



Characterization of $h\nu$ -Convex Sequences

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Abstract

Reconstructing a discrete object by means of X-rays along a finite set U of (discrete) directions represents one of the main task in discrete tomography. Indeed, it is an ill-posed inverse problem, since different structures exist having the same projections along all lines whose directions range in U . Characteristic of ambiguous reconstructions are special configurations, called switching components, whose understanding represents a main issue in discrete tomography, and an independent interesting geometric problem as well. The investigation of switching component usually bases on some kind of prior knowledge that is incorporated in the tomographic problem. In this paper, we focus on switching components under the constraint of convexity along the horizontal and the vertical directions imposed to the unknown object. Moving from their geometric characterization in windows and curls, we provide a numerical description, by encoding them as special sequences of integers. A detailed study of these sequences leads to the complete understanding of their combinatorial structure, and to a polynomial-time algorithm that explicitly reconstructs any of them from a pair of integers arbitrarily given.

Keywords Curl · Discrete tomography · $h\nu$ -convex set · polyomino · Projection · Switching-component · window · X-ray

Mathematics Subject Classification 52A30 · 68R01 · 52C30 · 52C45

1 Introduction

Reconstruction of finite discrete sets having prescribed projections is a challenging well-known inverse problem. In many cases, it shares its domain with Physics, Chemistry, Medical imaging and other areas where discrete models of real objects are employed. In discrete tomography, the object to be reconstructed is represented by a set of parallel slices, each modeling as a binary matrix the set of points of the object lying on the corresponding 2D plane. So, the basic problem is the reconstruction of a binary matrix from the collected projection data, that are obtained by means of X-rays in a set U of different directions.

Since the discrete sets consistent with a given set of projections are, in general, a huge number, and these can be

mutually very different, further constraints are required to select only those solutions that match the structures we are concerned with.

In the literature, the connectedness constraints have been primarily investigated, and many meaningful results have been obtained, both concerning the reconstruction and the enumeration of possible solutions in such a restricted space (see [24,25] for an overview). Differently, one can confine the investigation to a given lattice grid \mathcal{A} , so looking for uniqueness conditions inside \mathcal{A} , or in suitable subregions of \mathcal{A} (a non-exhaustive list of useful papers is [9,14–16,22,23,26,28], also including the corresponding reference sections).

In case different discrete sets Y_1 and Y_2 are *tomographically equivalent* with respect to a set U of directions, namely when Y_1, Y_2 can be reconstructed by means of the same X-rays with respect to U , then there exist specific patterns, called *switching components* which turn Y_1 into Y_2 . So, understanding the combinatorial and the geometric structure of the switching components is a main issue for a faithful reconstruction, largely investigated in discrete tomography, also in terms of *ghosts* or *bad configurations* (see, for instance [3,5,6,8,9,11,16,18–21,23,27]).

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The aim of the present paper falls in this area of research, and concerns finite discrete sets that are horizontally and vertically convex, and that share the same projections along those preferred directions, i.e., the horizontal and the vertical ones.

In particular, we are interested in the characterization of the switching components having a special geometrical condition, which defines a class of patterns called *hv-convex switching components*. It includes the classes of *regular* and of *irregular* switching components considered in [21], that, following [13], can be also defined in terms of *hv-convex windows* and *hv-convex curls*, respectively.

We show that these *hv-convex* switchings can be encoded by a sequence of integers. We provide a complete description of the combinatorial structures of such sequences, and an algorithm for their explicit constructions.

The paper extends and generalizes the results in [13]. It is structured as follows: in Sect. 2 we give the notations, and we provide the basic knowledge on switching components. We introduce windows and curls, then adding to both of them the *hv-convex* constraint, which leads to the basic Definition 5, where the notion of *hv-sequence* is stated. Section 3 contains one of the main results of our study, i.e. Theorem 17, which provides the characterization of the entries of the *hv-convex* sequences according to the number of points of the related *hv-convex* switching. In [13], only the easier case of *hv-sequences* related to windows had been focused, with a few sporadic examples concerning allowed or forbidden *hv-sequences* of curls. This led to [13, Problem 1], which asked for the characterization of all *hv-sequences*. By Theorem 17, we have completely solved Problem 1, and, as a particular case, we get also the characterization of the *hv-sequences* having just one odd repeated entry (Corollary 18), which provides the curl counterpart of [13, Theorem 3], holding for windows. In Sect. 4, we deeper investigate the geometric and the combinatorial structure of the *hv-convex* switchings. Thanks to Theorem 25 and Algorithm 1, *CompSeq*, we characterize and compute the exact positions, inside the *hv-sequence*, of the allowed entries, which frames in the proper context the sporadic cases considered in [13]. In Sect. 5, we provide a few concluding remarks, and outline some hints for possible further developments.

2 Notations and Preliminaries

We adopt the standard notation \mathbb{Z}^2 to indicate the lattice of points of the Euclidean two-dimensional space having integer coordinates. For each point $v \in \mathbb{Z}^2$, we denote by $L_h(v)$ and $L_v(v)$ the horizontal and vertical discrete lines intersecting v , respectively. A set of points $A \subset \mathbb{Z}^2$, considered up to translations, is *horizontally* (resp. *vertically*) *convex* if, for each $v \in A$, the set $L_h(v) \cap A$ (resp. $L_v(v) \cap A$) is connected.

A set that is both horizontally and vertically convex is shortly called *hv-convex*. A *polyomino* is a finite set of points in \mathbb{Z}^2 whose elements are 4-connected.

For a finite set A , we define its *horizontal* (resp. *vertical*) *projection* to be the integer vector $H(A)$ (resp. $V(A)$) counting the number of points of A that lie on each horizontal (resp. vertical) line intersecting A .

Given a point $v = (i, j) \in \mathbb{Z}^2$, the four following closed regions are defined (with the same notations as in [12,17]):

$$Z_0(v) = \{(i', j') \in \mathbb{Z}^2 : i' \leq i, j' \leq j\},$$

$$Z_1(v) = \{(i', j') \in \mathbb{Z}^2 : i' \geq i, j' \leq j\},$$

$$Z_2(v) = \{(i', j') \in \mathbb{Z}^2 : i' \geq i, j' \geq j\},$$

$$Z_3(v) = \{(i', j') \in \mathbb{Z}^2 : i' \leq i, j' \geq j\}.$$

A set $A \subset \mathbb{Z}^2$ is said to be *quadrant convex*, briefly *Q-convex*, (along the horizontal and vertical directions) if $Z_k(v) \cap A \neq \emptyset$ for all $k = 0, 1, 2, 3$ implies $v \in A$. Figure 1 shows simple examples of the introduced sets.

Lemma 1 *Let A be an *hv-convex* polyomino, and consider a point $v \in \mathbb{Z}^2$. If $w_1, w_2, w_3 \in A$ exist such that $Z_i(v) \cap \{w_1, w_2, w_3\} \neq \emptyset$ for all $i = 0, 1, 2, 3$, then $v \in A$.*

Proof By [7, Proposition 2.3], an *hv-convex* connected set is also *Q-convex*. The statement follows immediately by the *hv-convex* property of A . \square

2.1 Switching Components and the Uniqueness Problem

The following definition introduces the class of *hv-switchings*.

Definition 2 A pair $S = (S^0, S^1)$ of sets of points is a *hv-switching* if:

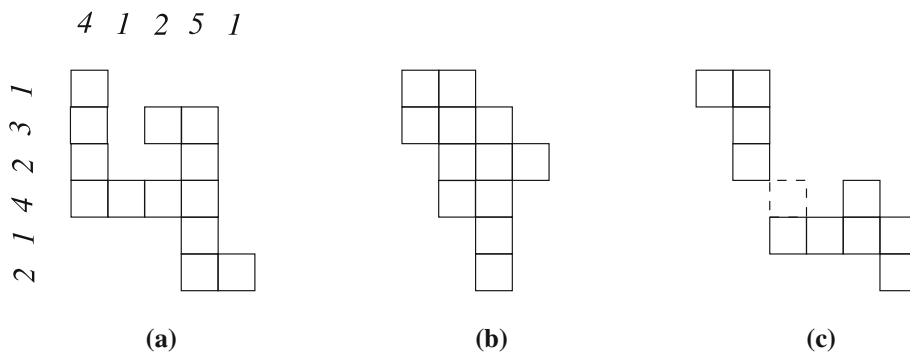
- $S^0 \cap S^1 = \emptyset$ and $|S^0| = |S^1|$;
- $H(S^0) = H(S^1)$ and $V(S^0) = V(S^1)$, i.e., S^0 and S^1 have the same horizontal and vertical projections.

Each set S^0 and S^1 is indicated as *hv-switching component*.

A discrete set A contains a *hv-switching component* if $S^0 \subseteq A$ and $S^1 \cap A = \emptyset$. In this case, we consider $A = Y \cup S^0$, with Y being a (possibly empty) discrete set; we define the set $A' = Y \cup S^1$ as the *dual* of A , and we say that the switching S is *associated* to A and A' .

A classical result in [27] states that if A_1 and A_2 are two discrete sets sharing the same horizontal and vertical projections, then A_2 is the dual of A_1 with respect to a *hv-switching*. So, for any point $v \in S^0$ (resp. $v \in S^1$), there exist points $w_1, w_2 \in S^1$ (resp. $w_1, w_2 \in S^0$) such that $w_1 \in L_h(v)$ and $w_2 \in L_v(v)$.

