ideal (1)), we obtain that  $\langle \mathcal{B} \rangle^{\text{sat}}$  defines a module  $\mathbb{K}[x]/\langle \mathcal{B} \rangle^{\text{sat}}$  having constant Hilbert polynomial  $p(t) (= p(r))$ .

Let us know suppose that the map  $\mathcal{B} \rightarrow \langle \mathcal{B} \rangle^{\text{sat}}$  is well-defined for any Hilbert polynomial of degree  $d-1$  and let  $p(t)$  be a Hilbert polynomial of degree  $d$ .  $\mathcal{B} = \mathcal{B}_{(\geq 1)} \subset \mathcal{P}(n-1, r)$  realizes the condition of the theorem w.r.t. the Hilbert polynomial  $\Delta p(t)$  and  $\deg \Delta p(t) = d-1$ . Hence by the inductive hypothesis the ideal  $\langle \mathcal{B} \rangle^{\text{sat}} \subset \mathbb{K}[x_1, \dots, x_n]$

has Hilbert polynomial  $\Delta p(t)$ . Let  $\bar{p}(t)$  be the Hilbert polynomial of  $\langle \mathcal{B} \rangle^{\text{sat}}$  in  $\mathbb{K}[x_0, \dots, x_n]$ :  $\bar{p}(t) = p(t) + a$ , because  $\Delta \bar{p}(t) = \Delta p(t)$ .  $\langle \mathcal{B} \rangle^{\text{sat}}$  turns out to be the  $x_1$ -saturation of  $\langle \mathcal{B} \rangle^{\text{sat}}$ , so by Proposition 3.2 the Hilbert polynomial of  $\mathbb{K}[x_0, \dots, x_n]/\langle \mathcal{B} \rangle^{\text{sat}}$  differs by a constant from  $\bar{p}(t)$  and since  $|\mathcal{N}| = |\mathcal{N}_{(\geq 0)}| = p(r)$  it coincides with  $p(t)$ .  $\square$

## 4. THE ALGORITHM

Therefore to compute the saturated Borel-fixed ideals we can construct Borel sets with the prescribed property. The proof of Theorem 3.3 suggests to use a recursive algorithm: i.e. to determine the Borel sets in  $\mathcal{P}(n, r)$  corresponding to the Hilbert polynomial  $p(t)$ , we begin computing the Borel sets in  $\mathcal{P}(n-1, r)$  corresponding to the Hilbert polynomial  $\Delta p(t)$ .

Let  $\bar{\mathcal{B}} \subset \mathcal{P}(n-1, r)$  a Borel set corresponding to the Hilbert polynomial  $\Delta p(t)$  and let  $\bar{\mathcal{N}} = \bar{\mathcal{B}}^c$ . In order for  $\bar{\mathcal{B}}$  to be the restriction  $\mathcal{B}_{(\geq 1)}$  of a Borel set  $\mathcal{B} \subset \mathcal{P}(n, r)$  (where  $\mathcal{P}(n, r)$  contains one more variable smaller than variables in  $\mathcal{P}(n-1, r)$ ), each monomial that can be obtained by decreasing moves from a monomial in  $\bar{\mathcal{N}}$  has to belong to  $\mathcal{N} = \mathcal{B}^c$ . This extension of an order set  $\bar{\mathcal{N}} \subset \mathcal{P}(n-1, r)$  to an order set  $\mathcal{N} \subset \mathcal{P}(n, r)$  has an ideal interpretation.

**LEMMA 4.1** ([5, LEMMA 5.1]). *Let  $\bar{\mathcal{B}} \subset \mathcal{P}(n-1, r)$  be a Borel set and let  $\bar{\mathcal{N}} = \bar{\mathcal{B}}^c$ . Moreover let  $\mathcal{N} \subset \mathcal{P}(n, r)$  be the order set containing the monomials in  $\bar{\mathcal{N}}$  and all those obtained by decreasing moves from them,*

$$\mathcal{N} = \mathcal{P}(n, r) \setminus \left\{ (\langle \bar{\mathcal{B}} \rangle^{\text{sat}} \cdot \mathbb{K}[x_0, \dots, x_n])_r \right\}. \quad (5)$$