Fig. 1 **a** A discrete set, and the corresponding horizontal and vertical projections, that is neither horizontally nor vertically convex. **b** An hv -convex set. **c** A convex set that is not Q -convex (the dashed cell does not belong to the set)



If A_1 and A_2 are hv -convex sets, then, due to Lemma 1, for any $v \in S$ there exists one and only one $i \in \{0, 1, 2, 3\}$ such that $Z_i(v) \cap S$ consists of points all belonging to the same component of S as v . The quadrant $Z_i(v)$ is said to be the *free region* of v , or the *S -free region* of v in case we wish to emphasize that the free region relates to the switching S . We denote by $F(v)$ (or by $F_S(v)$) the free region of $v \in S$. Also, $F_i(S)$ denotes the subset of S consisting of all points having free region $Z_i(v)$, namely $F_i(S) = \{v \in S, F_S(v) = Z_i(v)\}$, $i \in \{0, 1, 2, 3\}$.

We have the following result from [12]

Lemma 3 Let $S = (S^0, S^1)$ be a hv -switching. Then, the following conditions are equivalent

$$\bigcup_{i=0}^3 F_i(S) = S. \tag{1}$$

$$\begin{aligned} v, w \in F_i(S), i \in \{0, 1, 2, 3\}, v \in S^0, \\ w \in S^1 \Rightarrow v \notin Z_j(w), \\ w \notin Z_j(v), j = i + 2 \pmod{4}. \end{aligned} \tag{2}$$

Let $S = (S^0, S^1)$ be a hv -switching.

The following definition introduces the subclass of hv -switching we wish to deal with.

Definition 4 A hv -switching S is said to be a hv -convex switching if the equivalent conditions of Lemma 3 hold.

In the literature, hv -convex switchings have been deeply studied when related to hv -convex polyominoes (for basic definitions and main results see [24,25]). In this context, it holds that if $S = (S^0, S^1)$ is a hv -switching associated to a pair of hv -convex polyominoes, then (1) holds, so S is an hv -convex switching. However, the converse is not necessarily true, namely two polyominoes P_1 and P_2 can exist such that one is the dual of the other with respect to an hv -convex switching S , and such that one or both of them are not hv -convex. An interesting case is Fig. 23 in [21], or Fig. 2.

In the sequel, we focus on the problem of characterizing the class of hv -convex switchings by integers' sequences, say hv -sequences, as defined below. Let us start by recalling the

following classification of hv -convex switchings introduced in [12,13].

A closed polygonal curve K in \mathbb{R}^2 is said to be a *squared spiral* if K consists of segments having, alternatively, horizontal and vertical direction. Their endpoints form the *set of vertices* of the polygonal. Assume to travel K according to a prescribed orientation. A vertex v of K is said to be a *counterclockwise point* if, crossing v , implies a counterclockwise change of direction. Differently, v is a *clockwise point*. Of course, by reversing the traveling orientation, clockwise and counterclockwise vertices mutually exchange.

Windows and curls.

A squared spiral W is said to be a *window* if it can be traveled by turning always clockwise, or always counterclockwise. Otherwise, the squared spiral is said to be a *curl*. Therefore, traveling a curl needs changes of turning direction. From now on, we will consider spirals whose vertices have integer coordinates.

We underline that the set of vertices of a window or a curl forms a hv -switching $S = (S^0, S^1)$ by considering the vertices alternatively belonging to S^0 and S^1 .

Figure 3 shows an example of window and curl pointing out the elements belonging to different hv -switching components.

Let S be a squared spiral (window or curl). Suppose to select any vertex v , and start traveling S moving from v clockwise (or counterclockwise). So, visiting one vertex after the other, we can group them in ordered lists of consecutive clockwise and counterclockwise points.

Definition 5 Once an orientation of a squared spiral S has been chosen, let v be a vertex where the orientation changes. Starting from v , compute the lengths k_1, k_2, \dots, k_n of the lists of consecutive vertices having, alternatively, the same and the opposite orientation as v . The resulting sequence of integers $\pi = (k_1, k_2, \dots, k_n)$ is said to be the *hv -sequence* associated to S .

The sequence π is considered cyclically arranged, namely its first entry can be selected arbitrarily among k_1, \dots, k_n , and then listing the others preserving the given order. If the sequence (k_1, k_2, \dots, k_n) is periodic, then we adopt the nota-

Fig. 2 **a** an hv -convex switching S where the two components S_0 and S_1 have different colors. **b** and **c** are two dual sets with respect to S such that **(b)** is not hv -convex while **(c)** is (Color figure online)

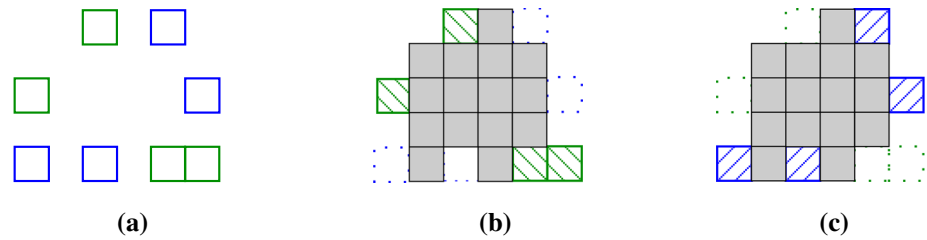
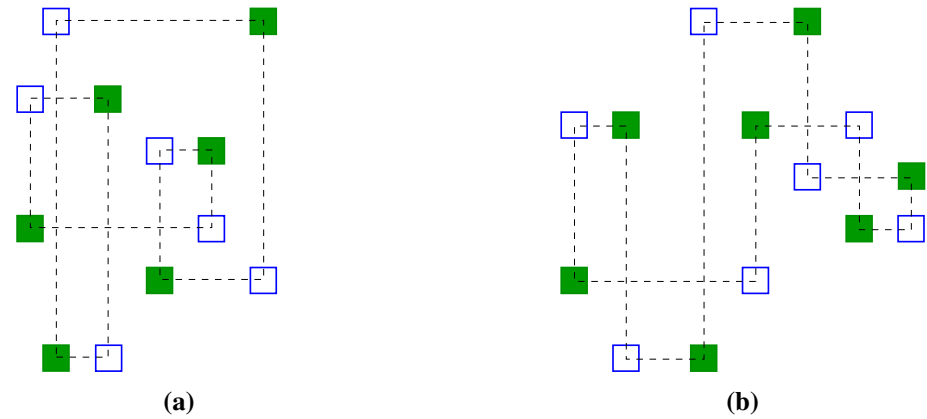


Fig. 3 Examples of squared spirals: **a** a window, and **b** a curl. In each spiral, the vertices belonging to the same component of the hv -switching have the same color and the same pattern (full colored and contour colored) (Color figure online)



tion $(k_1, \dots, k_{n'})_h$, to represent the h time repetition of the sequence $(k_1, \dots, k_{n'})$, with $n = n' \cdot h$; if $h = 1$, we choose to omit it (see Fig. 5 for examples).

It is worth noticing that a $(k_1, \dots, k_n)_h$ hv -sequence encoding a squared spiral, in general does not guarantee its hv -convexity, since the conditions in Lemma 3 do not automatically hold. Figure 4-(a) shows a (12) hv -convex windows, while, in Fig. 4b it is represented a $(7)_2$ hv -convex curl.

Note that the hv -switchings (a) (resp. (b)) in Figs. 3 and 4 are windows and curls of the same type, and have the same number of points, but only in Fig. 4 the hv -convex property holds.

Our aim is to investigate the structure of the hv -sequences that are associated to hv -convex switchings, shortly the hv -convex sequences. In particular, we are concerned with the following problem.

Problem. Characterize, and explicitly reconstruct, the hv -convex sequence associated to any given window or curl.

Remark 6 The notions of hv -convex windows and hv -convex curls already appeared in the literature. For instance, these have been considered in [21] in terms of regular and of irregular switching components. Also, the staircases studied in [1], and the staircase-like structures in [2], are special cases of windows.

A straightforward property of hv -convex switchings is:

Proposition 7 Each horizontal and vertical line intersecting an hv -convex switching $S = (S^0, S^1)$ contains precisely two points of S , one in S^0 and one in S^1 .

The easiest case concerns the hv -convex sequences associated to windows. Let us recall the following result from [13]:

Theorem 8 Let W be a window of size $n \geq 1$ and let $\{w_1, w_2, \dots, w_{4n}\}$ be the set of its vertices as encountered when traveling W from any fixed vertex, and with a given orientation. Then, W is an hv -convex switching if and only if a point $x \in \mathbb{R}^2$ exists such that $w_i \in Z_0(x) \cup Z_2(x)$ for all the odd indices, and $w_i \in Z_1(x) \cup Z_3(x)$ for all the even indices.

In [13, Theorem 3], a simple characterization had been given for what concerns the hv -convex sequences associated to hv -convex windows. For the convenience of the reader, we repeat below the proof.

Theorem 9 For each $n > 0$, an hv -convex switching S has hv -convex sequence $\pi = (4n)$ if and only if S is an hv -convex window.

Proof Suppose that S is a window. Then, there exists a point x as in Theorem 8. Note that all quadrants $Z_i(x)$, $i \in \{0, 1, 2, 3\}$ contain the same number of points of S . Since S has $4n$ vertices, and can be traveled by turning always in the same direction, then $\pi = (4n)$. Conversely, assume that an hv -convex sequence is $\pi = (4n)$. Since π consists of a single entry, then no change of direction occurs when traveling S , so S is a window, and, from the knowledge that π is an hv -convex sequence, we get that S is an hv -convex window. \square

Breaking the constraint of the existence of the point $x \in \mathbb{R}^2$ provides an easy way to define windows of whatever

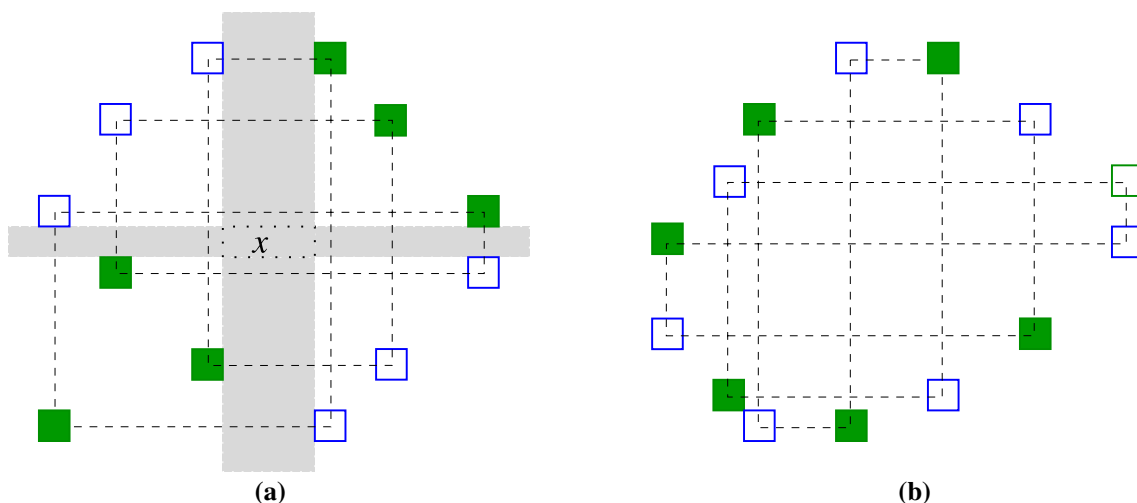


Fig. 4 **a** an *hv*-convex window. The central area, where the x points of Theorem 8 can be selected, and the four quadrants where only elements of one single component of the switching are present, are delimited by the gray shading. **b** a $(7)_2$ *hv*-convex curl is represented (Color figure online)

number (multiple of 4) of vertices that are *hv*-switchings without being *hv*-convex. In Figs. 3a and 4a, there are two examples of windows with the same number of vertices that are non-convex and convex *hv*-switchings.

Differently, the characterization of the *hv*-sequences associated to *hv*-convex switchings represented by curls, shortly *hv*-convex curls, represents a quite challenging task. We address this problem in the following section.

3 Structure of *hv*-convex Sequences Associated to Curls

We move from the following result in [12,13]:

Theorem 10 *Let C be an *hv*-convex curl and let v and w be two of its vertices with the same orientation. If precisely $2n > 0$ consecutive vertices between v and w exist having their opposite orientation, then C is not an *hv*-convex curl.*

Theorem 10 implies that the elements of each *hv*-convex sequence related to an *hv*-convex curl are odd numbers. For instance, Fig. 4, (a), provides an example where the only entry 7 appears, repeated twice. Differently, in Fig. 5, examples are shown of *hv*-convex sequences with two different entries.

3.1 Relating *hv*-convex Curls to Dual Polyominoes

As a first step, let us show how an *hv*-convex curl can be associated to a discrete set in order to produce two dual *hv*-convex polyominoes. It is well known that an *hv*-convex polyomino has the general form as those depicted in Fig. 5. The points of an *hv*-convex polyomino that touch the sides of its minimal

bounding rectangle are called *N(orth)*, *S(outh)*, *E(ast)* and *W(est)* feet, according to their positions.

We denote by *XY*-side the boundary of an *hv*-convex polyomino running from the *X* foot to the *Y* foot, with $X \in \{N, S\}$, and $Y \in \{E, W\}$. Also, $P(S)$ represents an arbitrary *hv*-convex polyomino associated with an *hv*-convex switching S .

We observe that the two components of an *hv*-convex switching that are associated to two dual *hv*-convex polyominoes, must lie on the border of the two polyominoes, in order to preserve horizontal and vertical convexity for both of them, as shown in Fig. 5.

Lemma 11 *Let π be an *hv*-convex sequence, S be the corresponding *hv*-convex switching, and $P(S)$ be an arbitrary *hv*-convex polyomino associated to S . Then, π depends only on the number of points of S lying on two consecutive sides of $P(S)$.*

Proof The *hv*-convex sequence π encodes the alternating lists of clockwise and counterclockwise oriented vertices of S . Since S is a squared spiral, then its vertices are the endpoints of horizontal and vertical sides, and, being S a switching component of $P(S)$, then its vertices all belong to the four sides of $P(S)$. Since the vertices of S correspond pairwise along horizontal and vertical lines, all of them are completely determined just by the subset of vertices belonging to two consecutive sides of $P(S)$. \square

Corollary 12 *Let π be an *hv*-convex sequence associated to an *hv*-convex switching S . Let a and b be the number of points of S lying, respectively, on the *NE*-side, and on the *SE*-side of $P(S)$. If $\pi = (k_1, \dots, k_n)_2$, then $k_1 + \dots + k_n = a + b$.*

Proof Since $\pi = (k_1, \dots, k_n)_2$, then S has $2(k_1 + \dots + k_n)$ vertices. By Lemma 11, the number of vertices of S is twice

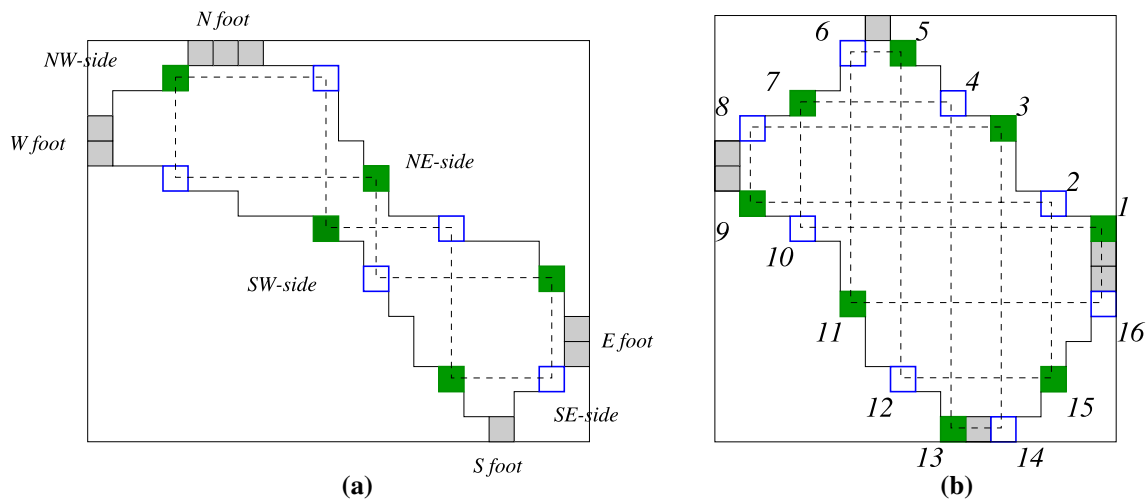


Fig. 5 **a** An hv -convex polyomino where the four feet are highlighted. A $(3, 1)_2$ hv -convex curl is shown (dashed line), and it is associated to two dual hv -convex polyominoes when adding either the full colored

green points or the contour colored blue points. **b** The polyomino is associated to a $(5, 3)_2$ hv -convex curl (dashed line). The labeling of the points of the curl is also provided (Color figure online)

the number of vertices lying on the union of the NE -side and the SE -side of $P(S)$. Therefore, we get $k_1 + \dots + k_n = a + b$. \square

Figure 5 also shows two examples of how hv -convex sequences relate to the points on the sides of a corresponding hv -convex polyomino. In Fig. 5, (a), the hv -convex sequence $\pi = (3, 1)_2$ is considered. According to Corollary 12, we must have $3 + 1 + 1 = 5$ points on two consecutive sides. Up to rotations, this allows four points lying on the NE -side, and one point on the SE -side. By symmetry, we also have four points on the NW -side and one point on the SW -side. In Fig. 5, (b), the sequence $(5, 3)_2$ allows five points to lie both on the NE -side and on the SW -side, and three points to lie both on the NW -side and on the SE -side.

Remark 13 A curl C can be uniquely associated to a couple of integers a and b that count the number of points that lie, respectively, on the NE -side, and on the SE -side of $P(C)$. Therefore, we can also adopt for C the notation $Curl(a, b)$.

We aim in finding an explicit way for the reconstruction of the whole hv -convex sequence related to $Curl(a, b)$ from the only knowledge of the integers a and b . To this, we introduce in the following section two important operators, acting on the labeled vertices of $Curl(a, b)$.