**PROOF.** Let us call  $\mathcal{B}$  the Borel set containing the monomials of degree  $r$  belonging to the ideal  $\langle \bar{\mathcal{B}} \rangle^{\text{sat}} \cdot \mathbb{K}[x_0, \dots, x_n]$ . Let  $x^\alpha = x_n^{\alpha_n} \cdots x_0^{\alpha_0}$  be a monomial of  $\mathcal{P}(n, r)$  and suppose  $\min x^\alpha = 0$ , i.e.  $\alpha_0 > 0$ . The monomial  $\alpha_0 \mathbf{e}_0^+(x^\alpha) = x_n^{\alpha_n} \cdots x_1^{\alpha_1 + \alpha_0}$  belongs to  $\mathcal{P}(n-1, r)$ , so either  $\alpha_0 \mathbf{e}_0^+(x^\alpha) \in \bar{\mathcal{B}}$  or  $\alpha_0 \mathbf{e}_0^+(x^\alpha) \in \bar{\mathcal{N}}$ .

If  $x_n^{\alpha_n} \cdots x_1^{\alpha_1 + \alpha_0} \in \bar{\mathcal{B}}$ , then  $x_n^{\alpha_n} \cdots x_2^{\alpha_2}$  is in  $\langle \bar{\mathcal{B}} \rangle^{\text{sat}}$  and so  $x^\alpha \in \langle \bar{\mathcal{B}} \rangle^{\text{sat}} \cdot \mathbb{K}[x_0, \dots, x_n]$ , otherwise  $x_n^{\alpha_n} \cdots x_1^{\alpha_1 + \alpha_0} \in \bar{\mathcal{N}}$  implies  $x^\alpha \in \mathcal{N}$ .  $\square$

By Proposition 3.2, we know that the Hilbert polynomial corresponding to a Borel set  $\mathcal{B}$  of the type  $\{(\langle \bar{\mathcal{B}} \rangle^{\text{sat}} \cdot \mathbb{K}[x])_r\}$  differs from the target Hilbert polynomial by a constant: to determine this constant we compare the value  $p(r)$  of the Hilbert polynomial  $p(t)$  in degree  $r$  with the cardinality of the order set  $\mathcal{N}$  obtained by decreasing moves from  $\bar{\mathcal{N}}$ .

**LEMMA 4.2.** *Let  $\bar{\mathcal{N}} \subset \mathcal{P}(n-1, r)$  be an order set and let  $\mathcal{N} \subset \mathcal{P}(n, r)$  be the order set defined from  $\bar{\mathcal{N}}$  by decreasing moves. Then,*

$$|\mathcal{N}| = \sum_{\substack{x^\alpha \in \bar{\mathcal{N}} \\ x^\alpha = x_n^{\alpha_n} \cdots x_1^{\alpha_1}}} (\alpha_1 + 1). \quad (6)$$

**PROOF.** Each monomial  $x^\alpha \in \bar{\mathcal{N}}$  imposes the belonging to  $\mathcal{N}$  (in addition to itself) of the monomials obtained from it applying the decreasing moves  $\{\mathbf{e}_1^-, 2\mathbf{e}_1^-, \dots, \alpha_1 \mathbf{e}_1^-\}$ .  $\square$

There are three possibilities:

- $p(r) - |\mathcal{N}| < 0$ ,  $\bar{\mathcal{N}}$  imposes too many monomials outside the ideal, so the hyperplane section defined by  $\langle \bar{\mathcal{B}} \rangle^{\text{sat}}$  has to be discarded (there exist no Borel-fixed ideals corresponding to  $p(t)$  with such a hyperplane section);

- $p(r) - |\mathcal{N}| = 0$ ,  $\langle \overline{\mathcal{B}} \rangle^{\text{sat}} \subset \mathbb{K}[x_0, \dots, x_n]$  is one of the ideals sought;
- $p(r) - |\mathcal{N}| > 0$ , applying repeatedly Lemma 3.1 we determine the ideals we are looking for.