3.2 The Horizontal and the Vertical Shift Operators

From now on, let us label the n points of an hv -convex curl C by $1, 2, \dots, n$, starting from the lowermost point in the NE -side of $P(C)$, and moving counterclockwise, as depicted in Fig. 5b.

Furthermore, let a (resp. b) be the number of points of C lying on the NE -side and SW -side (resp. NW -side and SE -side) of $P(C)$, and assume w.l.g. (by reflection along the horizontal axis) $a \geq b$. We identify each point $x \in C$ with its label $(\text{mod } (2a + 2b))$, and we define its horizontal and vertical shifted version $H(x)$ and $V(x)$, respectively, as follows:

$$\begin{aligned} H(x) &= 2a + 1 - x \pmod{(2a + 2b)}; \\ V(x) &= 1 - x \pmod{(2a + 2b)}. \end{aligned}$$

We recall that in case $a = b$, the corresponding switching component is a window whose sequence is of the form $(4h)$.

Lemma 14 Let $a, b \in \mathbb{N}$, with $a > b$. For each point $x \in C = Curl(a, b)$, the point of C with label $H(x)$ lies in the same row of x , while the point of C with label $V(x)$ lies in the same column of x .

Proof According to the positions of x on the sides of $P(C)$, four cases arise:

(i) x lies on the NE -side of $P(C)$: the label of the point y sharing the same row can be obtained by adding to the uppermost point of the NE -side, whose label is a , the $a - x$ points needed to reach again the row of x , so the label of y is $2a + 1 - x = H(x)$. On the other hand, the label of the point z sharing the same column with x is obtained by moving back of $x - 1$ from the last label $2a + 2b$, so $1 - x \pmod{(2a + 2b)}$ is the label of $V(x)$.

(ii) x lies on the NW -side of $P(C)$: the same argument as in case i) provides the label of y sharing the same row is again $2a + 1 - x = H(x)$. The label of the point z sharing

the same column with x is obtained by adding to the label $a + b + 1$ the number of points from $a + b$ to x , i.e., $a + b - x$. So the label of z turns out to be $a + b + 1 + a + b - x = 1 - x \pmod{(2a + 2b)} = V(x)$.

Cases (iii) and (iv), i.e., when x lies either on the SW -side or on the SE -side of $P(C)$, can be treated analogously. The reader can check these remaining cases on Fig.5, (b). \square

Lemma 15 *Let π be an hv -convex sequence corresponding to an hv -convex curl C . For every odd $\alpha = 2h + 1 \in \pi$, there exists a consecutive sequence x_1, \dots, x_α of points of C traveled with the same orientation such that the label of x_α is $2hb + x_1 \pmod{(2a + 2b)}$ if x_1 and x_2 lie in the same row, while if they lie in the same column, the label is $2ha + x_1 \pmod{(2a + 2b)}$.*

Proof Let us first assume that the path from x_1 to x_2 is horizontal. The points x_1, \dots, x_α have labels

$$x_1, H(x_1), VH(x_1), HVH(x_1), (VH)^2(x_1), \dots, (VH)^h(x_1).$$

By Lemma 14, it holds

$$\begin{aligned} VH(x_1) &= V(2a + 1 - x_1) \\ &\pmod{(2a + 2b)} = 1 - 2a - 1 + x_1 \\ &\pmod{(2a + 2b)} = 2b + x_1 \pmod{(2a + 2b)}. \end{aligned}$$

So, the label of the point x_α turns out to be

$$\begin{aligned} (VH)^h(x_1) &= (VH)^{h-1}(2b + x_1) \\ &\pmod{(2a + 2b)} = \dots = 2hb + x_1 \\ &\pmod{(2a + 2b)}. \end{aligned}$$

On the other hand, if the path from x_1 to x_2 is vertical, the labels of the points x_1, \dots, x_α are

$$x_1, V(x_1), HV(x_1), VHV(x_1), (HV)^2(x_1), \dots, (HV)^h(x_1).$$

Again by Lemma 14, it holds

$$\begin{aligned} HV(x_1) &= H(1 - x_1) \\ &\pmod{(2a + 2b)} = 1 + 2a - 1 + x_1 \\ &\pmod{(2a + 2b)} = 2a + x_1 \\ &\pmod{(2a + 2b)}. \end{aligned}$$

So, the label of the point x_α turns out to be

$$\begin{aligned} (HV)^h(x_1) &= (HV)^{h-1}(2a + x_1) \\ &\pmod{(2a + 2b)} = \dots = 2ha + x_1 \pmod{(2a + 2b)}. \end{aligned}$$

\square

The next theorem provides a complete characterization of the combinatorial structure of $Curl(a, b)$ for any possible choice of a and b , also showing that the corresponding hv -convex sequences always consist of one or two different repeated entries.

Remark 16 For the sake of brevity, in what follows sentences of the form “the point x belongs to the XY -side of $P(C)$ ” are often simply replaced by $x \in XY$.

Theorem 17 *Let $a, b \in \mathbb{N}$, with $b < a$, and let π be the hv -convex sequence associated to $C = Curl(a, b)$. Then, π consists of N_α entries $\alpha = 2h + 1$, and N_γ entries $\gamma = 2h + 3$, where*

$$N_\alpha = 2(h + 1)a - 2(h + 2)b \tag{3}$$

$$N_\gamma = 2(h + 1)b - 2ha \tag{4}$$

$$h = \left\lfloor \frac{b}{a - b} \right\rfloor \tag{5}$$

Proof Let $x \in [1, b]$, and note that $V(x) = 1 - x = 2a + 2b + 1 - x \pmod{(2a + 2b)}$, so $V(x) \in SE$ (see Remark 16). Also, $H(x) = 1 + 2a - x$, so, if $x \leq a - b$, then $H(x) \in SW$. This means that each $x \in [1, a - b]$ has orientation different from $H(x)$, and has the same orientation of $V(x)$, say clockwise without loss of generality. Therefore, each such x is the endpoint of a path in $Curl(a, b)$. Note that $(HV)^h(x) \in NE$ if h is even, while $(HV)^h(x) \in SW$ if h is odd. Consequently, x is the starting point of a clockwise oriented path of minimal length $\alpha = 2h + 1$ if and only if the following conditions hold

- (i) $V((HV)^h(x)) \in SW$ if $h = 2k$, for some $k \geq 0$.
- (ii) $V((HV)^h(x)) \in NE$ if $h = 2k - 1$, for some $k \geq 1$.

Suppose that $h = 2k$ for some $k \geq 0$. We have

$$\begin{aligned} (HV)^h(x) &= 2ha + x \\ &\pmod{(2a + 2b)} = 4ka + x \\ &\pmod{(2a + 2b)} = 2k(a - b) + x, \end{aligned}$$

so $(HV)^h(x) \in NE$ provides $1 \leq 2k(a - b) + x \leq a$, that is

$$x \in [hb - ha + 1, hb - (h - 1)a]. \tag{6}$$

Condition (i) provides

$$\begin{aligned} a + b + 1 &\leq 1 - (2k(a - b) + x) \\ &\pmod{(2a + 2b)} \leq 2a + b, \end{aligned}$$

that is

$$\begin{cases} x \geq 1 - 2k(a - b) - 2a - b \pmod{(2a + 2b)} \\ \quad = 1 + (2k + 1)b - 2ka > (2k + 1)b - 2ka \\ x \leq -2k(a - b) - a - b \pmod{(2a + 2b)} \\ \quad = -2k(a - b) + a + b = (2k + 1)b - (2k - 1)a, \end{cases}$$

so that x ranges in the interval $((2k + 1)b - 2ka, (2k + 1)b - (2k - 1)a) = ((h + 1)b - ha, hb - ha + a + b]$. Matching with (6), and with $x \in [1, a - b]$, we obtain

$$\begin{cases} (h + 1)b - ha \geq 0 \\ (h + 1)b - ha < a - b \\ hb - ha + a + b \leq a - b, \end{cases}$$

which provides

$$\frac{h}{h + 1}a \leq b < \frac{h + 1}{h + 2}a. \tag{7}$$

This implies

$$\frac{b}{a - b} - 1 < h \leq \frac{b}{a - b},$$

and (5) follows.

Therefore, if $h = 2k$, and $x \in ((h + 1)b - ha, (h + 1)b - (h - 1)a]$, there exist $N_{\alpha_1} = (h + 1)a - (h + 2)b$ entries of π equal to $\alpha = 2h + 1$. Differently, if $h = 2k$, and $x \in [1, (h + 1)b - ha]$, any clockwise oriented path starting at x , must have length $2h' + 1 \neq \alpha$, and consequently $h' > h$, due to the minimal assumption on α . Let us consider $h' = h + 1$. Then, h' is odd, so $HV^{h'}(x) \in SW$. Since h is even, it is $-ha = hb \pmod{(2a + 2b)}$, so, also exploiting (7), we get

$$\begin{aligned} V((HV)^{h'}(x)) &= 1 - (2(h + 1)a + x) \\ &= -ha + (h + 2)b + 1 - x \\ &\leq -ha + (h + 2)b < -ha + (h + 1)a = a. \end{aligned}$$

Therefore, $V((HV)^{h'}(x)) \in NE$, which implies that the orientation of $V((HV)^{h'}(x))$ differs from that of $(HV)^{h'}(x)$. This means that any $x \in [1, (h + 1)b - ha]$ is the starting point of a path of length $\gamma = 2h' + 1 = 2h + 3$. Therefore, there exist $N_{\gamma_1} = (h + 1)b - ha$ entries of π equal to $\gamma = 2h + 3$.

When $x \in [1, a - b]$, we get $N_{\alpha_1} + N_{\gamma_1}$ entries of π , corresponding to disjoint paths of length α and γ , respectively. These paths include a number of points of $Curl(a, b)$ given by

$$\begin{aligned} &\alpha N_{\alpha_1} + \gamma N_{\gamma_1} \\ &= (2h + 1)[(h + 1)a - (h + 2)b] \\ &\quad + (2h + 3)[(h + 1)b - ha] = a + b. \end{aligned}$$

By symmetry, the same arguments can be applied to paths having starting points at $x \in [a + b + 1, 2a]$, which provides $N_{\alpha_2} = N_{\alpha_1}$ paths of length α , for $x \in (a + b + (h + 1)b - ha, 2a]$, and $N_{\gamma_2} = N_{\gamma_1}$ paths of length γ , for $x \in [a + b + 1, a + b + (h + 1)b - ha]$. These paths include $a + b$ additional points of $Curl(a, b)$. Consequently, we have $N_{\alpha} = N_{\alpha_1} + N_{\alpha_2} = 2[(h + 1)a - (h + 2)b]$ entries of π equal to $\alpha = 2h + 1$, and $N_{\gamma} = N_{\gamma_1} + N_{\gamma_2} = 2[(h + 1)b - ha]$ entries of π equal to $\gamma = 2h + 3$. These entries correspond to paths of $Curl(a, b)$ which include $2a + 2b$ different points. Since $Curl(a, b)$ contains exactly $2a + 2b$ points, then the sequence π cannot have any further entry, and the statement follows.

Suppose now that $x \in [1, a - b]$, and $h = 2k - 1$ for some $k \geq 1$. Then, we can write

$$\begin{aligned} (HV)^h(x) &= 2ha + x \pmod{(2a + 2b)} \\ &= 4ka - 2a + x \pmod{(2a + 2b)} \\ &= 2k(a - b) + 2b + x, \end{aligned}$$

so $(HV)^h(x) \in SW$ provides $1 + a + b \leq 2k(a - b) + 2b + x \leq 2a + b$, that is

$$\begin{aligned} x &\in [(2k - 1)b - (2k - 1)a + 1, (2k - 1)b - 2(k - 1)a] \\ &= [hb - ha + 1, hb - ha + a] \end{aligned} \tag{8}$$

Condition (ii) provides

$$1 \leq 1 - (2k(a - b) + 2b + x) \pmod{(2a + 2b)} \leq a,$$

that is

$$\begin{cases} x \geq 1 - 2k(a - b) - 2b - a \pmod{(2a + 2b)} \\ \quad = 1 - 2ka + 2kb + a > 2kb - (2k - 1)a \\ x \leq -2k(a - b) - 2b \pmod{(2a + 2b)} \\ \quad = 2kb - 2ka + 2a, \end{cases}$$

so that x ranges in the interval $(hb - ha + b, hb - ha + a + b]$. Matching with (8), and with $x \in [1, a - b]$, we get $x \in (hb - ha + b, a - b]$, under conditions (5). This provides again $N_{\alpha_1} = (h + 1)a - (h + 2)b$ as in the case when h is even.

Therefore, if $h = 2k - 1$, and $x \in (hb - ha + b, a - b]$, then there exist $N_{\alpha_1} = (h + 1)a - (h + 2)b$ entries of π equal to $\alpha = 2h + 1$.

Differently, if $h = 2k - 1$, and $x \in [1, hb - ha + b]$, any clockwise oriented path starting at x , must have length $2h' + 1 \neq \alpha$, and consequently $h' > h$, due to the minimal assumption on α . Let us consider $h' = h + 1$. Then, h' is odd, so $(HV)^{h'}(x) \in NE$. Being $h' = h + 1$ even, we can write

$$\begin{aligned} V((HV)^{h'}(x)) &= 1 - (2(h + 1)a + x) \\ \pmod{(2a + 2b)} &= (h + 2)b - ha + a + b + 1 - x. \end{aligned}$$

Since $x \in [1, hb - ha + b]$, then by (5), we get

$$\begin{aligned} V((HV)^{h'}(x)) &\geq a + 2b + 1 > a + b + 1 \\ V((HV)^{h'}(x)) &< a + a + b + 1 - x \leq 2a + b. \end{aligned}$$

Therefore, $V((HV)^{h'}(x)) \in SW$, which implies that the orientation of $V((HV)^{h'}(x))$ differs from that of $(HV)^{h'}(x)$. This means that any $x \in [1, (h + 1)b - ha]$, is the starting point of a path of length $\gamma = 2h' + 1 = 2h + 3$. Therefore, there exist

$$N_{\gamma_1} = (h + 1)b - ha$$

entries of π equal to $\gamma = 2h + 3$.

Consequently, for any $x \in [1, a - b]$, we have the same results as in the case h even. By symmetry, the same holds also for $x \in [a + b + 1, 2a]$, and the statement follows. \square

The following corollary characterizes the hv -sequences having just one odd repeated entry, so it represents a kind of curl counterpart of the case of windows according to Theorem 9.

Corollary 18 For any $a, b \in \mathbb{N}$, $a > b$, the entries of the hv -convex sequence π associated to $Curl(a, b)$ are all equal if and only if $\frac{b}{a-b}$ is an integer. In addition, the $2(a - b)$ entries in this case are all equal to $\frac{a+b}{a-b}$.

Proof If $\frac{b}{a-b}$ is an integer, then, by (5), $h = \frac{b}{a-b}$, so (4) provides $N_{\gamma} = 0$, and, by (3), $N_{\alpha} = 2ha + 2a - 2hb - 4b = 2(a - b)$, which proves the statement.

Conversely, suppose that π consists of just one repeated odd entry. Then, by Theorem 17, one of N_{α} or N_{γ} must be equal to 0. Note that $N_{\alpha} = 0$ implies $h = \frac{b}{a-b} - 1$, which contradicts (5). Therefore, it must be $N_{\alpha} \neq 0$, so $N_{\gamma} = 0$, and consequently $h = \frac{b}{a-b}$. This implies that $\frac{b}{a-b}$ is an integer, and that π has $N_{\alpha} = 2(h+1)a - 2(h+2)b = 2a - 2b$ entries equal to $\alpha = 2h + 1 = \frac{a+b}{a-b}$. \square

Remark 19 By Theorem 17 we get the simple procedure *CompSeqEntries*, that computes the values of α , γ , N_{α} and N_{γ} from arbitrarily chosen integers a and b , with $a > b$, namely

CompSeqEntries

- set $h = \lfloor \frac{b}{a-b} \rfloor$
- return $\alpha = 2h + 1$, $\gamma = 2h + 3$, $N_{\alpha} = 2(h + 1)a - 2(h + 2)b$, $N_{\gamma} = 2(h + 1)b - 2ha$

Example 20 Assuming $a = 28, b = 5$, the procedure *CompSeqEntries* in Remark 19, we have $h = 0, \alpha = 1, \gamma = 3, N_1 = 2a - 4b = 36$ and $N_3 = 2b = 10$. The related hv -convex sequence is $\pi = (3, 1^4, 3, 1^4, 3, 1^3, 3, 1^4, 3, 1^3)_2$ where 1^j means that the entry 1 is repeated j times. In Fig. 6, a part of the corresponding hv -convex switching C is shown.

Remark 21 The case $a > 2b$ is the only one that allows entries equal to 1 to appear in the hv -convex sequence.

Example 22 Assuming $a = 27, b = 19$, the procedure *CompSeqEntries* in Remark 19 finds $\frac{2}{3}a < b < \frac{3}{4}a$, so $h = 2$, and returns $\alpha = 5, \gamma = 7, N_5 = 6a - 8b = 10, N_7 = 6b - 4a = 6$.

4 The Structure of hv -convex Switchings

Theorem 17 shows that any hv -convex sequence consists of one or two repeated entries. Corollary 18 characterizes the case when the hv -convex sequence consists of just one repeated entry. Differently, the problem of understanding how the entries α and γ mutually alternate in π becomes relevant. To this, we need to know the labels of the endpoints of the paths having lengths α and γ .

Remark 23 In what follows, all the involved intervals consist of just integer labels, even if the simplified “continuous” notation $(a, b]$, (instead of $(a, b] \cap \mathbb{N}$) is adopted.

To address the above problem, we first apply *CompSeqEntries* as pointed out in Remark 19, in order to prove the following.