PROPOSITION 4.3 (Cf. [5, REMARK 5.3]). *Let  $p(t)$  be an admissible Hilbert polynomial with Gotzmann number  $r$  and let  $r_1$  be the Gotzmann number of  $\Delta p(t)$ .*

- (i) *Given a saturated Borel-fixed ideal  $J \subset \mathbb{K}[x_1, \dots, x_n]$  such that  $\mathbb{K}[x_1, \dots, x_n]/J$  has Hilbert polynomial  $\Delta p(t)$ , to pass from  $\{(J \cdot \mathbb{K}[x_0, \dots, x_n])_r\} \subset \mathcal{P}(n, r)$  to a Borel set corresponding to  $p(t)$ , we need to remove at most  $r - r_1$  monomials.*
- (ii) *We need to remove exactly  $r - r_1$  monomials if we consider the lexicographic ideal  $L \subset \mathbb{K}[x_1, \dots, x_n]$  corresponding to the polynomial  $\Delta p(t)$ .*

PROOF. (i) The minimal Hilbert polynomial having first difference equal to  $\Delta p(t)$  is  $\Sigma(\Delta p)(t)$  and the Gotzmann number of  $\Sigma(\Delta p)(t)$  coincides with the Gotzmann number of  $\Delta p(t)$ .

(ii) The Hilbert polynomial  $\Sigma(\Delta p)(t)$  is admissible so there exists the saturated lexicographic ideal  $L$  with such Hilbert polynomial. By construction  $L$  is the 1-lifting of the ideal defining its hyperplane section, therefore it has no generators involving the variable  $x_1$  and the ideal  $L \cap \mathbb{K}[x_1, \dots, x_n]$ , being generated by the same monomials, is still a lexicographic ideal.  $\square$

## 4.1 The pseudocode description

The recursive strategy naturally gives rise to a rooted tree where the nodes are all the Borel-fixed ideals met during the computation and the father of each node is the ideal of its hyperplane section (see Figure 1 for an example). This graph turns out to be a rooted tree because with a sequence of hyperplane section we obtain a unique ideal from any Borel-fixed ideal with the given Hilbert polynomial:

- the ideal  $(1) \subset \mathbb{K}[x_{d+1}, \dots, x_n]$  applying  $d+1$  sections, whenever the degree of the Hilbert polynomial  $d$  is smaller than  $n-1$ ;
- the ideal  $(x_n^c) \subset \mathbb{K}[x_{n-1}, x_n]$ , where  $c = \Delta^d p(r)$ , if  $d = n-1$ .

The leaves are the Borel-fixed ideal we are looking for and by definition this tree will have height  $d+2$  if  $d < n-1$  and  $d+1$  if  $d = n-1$ .

To compute the Borel-fixed ideals in  $\mathbb{K}[x_0, \dots, x_n]$  with Hilbert polynomial  $p(t)$ , the algorithm described in [5] requires to compute the ideals in  $\mathbb{K}[x_1, \dots, x_n]$  with Hilbert polynomial  $\Delta p(t)$  and so on, until the last recursive call that requires to compute the ideal representing the root of the tree. Now, starting from the root of the tree, the algorithm performs a Breadth-First-Search visit. This approach is not optimal, because it requires much space in the memory to store also all the ideal which are not leaves at maximal height and so which are necessary only temporarily. Moreover some of such ideals can be discarded immediately because too many conditions are imposed when embedded in a polynomial ring with one more variable. For this reason we prefer an algorithm that visits the tree with a Depth-First-Search approach. The function BORELFIXEDIDEALSGENERATOR (Algorithm 1) initializes the computation determining the root of the tree and starting the DFS visit.

```

1: BORELFIXEDIDEALSGENERATOR( $\mathbb{K}[x_0, \dots, x_n], p(t)$ )
Input:  $\mathbb{K}[x_0, \dots, x_n]$ , polynomial ring.
Input:  $p(t)$ , admissible Hilbert polynomial in  $\mathbb{P}^n$ .
Output: the set of all Borel-fixed ideals in  $\mathbb{K}[x_0, \dots, x_n]$ 
        with Hilbert polynomial  $p(t)$ .
2:  $d \leftarrow \deg p(t)$ ;
3:  $r \leftarrow \text{GOTZMANNNUMBER}(p(t))$ ;
4: if  $d = n - 1$  then
5:    $c \leftarrow \Delta^d p(t)$ ;
6:    $\mathcal{B} \leftarrow \{(x_n^c)_r\} \subset \mathcal{P}(2, r)$ ; //  $\mathbb{K}[x_{n-1}, x_n]$ .
7:   return BORELIDEALS( $\mathbb{K}[x_0, \dots, x_n], p(t), d, \mathcal{B}$ );
8: else
9:    $\mathcal{B} \leftarrow \mathcal{P}(n - d, r)$ ; //  $\mathbb{K}[x_{d+1}, \dots, x_n]$ .
10:  return BORELIDEALS( $\mathbb{K}[x_0, \dots, x_n], p(t), d + 1, \mathcal{B}$ );
11: end if