Lemma 24 Let $N_{\gamma} = 2(h + 1)b - 2ha$ be computed by *CompSeqEntries*. According to the parity of h , it holds

$$\begin{aligned} -h \geq 0 \text{ even: } &\frac{N_{\gamma}}{2} \pmod{(2a + 2b)} = -2ha - 2a - b \pmod{(2a + 2b)}; \\ -h \text{ odd: } &\frac{N_{\gamma}}{2} \pmod{(2a + 2b)} = -2ha - a \pmod{(2a + 2b)}. \end{aligned}$$

Proof $h \geq 0$ even: it holds

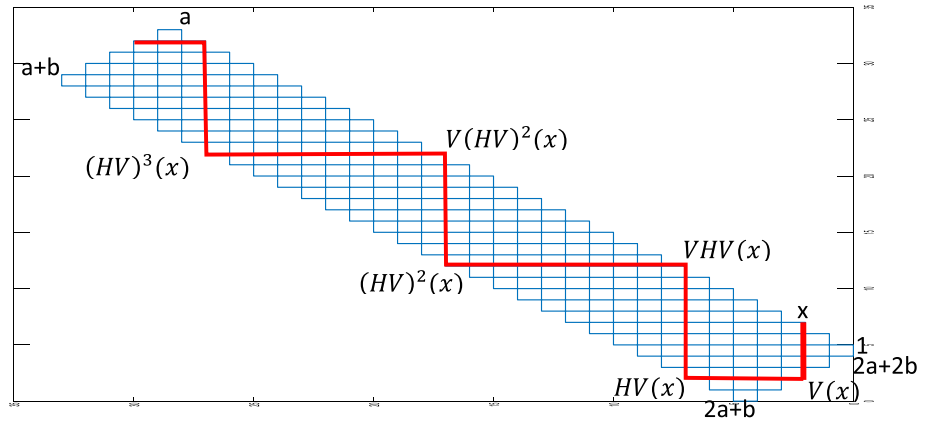
$$\begin{aligned} &-2ha - 2a - b \pmod{(2a + 2b)} \\ &= -2ha - 2a - b + (h + 2)(a + b) \\ &\pmod{(2a + 2b)} = (h + 1)b - ha = \frac{N_{\gamma}}{2}. \end{aligned}$$

h odd: it holds

$$\begin{aligned} &-2ha - a \pmod{(2a + 2b)} \\ &= -(h + 1)a - (h - 1)a - a \\ &\pmod{(2a + 2b)} = (h + 1)b - (h - 1)a - a \\ &\pmod{(2a + 2b)} = \frac{N_{\gamma}}{2}. \quad \square \end{aligned}$$

Theorem 25 Let us consider two integers a and b such that $a > b$. According to what defined in Theorem 17, let C and $P(C)$ be the hv -convex curl and the hv -convex polymino associated to a and b , and compute the values of $h, \alpha, \gamma, N_{\alpha}$ and N_{γ} . We define the following sets of (integer) labels:

Fig. 6 A hv -convex switching C with $a = 28$ and $b = 5$



$$I_\alpha: \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right] \cup \left(b, \frac{N_\alpha}{2} + b\right] \cup \left(a + b + \frac{N_\gamma}{2}, a + b + \frac{N_\alpha + N_\gamma}{2}\right] \cup \left(a + 2b, \frac{N_\alpha}{2} + a + 2b\right];$$

$$I_\gamma: \left[1, \frac{N_\gamma}{2}\right] \cup \left[a + 1 - \frac{N_\gamma}{2}, a\right] \cup \left[a + b + 1, a + b + \frac{N_\gamma}{2}\right] \cup \left[2a + b + 1 - \frac{N_\gamma}{2}, 2a + b\right].$$

It holds that the elements of I_α (resp. I_γ) are exactly (the labels of) the endpoints of the same-oriented paths of length α (resp. γ) of the hv -convex curl C .

Proof Let us assume w.l.g. that the point with label x lies in the NE -side of $P(C)$ and it is the first point of a path traveled clockwise. From this assumption, it immediately follows that $H(x) \in SW$, with $H(x) \leq H(1)$, i.e., $a + b + 1 \leq 1 + 2a - x \leq 2a$, and, consequently, $x \leq a - b = \frac{N_\alpha + N_\gamma}{2}$. So, such a path has a vertical first side. Lemma 15 assures that the last point of the path is $(HV)^h(x) = 2ha + x \pmod{(2a + 2b)}$ if and only if its length is $\alpha = 2h + 1$. The following two cases arise:

h even: $(HV)^h(x) \in NE$ and $V(HV)^h(x) \in SW$, between $V(a)$ and $2a + b$. It follows the inequality

$$2a + 2b - a + 1 \pmod{(2a + 2b)} \leq 1 - 2ha - x \pmod{(2a + 2b)} \leq 2a + b \pmod{(2a + 2b)},$$

and by Lemma 24

$$\begin{cases} x \leq -2ha - a - 2b \pmod{(2a + 2b)} = \frac{N_\gamma}{2} + a - b \\ = \frac{N_\gamma + N_\alpha}{2} \\ x \geq 1 - 2ha - 2a - b \pmod{(2a + 2b)} = 1 + \frac{N_\gamma}{2}. \end{cases}$$

As a consequence, the previous condition $x \in \left[1, \frac{N_\alpha + N_\gamma}{2}\right]$ restricts to $x \in \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right]$.

h odd: $(HV)^h(x) \in SW$, and $V(HV)^h(x) \in NE$, between a and $V(2a + b)$. It follows the inequality

$$1 - 2a - b \pmod{(2a + 2b)} \leq 1 - 2ha - x \pmod{(2a + 2b)} \leq a \pmod{(2a + 2b)},$$

and furthermore, $1 + b \leq 1 - 2ha - x \pmod{(2a + 2b)} \leq a$ and by Lemma 24

$$\begin{cases} x \leq -2ha - b \pmod{(2a + 2b)} = 2hb - b \pmod{(2a + 2b)} \\ x \geq 1 - 2ha - a \pmod{(2a + 2b)} = 1 + \frac{N_\gamma}{2}. \end{cases}$$

Since it also holds $h \leq 1$ and $2hb - b \geq b$, then the first inequality is always verified. Again, by the second inequality, the condition $x \in \left[1, \frac{N_\alpha + N_\gamma}{2}\right]$ restricts to $x \in \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right]$.

Summing up, if a path of length α has starting point x in the NE -side of $P(C)$, then it holds $x \in \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right]$.

Let us now focus on the final point $(HV)^h(x)$ of the path: it lies on the NE -side of $P(C)$ when h is even, and it lies on the SW -side of $P(C)$, otherwise. Again, let us analyze the two cases more precisely:

h even: by Lemma 24 it holds $2ha \pmod{(2a + 2b)} = -\frac{N_\gamma}{2} + b$, so $x \in \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right]$ implies

$$(HV)^h(x) \in \left(b, \frac{N_\alpha}{2} + b\right].$$

h odd: a similar computation as in the previous case, from $x \in \left(\frac{N_\gamma}{2}, \frac{N_\alpha + N_\gamma}{2}\right]$ leads to

$$(HV)^h(x) \in \left(a + 2b, \frac{N_\alpha}{2} + a + 2b\right].$$

In case the path is traveled counterclockwise and/or the starting point of the path lies on the *SW*-side of $P(C)$, then the previous results continue holding, up to a rearrangement of the starting and the ending points. These computations show that the point x can also belong to the interval $\left(a + b + \frac{N_\gamma}{2}, a + b + \frac{N_\alpha + N_\gamma}{2}\right]$, while the ending point $(HV)^h(x)$ can belong either to $\left(a + 2b, a + 2b + \frac{N_\alpha}{2}\right]$ or to $\left(b, b + \frac{N_\alpha}{2}\right]$ according to the parity of h . So, the extremal points of the paths of length α correspond to the set I_α .

Concerning the set I_γ , we suppose that the final point of a path starting in x is of the form $(HV)^{h+1}(x) = 2(h+1)a - x \pmod{(2a + 2b)}$, i.e., its length is $\gamma = 2h + 3$.

The same reasoning as for the case of the α length paths can be used after updating h with $h + 1$ and up to exchanging the *NW*-side and the *SE*-side of $P(C)$. More precisely:

h even: $(HV)^{h+1}(x)$ lies on the *SW*-side of $P(C)$ and $V(HV)^{h+1}(x)$ lies on the *NE*-side of $P(C)$ between and a . It follows the inequality

$$1 - 2a - b \pmod{(2a + 2b)} \leq 1 - 2(h + 1)a - x \pmod{(2a + 2b)} \leq a \pmod{(2a + 2b)},$$

and by Lemma 24

$$\begin{cases} x \leq -2ha + b \pmod{(2a + 2b)} = \frac{N_\gamma}{2} \\ x \geq 1 - 2ha - 3a \pmod{(2a + 2b)} = 1 + \frac{N_\gamma}{2} \\ \quad + b - a = 1 - \frac{N_\alpha}{2}. \end{cases}$$

Since $1 - \frac{N_\alpha}{2} < 0$ the second inequality is always verified, so the first inequality allows the condition $x \in \left[1, \frac{N_\alpha + N_\gamma}{2}\right]$ to restrict to $x \in \left[1, \frac{N_\gamma}{2}\right]$.

h odd: $(HV)^{h+1}(x) \in NE$, and $V(HV)^{h+1}(x) \in SW$, between $V(a)$ and $2a + b$. It follows the inequality

$$\begin{aligned} 1 - a \pmod{(2a + 2b)} \\ \leq 1 - 2(h + 1)a - x \pmod{(2a + 2b)} \leq 2a + b \pmod{(2a + 2b)}, \end{aligned}$$

and, by Lemma 24,

$$\begin{cases} x \leq -2ha - a \pmod{(2a + 2b)} = \frac{N_\gamma}{2} \\ x \geq -2ha - 4a - b \pmod{(2a + 2b)} = \frac{N_\gamma}{2} \\ \quad - 3a + b \pmod{(2a + 2b)} = \frac{N_\gamma}{2} - a + b. \end{cases}$$

Since $\frac{N_\gamma}{2} - a + b = -\frac{N_\alpha}{2} < 0$, then the second inequality is always verified. Again, by the first inequality, the condition $x \in \left[1, \frac{N_\alpha + N_\gamma}{2}\right]$ restricts to $x \in \left[1, \frac{N_\gamma}{2}\right]$.