```

**Algorithm 1:** The procedure that initializes the computation detecting the root of the tree.

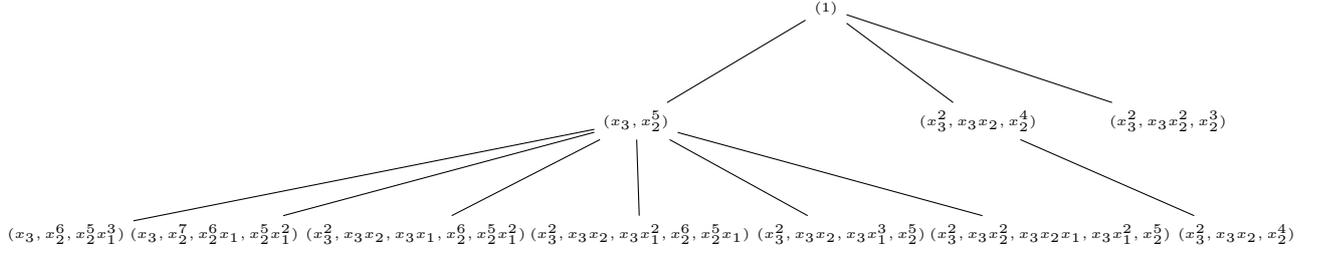
The core of the algorithm is described by the function BORELIDEALS (Algorithm 2). It takes a Borel set  $\mathcal{B}$  corresponding to a Borel-fixed ideal in  $\mathbb{K}[x_k, \dots, x_n]$  with Hilbert polynomial  $\Delta^k p(t)$ . If  $k = 0$  then  $\mathcal{B}$  corresponds to one of the searched ideal, so the function returns the saturation of the ideal generated by the monomials in  $\mathcal{B}$ . Otherwise, the function embeds  $\mathcal{B}$  in a poset with one more variable obtaining  $\mathcal{B}$ , it computes how many monomials do not belong to  $\mathcal{B}$  and how many monomials ( $q$ ) have to be removed. Then it calls the function REMOVE to compute all the Borel set that can be obtained from  $\mathcal{B}$  removing  $q$  monomials and finally for each of such new Borel set we have the recursive call with index  $k-1$ .

```

1: BORELIDEALS( $\mathbb{K}[x_0, \dots, x_n], p(t), k, \overline{\mathcal{B}}$ )
Input:  $\mathbb{K}[x_0, \dots, x_n]$ , polynomial ring.
Input:  $p(t)$ , admissible Hilbert polynomial in  $\mathbb{P}^n$ .
Input:  $k$ , integer s.t.  $0 \leq k \leq \deg p(t)$ .
Input:  $\overline{\mathcal{B}}$ , Borel set in  $\mathcal{P}(n - k, r)$  s.t. the ideal  $\langle \overline{\mathcal{B}} \rangle^{\text{sat}}$  in
         $\mathbb{K}[x_k, \dots, x_n]$  has Hilbert polynomial  $\Delta^k p(t)$ .
Output: the set of all saturated Borel-fixed ideals  $J$ 
        in  $\mathbb{K}[x_0, \dots, x_n]$  with Hilbert polynomial  $p(t)$  s.t.
         $\{J_r\}_{(r \geq k)} = \overline{\mathcal{B}}$ .
2: if  $k = 0$  then
3:   return  $\langle \overline{\mathcal{B}} \rangle^{\text{sat}}$ ;
4: else
5:    $\mathcal{B} \leftarrow \{(\overline{\mathcal{B}})^{\text{sat}} \cdot \mathbb{K}[x_{k-1}, \dots, x_n]\}_r \subset \mathcal{P}(n - k + 1, r)$ ;
6:    $q \leftarrow \Delta^{k-1} p(r) - |\mathcal{B}^c|$ ;
7:   if  $q \geq 0$  then
8:     oneMoreVariable  $\leftarrow$  REMOVE( $\mathcal{B}, q, 1$ );
9:     ideals  $\leftarrow \emptyset$ ;
10:    for all  $\tilde{\mathcal{B}} \in$  oneMoreVariable do
11:      ideals  $\leftarrow$  ideals  $\cup$ 
        BORELIDEALS( $\mathbb{K}[x_0, \dots, x_n], p(t), k - 1, \tilde{\mathcal{B}}$ );
12:    end for
13:    return ideals;
14:  else
15:    return  $\emptyset$ ;
16:  end if
17: end if

```

**Algorithm 2:** The core of the recursive strategy to compute Borel-fixed ideals in  $\mathbb{K}[x]$  with Hilbert polynomial  $p(t)$ .



**Figure 1:** The tree associated at the computation of the Borel-fixed ideals in  $\mathbb{K}[x_0, x_1, x_2, x_3]$  with Hilbert polynomial  $p(t) = 5t - 2$ .

Also the function REMOVE (Algorithm 3) uses a recursive strategy, i.e. given a Borel set  $\mathcal{B}$  and  $q$  monomials to remove, for each minimal element  $x^\alpha$  of  $\mathcal{B}$ , it calls itself on the Borel set  $\mathcal{B} \setminus \{x^\alpha\}$  to which we have to remove  $q - 1$  monomials. Whenever we consider a Borel set  $\mathcal{B}$  with (at least) two minimal elements  $x^\alpha$  and  $x^\gamma$  and we need to remove (at least) two monomials, the strategy just described would generate two times the same Borel set because  $(\mathcal{B} \setminus \{x^\alpha\}) \setminus \{x^\gamma\} = (\mathcal{B} \setminus \{x^\gamma\}) \setminus \{x^\alpha\}$ . To avoid this repetition we add, as argument of the function REMOVE, a monomial  $x^\beta$  and we require that any monomial removed from  $\mathcal{B}$  is greater than  $x^\beta$  w.r.t. a fixed term ordering (in Algorithm 3 we chose  $\text{DegLex}$ ). So whenever the function is called by BORELIDEALS we pass as argument the monomial 1 (any removal is admissible), whereas when REMOVE is called by itself, that is some removal has been already performed, we pass as argument the last monomial removed. In this way, assuming  $x^\alpha >_{\text{DegLex}} x^\gamma$ , the function REMOVE will only generate the Borel set  $(\mathcal{B} \setminus \{x^\gamma\}) \setminus \{x^\alpha\}$ .

```

1: REMOVE( $\mathcal{B}, q, x^\beta$ )
Input:  $\mathcal{B}$ , a Borel set.
Input:  $q$ , a non-negative integer.
Input:  $x^\beta$ , a monomial.
Output: the set of all Borel sets obtained from  $\mathcal{B}$  removing
in all the possible ways  $q$  monomials. The monomial  $x^\beta$ 
is to avoid repetitions: it will be 1 when the function
is called by BORELIDEALS, whereas it will be the last
monomial removed whenever the function is called by
itself.
2: if  $q = 0$  then
3:   return  $\{\mathcal{B}\}$ ;
4: else
5:   borelSets  $\leftarrow \emptyset$ ;
6:   minimalMonomials  $\leftarrow \text{MINIMALELEMENTS}(\mathcal{B})$ ;
7:   for all  $x^\alpha \in \text{minimalMonomials}$  do
8:     if  $x^\alpha >_{\text{DegLex}} x^\beta$  then
9:       borelSets  $\leftarrow \text{borelSets} \cup$ 
REMOVE( $\mathcal{B} \setminus \{x^\alpha\}, q - 1, x^\alpha$ );
10:    end if
11:  end for
12:  return borelSets;
13: end if

```

**Algorithm 3:** The function to remove monomials from a Borel set.

## 4.2 The implementation

<p>GOTZMANNNUMBER(<math>p(t)</math>)  <b>Input:</b> <math>p(t)</math>, a Hilbert polynomial.  <b>Output:</b> the Gotzmann number of <math>p(t)</math>.</p> <p>MINIMALELEMENTS(<math>\mathcal{B}</math>)  <b>Input:</b> <math>\mathcal{B}</math>, a Borel set.  <b>Output:</b> the set of minimal elements of <math>\mathcal{B}</math>.</p>
--