Summing up, if a γ length path starts from a point x in the *NE*-side of $P(C)$, then $x \in \left[1, \frac{N_\gamma}{2}\right]$. The last point of the path $(HV)^{h+1}(x)$ lies on the *SW*-side of $P(C)$ if h is even, and on the *NE*-side of $P(C)$ otherwise:

h even: by Lemma 24 it holds $2ha \pmod{(2a + 2b)} = -\frac{N_\gamma}{2} + b$, so $x \in \left[1, \frac{N_\gamma}{2}\right]$ implies

$$(HV)^{h+1}(x) \in \left[2a + b + 1 - \frac{N_\gamma}{2}, 2a + b\right].$$

h odd: a similar computation leads to

$$(HV)^{h+1}(x) \in \left[a + 1 - \frac{N_\gamma}{2}, a\right].$$

In case the path is traveled counterclockwise and/or the starting point of the path lies on the *SW*-side of $P(C)$, then the previous results continue holding, up to a rearrangement of the starting and the ending points. These computations show that the point x can also belong to the interval $\left[a + b + 1, a + b + \frac{N_\gamma}{2}\right]$, while the ending point $(HV)^{h+1}(x)$ can belong either to $\left[a + 1 - \frac{N_\gamma}{2}, a\right]$ or to $\left[2a + b - \frac{N_\gamma}{2}, 2a + b\right]$ according to the parity of h . So, the extremal points of the paths of length γ correspond to the set I_γ . \square

Denote by π_j the subsequence of π consisting of its first j -entries.

The following theorem provides the answer to how the entries α and γ mutually alternate in the *hv*-convex sequence S .

Theorem 26 *Let $a, b \in \mathbb{N}$, with $a > b$, and let $d = \gcd(a, b)$. Then, the *hv*-convex sequence π associated to $\text{Curl}(a, b)$ is $\pi = (\pi_{\frac{a-b}{d}})_{2d}$*

Proof Let us consider the function $f_\lambda(x) = (HV)^\lambda(x) = 2\lambda a + x$, where

$$\begin{aligned} \lambda &= \frac{N_\alpha}{2}h + \frac{N_\gamma}{2}(h + 1) + \frac{N_\alpha + N_\gamma}{4} && \text{if } \frac{N_\alpha + N_\gamma}{2} \text{ is even} \\ \lambda &= \frac{N_\alpha}{2}h + \frac{N_\gamma}{2}(h + 1) + \frac{1}{2} \left(\frac{N_\alpha + N_\gamma}{2} - 1\right) && \text{if } \frac{N_\alpha + N_\gamma}{2} \text{ is odd} \end{aligned}$$

Note that $\frac{N_\alpha + N_\gamma}{2} = a - b$, and

$$\begin{aligned} \frac{N_\alpha}{2}h + \frac{N_\gamma}{2}(h + 1) &= \frac{N_\alpha + N_\gamma}{2}h + \frac{N_\gamma}{2} \\ &= (a - b)h + (h + 1)b - ha = b, \end{aligned}$$

so that

$$\begin{aligned} \lambda &= b + \frac{1}{2} \frac{N_\alpha + N_\gamma}{2} && \text{if } a - b \text{ is even} \\ \lambda &= b + \frac{1}{2} \left(\frac{N_\alpha + N_\gamma}{2} - 1 \right) && \text{if } a - b \text{ is odd} \end{aligned}$$

Consequently, we get

$$f_\lambda(x) = 2a \left(b + \frac{1}{2} \frac{N_\alpha + N_\gamma}{2} \right) + x = a^2 + ab + x \tag{9}$$

if $a - b$ is even

$$\begin{aligned} f_\lambda(x) &= 2a \left(b + \frac{1}{2} \left(\frac{N_\alpha + N_\gamma}{2} - 1 \right) \right) + x \\ &= a^2 + ab - a + x && \text{if } a - b \text{ is odd} \end{aligned} \tag{10}$$

Let $j \in \left[1, \frac{N_\alpha + N_\gamma}{2} \right]$, and let x_j be the starting point of the j -th walked path. Then, the label of the starting point of the $(a - b + j)$ -th traveled path is $y_j = f_\lambda(x_j) \pmod{(2a + 2b)}$ if $a - b$ is even, and $y_j = V(f_\lambda(x_j)) = 1 - f_\lambda(1) \pmod{(2a + 2b)}$ if $a - b$ is odd.

Assume $d = 1$, and compute $f_\lambda(x_j)$ in the different allowed cases. We have

- (i) If $a - b$ is even, then, being $d = 1$, both a and b are odd integers, so $a > b$ implies $a > 1$, and consequently $a^2 + ab = (a - 1)(a + b) + a + b = a + b \pmod{(2a + 2b)}$, and we can write

$$\begin{aligned} f_\lambda(x_j) &\pmod{(2a + 2b)} \\ &= a + b + x_j \pmod{(2a + 2b)}, \end{aligned} \tag{11}$$

- (ii) If $a - b$ is odd, we consider two subcases.

- (a) a is even and b is odd. We have

$$\begin{aligned} a^2 + ab - a &= a(a + b) - a \\ &= -a \pmod{(2a + 2b)} = a + 2b. \end{aligned}$$

Therefore, we get

$$\begin{aligned} 1 - f_\lambda(x_j) &\pmod{(2a + 2b)} \\ &= a + 2b + x_j \pmod{(2a + 2b)}. \end{aligned} \tag{12}$$

- (b) a is odd and b is even. We have

$$\begin{aligned} a^2 + ab - a &= (a - 1)(a + b) + b \\ &= b \pmod{(2a + 2b)}. \end{aligned}$$

Consequently, we can write

$$\begin{aligned} 1 - f_\lambda(x_j) &\pmod{(2a + 2b)} \\ &= b + x_j \pmod{(2a + 2b)}. \end{aligned} \tag{13}$$

In all the cases, by Theorem 25 it follows that $x_j \in I_\alpha$ if and only if $y_j \in I_\alpha$, and $x_j \in I_\gamma$ if and only if $y_j \in I_\gamma$. Consequently, the entries of π repeat after the first $\frac{N_\alpha + N_\gamma}{2} = a - b$, which implies that $\pi = (\pi_{a-b})_2$. If $d > 1$, let $a' = \frac{a}{d}$ and $b' = \frac{b}{d}$. Then, with the same arguments, we obtain that the sequence associated to $Curl(a', b')$ is $(\pi_{a'-b'})_2$. Therefore, the sequence π associated to $Curl(a, b)$ is obtained by repeating $(\pi_{a'-b'})_2$ for d times, namely $\pi = (\pi_{a'-b'})_{2d}$, and the statement follows. \square

Theorems 25 and 26 lead to algorithm *CompSeq*, which returns the hv -convex sequence π consisting of entries α and γ , and associated to $Curl(a, b)$, for any pair of integers $a, b, a > b$. The algorithm also exploits the procedure *CompSeqEntries* defined in Remark 19.

Algorithm 1: CompSeq

Input: two integers a and b , with $a > b$.

Output: The hv -convex sequence π associated to a and b (according to Theorem 25 and Theorem 26).

```

compute  $d = \text{gcd}(a, b)$ ;
 $a = \frac{a}{d}, b = \frac{b}{d}$ ;
compute  $h, \alpha, \gamma, I_\alpha$  and  $I_\gamma$  %CompSeqEntries in Remark 19;
create a vector  $\pi$  of length  $\frac{N_\alpha + N_\gamma}{2}$  and set  $x = 1$ ;
for  $i = 1 : \frac{N_\alpha + N_\gamma}{2}$  do
  if  $x \in I_\alpha$  then
     $\pi(i) = \alpha$ ;
     $y = (VH)^h(x)$ ;
  else
     $\pi(i) = \gamma$ ;
     $y = (HV)^{h+1}(x)$ ;
  end if
  if  $i$  even then
     $x = H(y)$ ;
  else
     $x = V(y)$ ;
  end if
end for
return  $(\pi)_{2d}$ . % we stress that  $\pi$  equals  $\pi_{\frac{a-b}{d}}$  of Theorem 26

```

Example 27 Let us show the run of *CompSeq* on input $a = 27$ and $b = 19$ (see Fig. 7). *CompSeqEntries*(27, 19) outputs $h = 2, \alpha = 5, \gamma = 7, N_\alpha = 10$, and $N_\gamma = 6$. Then, we have

$$\begin{aligned} I_\alpha &= (3, 8] \cup (19, 24] \cup (49, 54] \cup (65, 70] \text{ and} \\ I_\gamma &= [1, 3] \cup [25, 27] \cup [47, 49] \cup [71, 73]. \end{aligned}$$