**Algorithm 4:** Auxiliary functions

The key of an efficient implementation is how we realize the Borel set. The guidelines are

- slim structure, in order to take up less memory as possible (since the number of final objects can be huge);
- quick implementation of the basic operations:
  1. embedding of a Borel set in a poset with one more variable (lines 5-6 Algorithm 2);
  2. computation of the minimal elements of a Borel set (line 6 Algorithm 3);
  3. removal of a monomial from a Borel set (line 9 Algorithm 3);
  4. computation of the saturation of the ideal generated by a Borel set (line 3 Algorithm 2).

The most compact way to describe a Borel set is to store its minimal elements. This is efficient also to evaluate the belonging to the Borel set of any other monomial, because we can use the following proposition.

**PROPOSITION 4.4** ([10, PROPOSITION 2.3]). *Let  $x^\alpha, x^\beta$  be two monomials in  $\mathbb{K}[x_0, \dots, x_n]_r$ .*

$$x^\alpha >_B x^\beta \iff \sum_{j=i}^n (\alpha_j - \beta_j) \geq 0, \forall i = 0, \dots, n. \quad (7)$$

Indeed if a monomial  $x^\gamma$  belongs to the Borel set, there exists a minimal element  $x^\alpha$  such that  $x^\gamma \geq_B x^\alpha$ , whereas if  $x^\gamma$  does not belong there exists a minimal element  $x^\beta$  such that  $x^\gamma <_B x^\beta$ .

In this way, the computation of the set  $S$  of minimal elements is immediate. When a minimal monomial  $x^\alpha$  is removed, we computed all the monomials that can be obtained from  $x^\alpha$  with an increasing elementary move and for each of them we check if it belongs to the Borel set corresponding to

$S \setminus \{x^\alpha\}$ : if not we add such monomial to the set of minimal elements.

The description of the Borel set by means of its minimal elements turns out to be efficient also when we need to pass to a poset with one more variable. Indeed given the Borel set  $\mathcal{B} \subset \mathcal{P}(n-k, r)$  with minimal elements stored in the set  $\overline{M}$ , by Lemma 4.1 we know that the Borel set  $\mathcal{B} = \{(\langle \mathcal{B} \rangle^{\text{sat}} \cdot \mathbb{K}[x_{k-1}, \dots, x_n])_r\} \subset \mathcal{P}(n-k+1, r)$  is described by the set of minimal elements

$$M = \{x_n^{\alpha_n} \cdots x_{k+1}^{\alpha_{k+1}} x_k^{\alpha_k} \mid x_n^{\alpha_n} \cdots x_{k+1}^{\alpha_{k+1}} x_k^{\alpha_k} \in \overline{M}\}. \quad (8)$$

A not trivial task is to compute the number of monomials that do not belong to the Borel set after the embedding in a poset with one more variable. In principle it consists in the computation of the Hilbert polynomial and in its evaluation at degree  $r$ . But the algebraic approach (used in [12]) slows down the algorithm because any software dedicated to the study of polynomial ideals computes the Hilbert polynomial with Gröbner basis tools, which in this case are unnecessary.

Given a Borel set  $\mathcal{B} \subset \mathcal{P}(n-k, r)$ , let  $\mathcal{N} = \mathcal{B}^c$ . If we decompose  $\mathcal{N}$  according to the power of the smallest variable  $x_k$ , i.e.

$$\overline{\mathcal{N}} = \prod_{i=0}^r \overline{\mathcal{N}}_i, \quad \overline{\mathcal{N}}_i = \{x^\alpha \in \overline{\mathcal{N}} \text{ s.t. } \alpha_k = i\}. \quad (9)$$

We can compute the cardinality of  $\mathcal{N} = \mathcal{B}^c$  rewriting (6) of Lemma 4.2 as

$$\begin{aligned} |\mathcal{N}| &= \sum_{\substack{x^\alpha \in \overline{\mathcal{N}} \\ x^\alpha = x_n^{\alpha_n} \cdots x_k^{\alpha_k}}} (\alpha_k + 1) = \\ &= \sum_{i=0}^r \sum_{\substack{x^\alpha \in \overline{\mathcal{N}}_i \\ x^\alpha = x_n^{\alpha_n} \cdots x_k^i}} (i+1) = \sum_{i=0}^r (i+1) |\overline{\mathcal{N}}_i|. \end{aligned}$$