The *for* loop runs from 1 to $a - b = 8$ setting the elements in π as follows:

$i = 1$: initializes $x = 1$ and, since $1 \in I_\gamma$, it sets $\pi(1) = 7$. Then, $y = (HV)^3(1) = 71$ and, since $i = 1$ is odd $x = V(71) = 1 - 71 = -70 \pmod{92} = 22$;

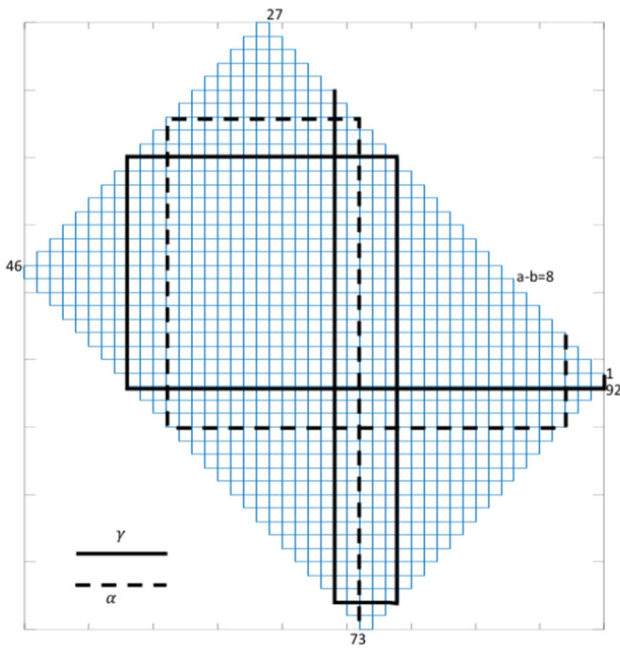
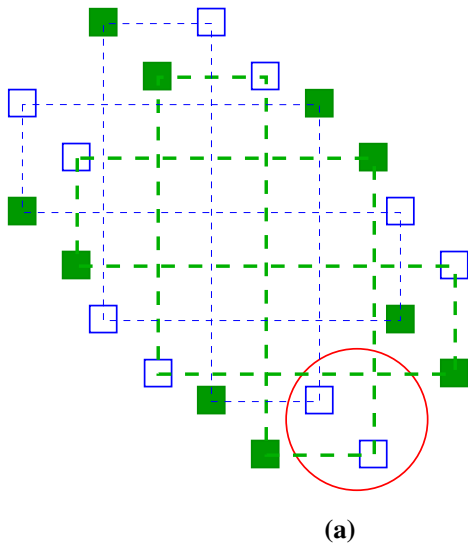


Fig. 7 The Curl(27,19), with a few labeled points and two paths of lengths α (dashed line) and γ (solid line) shown

$i = 2$: since $22 \in I_\alpha$, then $\pi(2) = 5$. Then, $y = (VH)^2(22) = 98 \pmod{92} = 6$ and, since $i = 2$ is even, $x = H(6) = 55 - 6 = 49$;
 $i = 3$: since $49 \in I_\gamma$, then $\pi(3) = 7$. Then, $y = (HV)^3(49) = 70 + 49 \pmod{92} = 27$ and, since $i = 3$ is odd, $x = V(27) = 1 - 27 \pmod{92} = 66$;
 $i = \dots$:



(a)

The *for* loop ends providing the vector $\pi_8 = (7, 5, 7, 5, 5, 7, 5, 5)$, and the final output is $\pi = (\pi_8)_2$.

Remark 28 Let C_1 and C_2 be two *hv*-convex curls having the same *hv*-convex sequence π . Then, C_1 and C_2 can be merged into one single *hv*-convex curl whose *hv*-convex sequence is $(\pi)_2$.

The proof is straightforward. Figure 8 shows an example of the merging process: the curls C_1 and C_2 are placed one next to the other as shown in Fig. 8, (a), and two of their neighbors points belonging to the same switching component (in Fig. 8 the points inside the red circle) exchange one of their coordinates.

We underline that the merging process can be iterated when more than two *hv*-convex curls are involved. Again the process ends up in one single *hv*-convex curl.

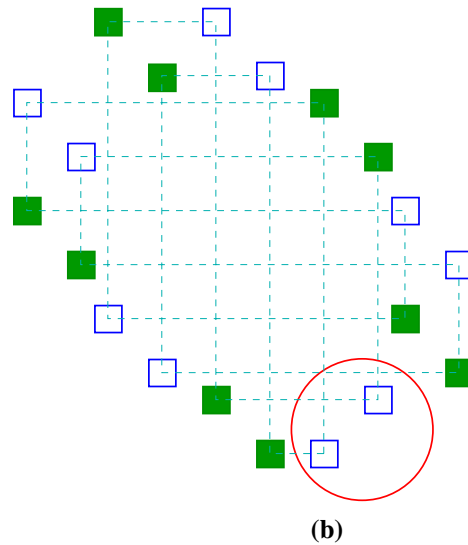
Example 29 Assume $a = 24$ and $b = 15$, so $CompSeq(24, 15)$ computes $d = 3$, updates a, b with 8 and 5, respectively, and exploits $CompSeqEntries(8, 5)$, which outputs $h = 1, \alpha = 3, \gamma = 5, N_\alpha = 2$, and $N_\gamma = 4$. Then, we have

$$I_\alpha = (2, 3] \cup (5, 6] \cup (15, 16] \cup (18, 19] \text{ and}$$

$$I_\gamma = [1, 2] \cup [7, 8] \cup [14, 15] \cup [20, 21].$$

The *for* loop runs from 1 to $a - b = 3$ setting the elements in π as follows:

$i = 1$: initializes $x = 1$ and, since $1 \in I_\gamma$, it sets $\pi(1) = 5$. Then, $y = (HV)^2(1) = 7$ and, since $i = 1$ is odd $x = V(7) = 1 - 7 = -6 \pmod{26} = 20$;



(b)

Fig. 8 Two distinct *hv*-convex curls each having sequence $(5)_2$ denoted by full green colored squares (bold dashed lines) and contour blue colored squares (thin dashed lines). After exchanging one of the coordi-

ates of the two neighbor points inside the circle, we merge them into one single curl that preserves the *hv*-convexity and whose sequence is $(5)_4$ (Color figure online)

$i = 2$: since $20 \in I_\gamma$, then $\pi(2) = 5$. Then, $y = (VH)^2(20) = 40 \pmod{26} = 14$ and, since $i = 2$ is even, $x = H(14) = 17 - 14 = 3$;
 $i = 3$: since $3 \in I_\alpha$, then $\pi(3) = 3$, $y = HV(3) = 19$ and the *for* loop stops.

The vector $\pi_3 = (5, 5, 3)$ is obtained, and, being $d = 3$, the final output is $\pi = (\pi_3)_6$

5 Concluding Remarks

In this paper, we have considered *hv*-convex switching components, namely the switching components under the prior knowledge that the object to be reconstructed is convex along the horizontal and the vertical directions. These are separated in two different geometric classes, windows and curls respectively (see [12,13]), and can be encoded as special sequences of integer numbers, called *hv*-sequences. We have provided a detailed investigation of *hv*-sequences, so obtaining a complete combinatorial description. In case of windows, the related *hv*-sequences are quite simple, since reduce to a single multiple of 4, that can be easily determined by the geometric structure of the switching component. Differently, for curls, *hv*-sequences consist of alternating strings of odd entries α, γ . The numerical values of α and γ , as well as the lengths of each string occurring in a given *hv*-sequence, can be exactly computed by exploiting two integer numbers a, b , that are related to the geometric structure of the curl (denoted by $Curl(a, b)$ for such a reason). A polynomial-time algorithm in a, b has been provided, which explicitly reconstructs the *hv*-sequence related to $Curl(a, b)$, and several examples have been shown and discussed in detail.

The knowledge of the combinatorial structure of *hv*-convex sequences, and consequently the geometric properties of the *hv*-convex switching components, is of main relevance for the faithful reconstruction of *hv*-convex polyominoes.

Furthermore, it is well known that small modifications of the projections may lead to dramatic changes in the reconstructed object. The reconstruction process is called stable if a small amount of noise can only lead to small differences in the reconstruction. So, it is of primary relevance to understand both the structure and the size of the switching components of a reconstructed object to estimate how far it is from the original one, when noisy projections are provided (the reconstruction process is called stable if a small amount of noise produces small differences in the reconstruction).

We feel that the obtained results provide additional knowledge on the geometric and combinatorial properties of the switching components, which could be exploited and incorporated in the tomographic problem in view of a reduction of the number of allowed solutions. Also, this could reveal of

interest when used in combination with the system of clauses involved in the P -time reconstruction process defined in [4], in order to possibly reduce them to a 2-SAT formula.

A final possible challenging problem consists in relating the geometrical and the combinatorial structure of the *hv*-convex switching components and sequences to the degree of convexity of the associated *hv*-convex polyominoes, according to the definition in [10]. In particular, it is interesting to investigate the possible co-existence of *hv*-convex windows and curls in a same *hv*-convex polyomino, aiming at deriving some limitations on the number and the type of allowed switching components (see also [21]), in view of a solution of the computational complexity problem of the reconstruction of digital convex polyomino from two projections.

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