Hence we add to the structure describing a Borel set an array of  $r+1$  integers, such that the  $i$ -th index is equal to the number of monomials not belonging to the Borel set with the power of the smallest variable equal to  $i$ . It is easy also to deduce the array corresponding to a Borel set after the embedding in a poset with one more variable, indeed any  $x^\alpha \in \overline{\mathcal{N}}_i$  implies the belonging to  $\mathcal{N}$  of  $i+1$  monomials

$$\begin{aligned} x^\alpha &= x_n^{\alpha_n} \cdots x_k^{\alpha_k} x_{k-1}^0 \in \mathcal{N}_0, \\ \mathbf{e}_k^-(x^\alpha) &= x_n^{\alpha_n} \cdots x_k^{\alpha_k-1} x_{k-1} \in \mathcal{N}_1, \\ &\vdots \\ \alpha_k \mathbf{e}_k^-(x^\alpha) &= x_n^{\alpha_n} \cdots x_k^0 x_{k-1}^{\alpha_k} \in \mathcal{N}_{\alpha_k} \end{aligned}$$

so that

$$|\mathcal{N}_j| = \sum_{i=0}^j |\overline{\mathcal{N}}_i|.$$

Moreover when a minimal monomial  $x^\alpha$  is removed from a Borel set, we increase by one the index of the array corresponding to the power of the smallest variable in  $x^\alpha$ .

The last operation which has to be as quick as possible is the computation of the saturation of the ideal. We exploit the algebraic approach of [12] and we add to the structure describing the Borel set the list of monomials that generate the saturated ideal corresponding to the Borel set. If

$x^\alpha = x_n^{\alpha_n} \cdots x_k^{\alpha_k}$  is a minimal monomial of  $\mathcal{B}$  then  $x^\alpha = x^\alpha|_{x_k=1} = x_n^{\alpha_n} \cdots x_{k+1}^{\alpha_{k+1}}$  is a generator of the saturated ideal. As shown in [12], removing  $x^\alpha$  from  $\mathcal{B}$  implies that  $\langle \mathcal{B} \setminus \{x^\alpha\} \rangle^{\text{sat}}$  is generated by the same generators of  $\langle \mathcal{B} \rangle^{\text{sat}}$  except for  $x^\alpha$  that is replaced by

$$x^\alpha \cdot x_m, \dots, x^\alpha \cdot x_1, \quad \text{where } x_m = \min x^\alpha.$$

### 4.3 Experimental results

The following experimental results are obtained with an implementation of the algorithm realized with the guidelines showed above, coded in `java` and ran on a MacBook Pro with a 2.4 GHz Intel Core 2 Duo processor. It can be tested by means of an applet available at [www.personalweb.unito.it/paolo.lella/HSC/borelGenerator.html](http://www.personalweb.unito.it/paolo.lella/HSC/borelGenerator.html). In the following tables we consider Hilbert polynomials of degree 0,1,2 (i.e. points, curves and surfaces) and projective spaces of increasing dimension and we reported the elapsed time (in seconds) of the computation and the number of Borel-fixed ideals obtained.

Time	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$p(t) = 5$	0.101	0.025	0.017	0.021
$p(t) = 10$	0.062	0.064	0.119	0.048
$p(t) = 15$	0.079	0.225	0.298	0.401
$p(t) = 20$	0.341	1.595	2.735	3.870
$p(t) = 25$	2.094	13.595	24.497	33.303

Ideals	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$p(t) = 5$	5	5	5	5
$p(t) = 10$	42	50	50	50
$p(t) = 15$	287	417	425	425
$p(t) = 20$	1732	3130	3263	3271
$p(t) = 25$	9501	21616	23158	23291

**Table 1: Experiments with constant Hilbert polynomials. The Gotzmann number coincides with the number of points.**

Time	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$5t + 1$ (11)	0.117	0.159	0.067	0.0621
$5t + 7$ (17)	0.502	1.480	2.312	3.290
$5t + 13$ (23)	10.513	56.652	91.456	128.341
$8t - 6$ (22)	0.987	2.623	4.138	5.852
$8t - 3$ (25)	3.008	14.128	22.960	32.300
$8t$ (28)	12.960	72.053	273.719	238.856

Ideals	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$5t + 1$ (11)	89	98	98	98
$5t + 7$ (17)	3028	4560	4587	4587
$5t + 13$ (23)	58124	123689	126962	127030
$8t - 6$ (22)	4171	6741	6837	6838
$8t - 3$ (25)	17334	32073	32848	32868
$8t$ (28)	68291	144660	149777	149976

**Table 2: Experiments with Hilbert polynomials of curves. The Gotzmann number is reported in brackets.**

Time	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$2t^2 + 8t - 46$ (16)	0.312	0.189	0.304	0.516
$2t^2 + 8t - 42$ (20)	0.103	0.338	0.558	0.883
$2t^2 + 8t - 38$ (24)	0.741	3.167	5.237	7.055
$4t^2 - 12t + 10$ (20)	0.147	0.280	0.377	0.561
$4t^2 - 12t + 14$ (24)	0.953	3.909	6.007	8.588
$4t^2 - 12t + 18$ (28)	9.066	50.071	82.592	112.237

Ideals	$n = 5$	$n = 10$	$n = 15$	$n = 20$
$2t^2 + 8t - 46$ (16)	834	38	38	38
$2t^2 + 8t - 42$ (20)	481	670	671	671
$2t^2 + 8t - 38$ (24)	4774	8393	8476	8476
$4t^2 - 12t + 10$ (20)	631	856	857	857
$4t^2 - 12t + 14$ (24)	6394	10986	11082	11082
$4t^2 - 12t + 18$ (28)	51527	112852	115295	115332

**Table 3:** Experiments with Hilbert polynomials of degree 2. The Gotzmann number is reported in brackets.

## 5. REFERENCES

- [1] D. Bayer and M. Stillman. A criterion for detecting  $m$ -regularity. *Invent. Math.*, 87(1):1–11, 1987.
- [2] C. Bertone, F. Cioffi, P. Lella, and M. Roggero. Upgraded methods for the effective computation of marked schemes on a strongly stable ideal. *ArXiv e-prints*, 2011. Available at [arxiv.org/abs/1110.0698](https://arxiv.org/abs/1110.0698).
- [3] C. Bertone, P. Lella, and M. Roggero. Borel open covering of Hilbert schemes. *ArXiv e-prints*, 2011. Available at [arxiv.org/abs/0909.2184](https://arxiv.org/abs/0909.2184).
- [4] J. Brachat, P. Lella, B. Mourrain, and M. Roggero. Low degree equations defining the Hilbert scheme. *ArXiv e-prints*, 2011. Available at [arxiv.org/abs/1104.2007](https://arxiv.org/abs/1104.2007).
- [5] F. Cioffi, P. Lella, M. G. Marinari, and M. Roggero. Segments and Hilbert schemes of points. *Discrete Mathematics*, 311(20):2238 – 2252, 2011.
- [6] F. Cioffi and M. Roggero. Flat families by strongly stable ideals and a generalization of Gröbner bases. *J. Symb. Comput.*, 46(9):1070–1084, 2011.
- [7] G. Gotzmann. Eine Bedingung für die Flachheit und das Hilbertpolynom eines graduierten Ringes. *Math. Z.*, 158(1):61–70, 1978.
- [8] M. L. Green. Generic initial ideals. In *Six lectures on commutative algebra*, Mod. Birkhäuser Class., pages 119–186. Birkhäuser Verlag, Basel, 2010.
- [9] R. Hartshorne. Connectedness of the Hilbert scheme. *Inst. Hautes Études Sci. Publ. Math.*, (29):5–48, 1966.
- [10] P. Lella. A network of rational curves on the Hilbert scheme. *ArXiv e-prints*, 2010. Available at [arxiv.org/abs/1006.5020](https://arxiv.org/abs/1006.5020).
- [11] P. Lella and M. Roggero. Rational components of Hilbert schemes. *Rend. Semin. Mat. Univ. Padova*, 126:11–45, 2011.
- [12] D. Moore and U. Nagel. Algorithms for strongly stable ideals. *ArXiv e-prints*, 2011. Available at [arxiv.org/abs/1110.4080](https://arxiv.org/abs/1110.4080).